THE NUMBER OF SEMIGROUPS OF ORDER n

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ABSTRACT. The number of semigroups on *n* elements is counted asymptotically for large *n*. It is shown that "almost all" semigroups on *n* elements have the following property: The *n* elements are split into sets *A*, *B* and there is an $e \in B$ so that whenever $x, y \in A, xy \in B$, but if x or y is in B, xy = e.

1. The problem. Fix a labelled *n*-element set $[n] = \{1, ..., n\}$. A semigroup SG on [n] is an associative binary operation (denoted by concatenation). Let S(n) denote the number of semigroups on [n]. We find an asymptotic approximation to S(n). Let

(1.1)
$$f(t) = \binom{n}{t} t^{1+(n-t)^2}.$$

THEOREM 1.

(1.2)
$$S(n) = \left[\sum_{t=1}^{n} f(t)\right](1+o(1)).$$

Define $t_0 = t_0(n)$ as that t which maximizes f(t). One can show

(1.3)
$$t_0 \sim n/(2 \ln n)$$

and f(t) has a sharp peak at t_0 . Equation (1.2) simplifies to

(1.4)
$$S(n) = f(t_0)(1 + o(1))$$

except in those "rare" instances when $f(t_0 + 1)$ or $f(t_0 - 1)$ are "near" $f(t_0)$.

For the remainder of the paper we adopt the convention that any inequality about functions of n is meant to be true only for all n sufficiently large, where how large depends on the statement.

2. Construction of the semigroups. Let $A \subseteq [n]$, |A| = n - t, $t \ge 1$. Set $B = A^0$. We construct a family S(A) of semigroups on [n] as follows:

(i) Select $e \in B$.

(ii) For $x, y \in A$ define xy to be an arbitrary member of B.

(iii) For $x \in B$ or $y \in B$ define xy = e.

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Then for all $x, y, z \in [n], xy \in B$ so (xy)z = e. Similarly, x(yz) = e. This yields

(2.1)
$$|S(A)| = t^{1+(n-t)^2}$$

semigroups. Call a semigroup SG of type T (for trivial) if it is of the above form for some A, e.

Let $S_1(A)$ denote those semigroups $SG \in S(A)$ such that for no $a \in A$ is ax = xa = e for all $x \in [n]$. Then

(2.2)
$$|S_1(A)| = (1 + o(1))|S(A)|$$

and the $S_1(A)$ are disjoint. Hence, there are $(\sum_{t=0}^n f(t))(1 + o(1))$ semigroups of type T. This implies

(2.3)
$$S(n) \ge \left[\sum_{t=0}^{n} f(t)\right](1+o(1)).$$

3. An upper-bound. Let $\mathfrak{T}(A)$ denote the family of semigroups for which A is a minimal (in cardinality) set of generators. (I.e. all $x \in [n]$ may be expressed as $x = a_1 \cdots a_k$ for some $k \ge 1, a_1, \ldots, a_k \in A$.) Set N(A) $= |\mathfrak{I}(A)|$. Then

(3.1)
$$S(n) \leq \sum_{A \subseteq [n]} N(A).$$

If $a_1, a_2 \in A$ then

$$(3.2) a_1 a_2 \in \{a_1, a_2\} \cup B,$$

as otherwise $A - \{a_1, a_2\}$ would be a smaller set of generators. We bound N(A)by constructing all $SG \in \mathfrak{T}(A)$ as follows:

(i) For $a_1, a_2 \in A$ choose $a_1 a_2 \in \{a_1, a_2\} \cup B$. (This may be done in $(t+2)^{(n-t)^2}$ ways.)

(ii) Let A' be a minimal (in cardinality) subset of A so that $A'A \supseteq AA \cap B$. For each $b \in AA \cap B$ there exist a_b' , a_b so that $b = a_b'a_b$. The set X $= \{a'_b : b \in B\}$ satisfies $XA \supseteq AA \cap B$, so that

$$(3.3) |A'| \leq |X| \leq |B| = t.$$

Since $A' \subseteq A$, $|A'| \leq \min(t, n - t)$.

Now for all $a \in A'$, $b \in B$ choose ab arbitrarily. (This may be done in at most $n^{|A'||B|} \leq n^{t \min(t,n-t)}$ wavs.)

Claim 1. SG is determined by $AA \cup A'B$.

PROOF. As all $x, y \in [n]$ are finite products of A's, it suffices to show that all $a_1 a_2 \cdots a_k$ are determined. We show this by induction, it being trivial for $k \leq 2$. For k > 2 set $a_1 a_2 = b$. If $b \in A$, $a_1 a_2 \cdots a_k = b a_3 \cdots a_k$ is determined by induction. If $b \notin A, b \in B \cap AA$ so b = cd for some $c \in A', d$ $\in A$. Then $a_1 \cdots a_k = c(da_3 \cdots a_k)$. Now $x = da_3 \cdots a_k$ is determined by induction and cx is determined since A'[n] is determined.

228

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(3.4)
$$N(A) \leq (t+2)^{(n-t)^2} n^{t \min(t,n-t)}.$$

Thus there are at most $\binom{n}{t}(t+2)^{(n-t)^2}n^{t\min(t,n-t)}$ semigroups having a minimal set of generators of cardinality t.

From (3.1), (3.4),

(3.5)
$$S(n) \leq \sum_{t=1}^{n} {\binom{n}{t}} (t+2)^{(n-t)^2} n^{t \min(t,n-t)}$$

so that

(3.6)
$$S(n) = [n/(2e + o(1))\ln n]^{n^2}.$$

Let $\varepsilon > 0$ be fixed and small ($\varepsilon = 10^{-6}$ will do). By elementary calculations and (2.3),

$$\sum' \binom{n}{t} (t+2)^{(n-t)^2} n^{t \min(t,n-t)} = o(S(n))$$

where \sum' runs over t, $|t - t_0| > \epsilon n/\ln n$. Thus: Almost all semigroups are in $\mathfrak{T}(A)$ for some A, |A| = n - t, $|t - t_0| \leq \epsilon n/\ln n$. We restrict our attention to t in this range for the duration of this paper.

4. The easy case. Our ultimate objective is to show

(4.1)
$$N(A) = t^{1+(n-t)^2}(1+o(1))$$

for |A| = n - t, $|t - t_0| < \epsilon n/\ln n$. In this section we show the corresponding result for a restricted class of semigroups.

We say that a semigroup in $\mathfrak{T}(A)$ has property E (for easy) if $a_1, a_2 \in A$ imply $a_1 a_2 \neq a_1$ and $a_1 a_2 \neq a_2$. Let $\mathfrak{T}^*(A)$ denote the subclass of $\mathfrak{T}(A)$ satisfying property E and $N^*(A) = |\mathfrak{T}^*(A)|$. We shall show

(4.2)
$$N^*(A) = t^{1+(n-t)^2}(1+o(1))$$

for |A| = n - t, $|t - t_0| < \epsilon n / \ln n$. For $SG \in \mathfrak{I}^*(A)$ set

$$G_{ax} = \{b \in B: ab = x\} \quad \text{for } a \in A, x \in [n].$$

Set

$$F(a) = \max_{x \in [n]} |G_{ax}|$$

and

$$\delta = t - \min_{a \in A} F(a).$$

The number δ defined above may be thought of as follows. Consider the rows aB, $a \in A$. Each row has a most frequent entry. δ gives a uniform upper bound (over $a \in A$) for the number of entries in aB not equal to the most frequent one. Trivial semigroups have $\delta = 0$. Let $\mathbb{T}^*_{\delta}(A)$ be those semigroups $\lim_{x \to a} \mathbb{T}^*_{\delta}(A)$ be those semigroups $\lim_{x \to a} \mathbb{T}^*_{\delta}(A)$ be those semigroups have $\delta = 0$.

229

We now bound $N_{\delta}^{*}(A)$. The semigroups in $\mathfrak{T}_{\delta}^{*}(A)$ may be constructed by the following procedure.

(i) Pick $g \in A$ so that $t - \delta = F(g)$. (This may be done in $n - t \leq n$ ways.) (ii) Set A' equal a minimal subset of A containing g so that $A'A \supseteq AA \cap B$. fix A'. (Since $|A'| \leq t$, this may be done in at most n^t ways.)

(iii) Determine A'B. For each $a \in A'$, $F(a) \ge t - \delta$; so there exists $x \in [n]$, $|G_{ax}| \ge t - \delta$. Now aB may be determined by selecting x (in at most n ways), G_{ax} (in $\sum_{i=t-\delta}^{t} {t \choose i} \le 1 + t^{\delta} \le n^{\delta}$ ways), and determining ab for $b \in B$ $-G_{ax}$ (in $n^{|B|-|G_{ax}|} \le n^{\delta}$ ways.) Thus each row has at most $n^{2\delta+1}$ possibilities, and therefore A'B has at most $n^{(2\delta+1)|A'|} \le n^{(2\delta+1)t}$ possibilities.

(iv) Determine gA. For each $a \in A$ we have $ga \in B$ (by property E). There are at most t^{n-t} possible gA.

(v) Define an equivalence class on A by $a \equiv b$ if ga = gb. Since $gA \subseteq B$ there are at most t equivalence classes and we select representatives $g = a_1$, ..., a_s , $s \leq t$, of the classes in an arbitrary fashion. For $2 \leq i \leq s$ we determine a_iA , each in at most t^{n-t} ways, for a total of $t^{(n-t)(s-1)}$ possibilities.

(vi) Determine the remainder of AA. Let $a, c \in A, a \neq a_1, \ldots, a_s$. Then $a \equiv a_i$ for some $i, 1 \leq i \leq s$. Then

$$g(ac) = (ga)c = (ga_i)c = g(a_ic)$$
 and $ac \in G_{g,g(a_ic)}$

Now $a_i c$ has already been determined, so $g(a_i c)$ has been determined, as has $G_{g,g(a,c)}$. From the definition of δ ,

$$|G_{g,g(a_ic)}| \leq t - \delta.$$

The product *ac* can be determined in at most $t - \delta$ ways. This holds for (n-t)(n-t-s) pairs (a, c), giving $(t-\delta)^{(n-t)(n-t-s)}$ possibilities.

By Claim 1 the semigroup is now determined.

(4.4)
$$N_{\delta}^{*}(A) \leq nn^{t} n^{(2\delta+1)t} t^{(n-t)s} (t-\delta)^{(n-t)(n-t-s)}$$

(4.5)
$$\leq n^{1+(2\delta+2)t}(1-\delta/t)^{n^2(1+o(1))}t^{(n-t)^2}$$

and

(4.6)
$$\sum_{\delta=1}^{t-1} N_{\delta}^{*}(A) = o\left[t^{1+(n-t)^{2}}\right]$$

by an elementary calculation (using $t \sim n/(2 \ln n)$). That is, almost all semigroups in $\mathfrak{T}^*(A)$ have $\delta = 0$. Intuitively, for $\delta \ge 1$, $N^*_{\delta}(A) = o(t^{(n-t)^2})$ because for most $a, c \in A$ the product ac can take at most $(t - \delta)$ values versus t in the $\delta = 0$ case.

Call an $SG \in \mathfrak{T}^*(A)$ absurd if $\delta = 0$ but is not trivial. Since $\delta = 0$ each aB, $a \in A$, is constant.

Claim 2. If $SG \in \mathfrak{T}^*(A)$ and AB = constant, then SG is trivial.

PROOF. Say AB = e. For any $a \in A$, $aa \in B$ so aaa = e. Now $e \notin A$, since, if it were, $A - \{e\}$ would generate e and, therefore, [n], contradicting the minimality of |A|. So $e \in B$. Let $b \in B$, $x \in [n]$. Since A generates [n], b License or copyright restrictions may apply to redistribution; see https://www.ams.org/ournal-terms-of-use

 $= a_1 \cdots a_s \ (s \ge 2), \ x = a_{s+1} \cdots a_{s+t} \ (t \ge 1)$ so $bx = a_1 \cdots a_{s+t} = e$ as $s + t \ge 3$. That is, B[n] = e so SG is trivial.

Claim 3. Let $SG \in \mathfrak{T}^*(A)$, $\delta = 0$, $a_1, a_2 \in A$, $a_1B = e_1 \neq e_2 = a_2B$. Then $a_1A \cap a_2A = \emptyset$.

PROOF. Suppose $x, y \in A$, $a_1 x = a_2 y$. For any $z \in A$, $xz, yz \in B$ so $e_1 = a_1(xz) = (a_1 x)z = (a_2 y)z = a_2(yz) = e_2$, a contradiction.

The conditions on absurd $SG \in \mathfrak{T}^*(A)$ imposed by Claim 3 are sufficiently stringent that we may easily show (details omitted) that they are "small" in number (i.e. $o(t^{1+(n-t)^2})$). Hence, almost all $SG \in \mathfrak{T}^*(A)$ are trivial, yielding (4.2).

5. The general case-outline. Theorem 1 is implied by (4.1). We give a brief outline of the proof of (4.1).

We let $\mathfrak{T}(A, \delta, ...)$ and $N(A, \delta, ...)$ denote the set, and number, of semigroups with parameters A, δ, \ldots . By $N(A, \delta, \ldots)$ "small" we always mean in comparison to $t^{(n-t)^2}$.

For a particular counting scheme let $\nu(a)$ denote the number of possible rows aA and $\mu(a) = \nu(a)/t^{n-t}$. By (3.2) all $\nu(a) \leq (t+2)^{n-t}$ so $\mu(a) \leq n^4$.

For $a \in A$ set

$$S_a = \{x \in A : ax = x\}$$
 and $L = \{a \in A : |S_a| \ge .01n\}, \quad l = |L|.$

If $a \in L$ there are less than 2^n choices for S_a and $n^{.99n}$ choices for $a([n] - S_a)$. There are less than $2^n n^{.99n} = t^{(.99+o(1))n}$ choices for a[n]. We may determine $SG \in \mathfrak{T}(A, L)$ by determining LA, then (A - L) A, then defining A' and determining A'B. If l is "large", N(A, L) is "small". Most $SG \in \mathfrak{T}(A)$ have l = o(n) -which we assume for the duration.

We modify the bounding of $N_{\delta}^{*}(A)$ in §4 by (4.4) to $N(A, L, \delta)$. We need a slightly different definition for δ :

(5.1)
$$\delta = t - \min_{a \in A - L} F(a).$$

We fix $g \in A - L$ so that $F(g) = t - \delta$. We define A' as the minimal set so that $(A' \cup L)A \supseteq AA \cap B$. We bound $N(A, L, \delta)$ by first determining L[n] and then following steps (i), ..., (vi) of §4. The equivalence classes, in step (v), are defined on A - L and there will be $s \leq t + 1 + .01n$ such classes (when ga = a, $\{a\}$ is an equivalence class). We find

(5.2)
$$N(A, L, \delta) \leq t^{(.99+o(1))nl} nn^{t} n^{(2\delta+1)t} \cdot (t+2)^{s(n-t)} (t+2-\delta)^{(n-t-l-s)(n-t)}.$$

For $\delta \ge 3$, $N(A, L, \delta)$ is "small". We assume $\delta \le 2$ for the duration. Note this implies $|aB| \le 3$ for all $a \in A - L$.

Set

$$I = \{a \in A - L: aa = a\},\$$

$$V = \{a \in A - L - I: \exists x \in A - L, a \neq x, ax = a \text{ or } xa = a\},\$$

$$|I| = i|V| = v.$$

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For $a \in I$, $x \in A - \{a\}$,

$$ax = (aa)x = a(ax) \in \{a, x\} \cup aB,$$

at most 5 possibilities, so $\nu(a) \leq 5^{n-t}$, $\mu(a) = t^{n(-1+o(1))}$.

Suppose $a \in V$, $x \in A - L$, ax = a or xa = a and xA has already been determined.

Case 1. xa = a. For all $c \in A$,

$$ac = (xa)c = x(ac) \in \{xa, xc\} \cup xB$$

at most 5 possibilities. Then $\nu(a) \leq 5^{n-t}$, $\mu(a) \leq t^{n(-1+o(1))}$. Case 2. ax = a. For $c \in A$, $xc \neq c$ we have

$$ac = (ax)c = a(xc) \in \{ax\} \cup aB,$$

at most 4 possibilities. We determine aB (at most n' ways), then aA. Since $x \notin L$, at most .01n c's have xc = c. For these there are n possible ac. Then $\nu(a) \leq n' 5^n n^{.01n} = n^{n(.01+o(1))}$, so $\mu(a) \leq t^{n(-.99+o(1))}$.

If *i* is large we may bound N(A, I) by determining *IA*, then (A - I)A, then A'A. We assume i = o(n) for the duration.

If v is large we may bound N(A, V) by determining (A - V)A, then VA, then A'A. (A technical problem arises. For $a \in V$, ax = a or xa = a, we want to determine aA after xA. But perhaps $x \in V$. One may order V (trying any ordering and its reverse) so at least half the $a \in V$ come after their respective x.) We may assume v = o(n) for the duration.

We bound $N(A, L, I, V, \delta)$, where l, i, v = o(n). We determine xA for $x \in L \cup I \cup V$ -and when xA is needed before $aA, a \in V$. We define equivalence classes on the remaining $x \in A$, determining xA first for x a representative, then for the remainder of A. We achieve an expression analogous to (5.2) with a factor of $(t - \delta)^{m(n-1)}$ where m, the number of "remaining" A, is at least (.99 + o(1))n. The expression is "small" for $\delta > 0$. We assume $\delta = 0$.

If $L \cup I \cup V \neq \emptyset$ we have factors $\mu(a) \leq t^{(-.99+o(1))n}$ that are not adequately counterbalanced so that $N(A, \delta, L, I, V)$ is "small". Most semigroups have $\delta = 0$, $L = I = V = \emptyset$ and are trivial.

(A final note on "filling in details". One shows $N(A, \delta, L, ...) \leq t^{-.99n} t^{(n-t)^2}$ and there are less than, say, 5^n possible $\delta, L, ...$ so that $\sum N(A, \delta, L, ...)$ taken over all $\delta, L, ...,$ except $\delta = 0, L = \cdots = \emptyset$, is small.)

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232