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## THE TOEPLITZ PACRAGE USERS' GUIDE*

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October 1983
*The work of the American participants was supported in part by the National Science Foundation and in part by the Applied Mathematical Sciences Research Program (KC-04-02) of the Office of Energy Research of the U.S. Department of Energy under Contract W-31-109-Eng-38. The work of the Soviet participants was supported by the State Committee for Science and Technology (U.S.S.R.).
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## TEE TOEPLITZ PACRAGE USEES' GUIDB

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#### Abstract

The TOEPLITZ package is a collection of Fortran subroutines for the numerical solution of systems of linear equations with coefficient matrices of Toeplitz or circulant form. This report provides a description of the algorithms and software in the package and includes program listings.


## INTRODUCTION

## 1. Overview of the TOEPLITZ Package

The TOEPLITZ package is a collection of Fortran subroutines for solving linear systems

$$
A x=b,
$$

where $A$ is a Toeplitz matrix (see subsection l.l of Chapter l), a circulant matrix (see subsection 1.2 of Chapter 1 ), or has one or several block structures based on Toeplitz or circulant matrices. Included also is capability for orthogonal factorization of a column-circulant matrix (see subsection 1.4 of Chapter 1).

Such systems arise in problems of electrodynamics, acoustics, mathematical statistics, algebra, in the numerical solution of integral equations with a difference kernel, and in the theory of stationary time series and signals (see, e.g., $[5,7,9,17,20,25,26]$ ). Circulant matrices play an important role in the theory of circular convolutions [13]. Block-Toeplitz matrices have recently begun to play a significant role as the applicability of multichannel time series increases $[22,30]$.

Although the theoretical and practical significance of Toeplitz matrices was recognized early in this century $[23,28,31]$, computational aspe tr were nol. studied until more recently. The most influential and fundamental paper on algorithmic aspects was Levinson's extension to the discrete case of

Wiener's basic work on filtering [19,29]. It was here that the technique of bordering and recursion on the order of the system was first shown to be an effective way to produce efficient algorithms for Toeplitz systems. Levinson's algorithm is an $\rho\left(M^{2}\right)$ method for solving an order $M$ positive-definite symmetric Toeplitz system of equations. Trench later used the same 1 deas to show how bordering could be exploited for general Toeplitz systems [24]. Trench's work was made more explicit and generalized by Zohar [32,33].

These $0\left(M^{2}\right)$ algorithms for Toeplitz systems are currently the most practical methods for such problems. They have simple descriptions as programs, they use simple storage and control structures, and ercor analyses are avallable for some of them $[8,10,11]$.

The algorithms in this package for circulant matrices appear to have been known classically (see [13]). Toeplitz matrices of the second level are discussed in $[4,21,22,27]$; the algorithms are essentially the same as those in this package.

Toeplitz matrices arising in time series and signal processing are quite often covariance matrices that occur in normal equations for linear leastsquares problems. The coefficient matrices in these problems often have column-circulant structures that lend themselves to efficient methods for problem solution by orthogonalization. These methods are usually called "lattice methods" in the signal processing iiterature [12,14,18]; one such method [i2] is implemented in the TOEPLITZ package.

The TOEPLITZ package has an intentional similarity to LINPACK [15] in the format of the Fortran source, in the comments, and in the subroutine naming conventions. All names consist of four, five, or six letters (depending on the level of block structure of the matrix A) in the forms XSL\#, XYSL\#, or XYZSL\# for the system solving subroutines and CQR\# for the orthogonal factorization subroutines.* When A has no special block structure (see Chapter 1), the letter in the $X$ position specifies the type of the matrix:

T Toeplitz
C Circulant.

[^0]When A has a two-level block structure (see Chapter 2), the letters in the XY positions specify the type of the matrix:

TG Block-Toeplitz where the blocks are general matrices
CT Block-circulant where the blocks are Toeplitz matrices
CC Block-circulant where the blocks themselves are circulant matrices

CG Block-circulant where the blocks are general matrices.

When $A$ has a three-level block structure (see Chapter 3), the letters in the XYZ positions specify the type of the matrix:

CTG Block-circulant where the blocks are two-level TG-type matrices
CCT Block-circulant where the blocks are two-level CT-type matrices
CCC Block-circulant where the blocks are two-level CC-type matrices
CCG Block-circulant where the blocks are two-level CG-type matrices.
By permuting corresponding rows and colimns, one can transform any twolevel XY-type matrix to YX-type (see Tyrtyshnikov [25]). Similarly, one can interchange any two levels of a three-level XYZ-type matrix. These circumstances effectively extend the capability of the TOEPLITZ package to additional matrix types.

The fixed letters $S L$ indicate that the routine solves a linear system, while the letters $Q R$ indicate that the routine performs an orthogonal factorizatica.

The last letter in the \# position specifies the matrix data type. Standard Fortran allows the use of three such types:

$$
\begin{array}{ll}
\text { S } & \text { REAL } \\
\text { D } & \text { DOUBLE PRECISION } \\
\text { C } & \text { COMPLEX. }
\end{array}
$$

In addition, some Fortran systems allow a double precision complex type:
2 DOUBLE COMPLEX.

## 2. The Leading Array Dimension Parameter

Those members of the TOEPLITZ package that process a two-dimensional array include in their calling sequences the parameter LDA (or LDQ,LDS) to
communicate the leading dimension of the array. "Leading dimension" refers to the DIMENSION statement storage allocation for the array and should be distinguished from the order of the linear system. The inclusion of this parameter enables flexibility in processing systems of varying order without the bother of changing the DIMFNSION statement for the coefficient matrix.

For example, if the array $A$ has been declared " $A(50,20)$ " in the DIMENS ION statement, then simply enter the statement "LDA $=50^{\prime \prime}$ into the body of the program before the call to the TOEPLTTZ package subroutine.

## 3. Development of the TOXPLITZ Package

In offering the TOEPLITZ package to the international computing communty, it is appropriate to note that this software is the result of collaboratior among scientists in the United States and the Soviet Union. Hence, in addition to the intrinsic usefulness of the package, the software in its present form demonstrates the possibilities inherent in Soviet-American collaboration in the development of scientific software. The work was caried out under the auspices of the agreement between the U.S.A. and the U.S.S.R. on Scientific and Technological Cooperation in the field of Application of Computers to Economics and Management, subtopic Mathematical Software.

This collaborative effort was initiated at the Numerical Software Workshop which took place at the National Science Foundation (NSF) in Washington, D.C. in December of 1975. The general framework of joint efforts was discussed during that workshop by D. Aufenkamp of NSF, W. Cody of the Applied Mathematics Division, Argonne National Laboratory (AMD-ANL), and O. Arushanian of the Science Research Computing Center, Moscow State University (SRCC-MSU), then visiting Pennsylvania State University for the year. Further steps were discussed during a meeting which took place at Penn State in February of 1976 involving D. Aufenkamp (NSF), J. Boyle (AMD-ANL), W. Cowell (AMD-ANL), and 0. Arushanian (SRCC-MSU), and during a short visit by 0 . Arushanian to J. Bunch, University of California at San Diego (UCSD). In accordance with plans agreed upon during these meetings and approved in the meeting of coordinators and experts on the topic "Theoretical Foundations of Software for Application in Economica and Management" which took place in Moscow in June of 1976, Long-term visits of American scientista to the U.S.S.R. in 1976 and 1978 and of Soviet acientiats to the U.S.A. In 1978 and 1979 were arranged to
exchange information and to carry out joint work on numerical software development. These joint efforts came to be known as the SALAR (Soviet-American Libraries and Algorithms Research) project. Results of accomplished works have appeared in 25 papers (see [1] and [2]) and were presented at the IFIP Congress in August of 1977 in Toronto, Canada (see [3]).

The contributions from the U.S.A. side were made by J. Boyle, K. Dritz, W. Cowell, and $B$. Garbow of $A M D-A N L$ (now redesignated MCS-ANL), J. Bunch of UCSD, D. Sorensen (now of MCS-ANL), W. Miller (now of the University of Arizona), and $C$. Moler of the University of New Mexico. The contributions from the U.S.S.R. side were mace by V. Voevodin (now of the Academy of Sciences, State Committee for Science and Technology), 0. Arushanian, M. Samarin, E. Nikolaev, V. Morozov, Y. Kuchevskiy, E. Tyrtyshifkov, N. Bogomolov, and V. Borisov of SRCC-MSU.

The SALAR project had a number of objectives. First of all, it represented joint research into the methodology and practical aspects of producing mathematical software, namely, numerical libraries and packages. This main objective dictated the necessity of also investigating systems aspects of mathematical software development, which include the study of transportability problems, tailoring of programs to user requests, abstract formulation of numerical algorithms, and program transformation and generation systems. Methodological questions associated with the joint systematization, testing, and certification of mathematical software packages were also of great importance in the SALAR project. Research in numerical algorithms development was conducted mostly in linear algebra on problems such as updating algorithms for matrix decomposition and solving special types of linear systems.

The TOEPLITZ package was produced as a part of the SALAR project and can be considered as a practical result of previous investigations. The routines were originally written in 1978 at Moscow State University by E. Tyrtyshnikov [25] on the basis of the theoretical results of W. Trench [24] and S. Voevodina [27], and on his own resea ch. A preliminary version of the users' guide was written by Soviet and American scientists during a visit to Argonne National Laboratory (U.S.A.) made by Soviet scientists O. Arushanian and M. Samarin (of SRCC-MSU) in 1979. Multiple versions of TOEPLITZ subroutines and formatting of codes were obtained with the help of the TAMPR-system [3], produced by J. Boyle and K. Dritz of AMD-ANL. Modifications, commenting,
and test driver design were also accomplished during this Argonne visit. Scientific supervision over the development of the TOEPLITZ package at SRCCMSU was provided by V. Voevodin.

Further developmental work on the codes and preparation of this users' gaide were accomplished at Argonne in 1982. The added capability for orthogonalization of column-circulant matrices derives from a new algorithm of G. Cybenko [12] (of Tufts University). Cybenko also suggested an improved formulation of another of the algorithms, supplied background information included in the "Overview" section of this guide, and pointed us to many of the references.

In conclusion we wish to acknowledge the support of the National Science Foundation (U.S.A.) and the State Committee for Science and Technology (U.S.S.R.), executors of the Science and Technology Agreement. Special thanks are due to D. Aufenkamp (U.S.A.), B. Rameev (II.S.S.R.), and Y. Baraboshkin (U.S.S.R.) who created conditions in which our joint work could flourish. We also express our great gratitude : Judy Beumer (of MCS-ANL) who carefully typed the manuscript for ${ }^{\text {ti is }}$ usels' guide.

## 4. Availability of the TOEPL iTZ Fackage

The TOEPLITZ package is available on tape from the following sources.

| National Energy Software Center | IMSL, Inc. |
| :--- | :--- |
| Argonne National Laboratory | or $\operatorname{Slxth}$ Floor, NBC Kldg. |
| 9700 South Cass Avenue | 7500 Bellaire Blvd. |
| Argonne, IL 60439 | Houston, TX 77036-5085 |
| Phone: (312) 972-7250 | Phone: (713) 772-1927 |

The package includes both single precision and double precision versions of the programs, and testing aids are also provided on the tape (see The TOEPLITZ Package Implementation Gide, ANL 83-17).

Comments and questions regarding the TOEPLITZ package should be directed to

```
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```


## CHAPTER 1: TOEPLITZ AND CIRCUIANT MATRICBS

1. Stinicture and Representation

### 1.1. Toeplitz matrices (T-matrices)

A Toeplitz matrix, or $T$-matirix, $A$ is a real or complex square matrix whose elements along the main diagonal and along each co-diagonal are equal; thus $A$ has the representation

$$
A=\left(\begin{array}{ccccc}
a_{0} & a_{1} & a_{2} & \cdots & a_{M-1} \\
a_{-1} & a_{0} & a_{1} & \cdots & a_{M-2} \\
a_{-2} & a_{-1} & a_{0} & \cdots & a_{M-3} \\
\cdot & \cdot & \cdot & \cdots & \cdot \\
a_{-M+1} & a_{-M+2} & a_{-M+3} & \cdots & a_{0}
\end{array}\right)
$$

A $T$-matrix is completely specified by its first row and column.
In the TOEPLITZ package a $T$-matrix of order $M$ is represented by a singiy subscripted array of $2 * M-1$ elements which contains the first row of the matrix followed by its first colunn beginning with the second element:

$$
a_{0}, a_{1}, a_{2}, \ldots, a_{M-1}, a_{-1}, a_{-2}, \ldots, a_{-M+1}
$$

### 1.2. Gircilant Eatrices (C-matrices)

A circulant matrix, or C-matrix, A is a T-watrix, limited here to complex mode, with the further property that

$$
a_{-1}=a_{M-1}, \quad 1=1,2, \ldots, M-1 ;
$$

thus $A$ has the -epresentation

$$
A=\left(\left.\begin{array}{ccccc}
a_{0} & a_{1} & a_{2} & \cdots & a_{M-1} \\
a_{M-1} & a_{0} & a_{1} & \cdots & a_{M-2} \\
a_{M-2} & a_{M-1} & a_{0} & \cdots & a_{M-3} \\
\cdot & \cdot & \cdots & \cdots & \cdot \\
a_{1} & a_{2} & a_{3} & \cdots & a_{0}
\end{array} \right\rvert\,\right.
$$

A C-matrix is completely specified by its first row; each further row may be obtained from the previous one by a right cyclic siift.

In the TOEPLITZ package a C-matrix of order $M$ is represented by a singly subscripted array of $M$ elements which contains the first row of the matrix:

$$
a_{0}, a_{1}, a_{2}, \ldots, a_{M-1} .
$$

### 1.3. General matrices (G-matrices)

A general real or complex square matrix

$$
A=\left(\begin{array}{ccccc}
a_{11} & a_{12} & a_{13} & \cdots & a_{1 M} \\
a_{21} & a_{22} & a_{23} & \cdots & a_{2 M} \\
a_{31} & a_{32} & a_{33} & \cdots & a_{3 M} \\
\cdots & \cdots & \cdots & \cdots & \cdots \\
a_{M 1} & a_{M 2} & a_{M 3} & \cdots & a_{M M}
\end{array}\right)
$$

will be called a G-matrix.
In the TOEPLI'sZ package a G-matrix of order $M$ is represented by a singly subscripted array of $M * * 2$ elements which contains the successive columns of the matrix:

$$
a_{11}, a_{21}, a_{31}, \ldots, a_{M 1}, a_{12}, a_{22}, a_{32}, \ldots, a_{M 2}, \ldots,{ }_{1 M}, a_{2 M}, a_{3 M}, \ldots, a_{M M} .
$$

### 1.4. Colum-circulant matrices

The designation "column-circulant" will be given to a real or complex rectangular matrix $A$, with row order $M$ at least equal to its column oider $L$, whose first column is specified and each further column obtained from its predecessor by a downward cyclic shift; thea nas the representation

$$
A=\left(\begin{array}{ccccc}
a_{0} & a_{M-1} & a_{M-2} & \cdots & a_{M-L+1} \\
a_{1} & a_{0} & a_{M-1} & \cdots & a_{M-L+2} \\
a_{2} & a_{1} & a_{0} & \cdots & a_{M-L+3} \\
\cdot & \cdot & \cdot & & \vdots \\
\cdot & \cdot & \cdot & & \cdot \\
a_{M-1} & a_{M-2} & a_{M-3} & \cdots & \cdots \\
a_{M-L}
\end{array}\right)
$$

In the TOEPLITZ package a column-circulant mairix with $M$ rows is represented by a singly subscripted array of $M$ elements which contains the first column of the matrix:

$$
a_{0}, a_{1}, a_{2}, \ldots, a_{M-1}
$$

## 2. Solution of Linear Bquations with T-Matrices

### 2.1. Purpose

The TOEPLITZ subroutines in this section are designed to solve linear algebraic equations with $T$-matrices. Usage will be described for the single precision real version. Double precision, complex, and double precision complex versions are also avallable. Indeed, the complex version is called in solving two-level CT-matrix systems (see subsection 3.5 of Chapter 2).

### 2.2. Usage

Single precision real T-matrices. TSLS solves a linear system with a real Toeplitz matrix. The calling sequence is

CALL $\operatorname{TSLS}(A, X, R, M)$.

On entry,

A is a singly subscripted array of $2 \star \mathrm{M}-1$ elements which contains the first row of the $T$-matrix followed by its first column beginning with the second element. A is unaltered by TSLS.
$X$ is a singly subscripted array of $M$ elements which contains the right hand side of the system.
$R \quad$ is a singly subscripted array of $2 * M-2$ elements used for work space.
$M$ is the order of $A$ and the number of elements in $X$.

On return,
$X$ contains the solution of the system.

Double precision real T-matrices. The calling sequence of the double precision real T-matrix subroutine TSLD is the same as that of TSLS with A, X, and $R$ DOU3LE PRECISION variables.

Single precision complex T-matrices. The calling sequence of the single precision complex T-matrix subroutine TSLC is the same as that of TSLS with A, $X$, and $R$ COMPLEX variables.

Double precision complex T-matrices. In those computing systems where it is available, the calling sequence of the double precision complex $T$-matrix subroutine TSLZ is the same as that of TSLS with $A, X$, and $R$ DOUBLE COMPLEX variables.

### 2.3. Example

The following program segment illustrates the use of the single precision subroutine TSLS for real T-matrices. Examples of the use of TSLD, TSLC, and TSLZ could be obtained by changing the subroutine name and type declaration. The system is of order 4 with coefficiencs as follows.

$$
A=\left(\left.\begin{array}{llll}
1 & 2 & 3 & 4 \\
5 & 1 & 2 & 3 \\
6 & 5 & 1 & 2 \\
7 & 6 & 5 & 1
\end{array} \right\rvert\, \quad X=\left(\left.\begin{array}{l}
10 \\
11 \\
14 \\
19
\end{array} \right\rvert\,\right.\right.
$$

```
    REAL A(7),X(4),R(6)
    INTEGER M,I
    DATA A(1)/1.0/,A(2)/2.0/,A(3)/3.0/,A(4)/4.0/,
    * A(5)/5.0/,A(6)/6.0/,A(7)/7.0/
    DATA X(1)/10.0!,X(2)/11.0/,X(3)/14.0/,X(4)/19.0/
    M = 4
    CALL TSLS(A,X,R,M)
DO 10 I = 1, M
    WRITE(...,...) X(I)
10 CONTINUE
    STOP
    END
```

The solution of the system is

$$
X=(1.0,1.0,1.0,1.0)
$$

### 2.4. Algorithm

The algorithm for the solution of a system of linear algebraic equations

$$
\begin{equation*}
A x=b \tag{1}
\end{equation*}
$$

with a T-matrix A of order $M$ comprises a sequence of $M$ steps. At the (k+1)-st step the solution of the system

$$
\begin{equation*}
A_{k} y_{k}=d_{k} \tag{2}
\end{equation*}
$$

is determined. Here

The vector $y_{k}$ is calculated by recurrence from $y_{k-1}$. The final result of the recurrent process is the solution of system (1), namely, $x=y_{M-1}$.

At step $1, y_{0}=b_{0} / a_{0}$. At step $k+1$, let us consider the unknown vector $y_{k}$ to be the sum of two vectors, one of which, augmented by a zero, was determined at the $k-t h$ step:

$$
\left(\begin{array}{c}
y_{0, k}  \tag{3}\\
y_{1, k} \\
\vdots \\
y_{k-1, k} \\
y_{k, k}
\end{array}\right)=\left(\begin{array}{c}
y_{0, k-1} \\
y_{1, k-1} \\
\vdots \\
y_{k-1, k-1} \\
0
\end{array} \left\lvert\,+\left(\begin{array}{c}
z_{0, k} \\
z_{1, k} \\
\vdots \\
z_{k-1, k} \\
z_{k, k}
\end{array}\right)\right.\right.
$$

Substituting this sum into equation (2) and taking into account that the vector $y_{k-1}$ satisfies the equation

$$
A_{k-1} y_{k-1}=d_{k-1},
$$

we see that the unknown vector $z_{k}$ from (3) with elements

$$
z_{0, k}, z_{1, k}, \ldots, z_{k, k}
$$

is the solution of the system

$$
A_{k} z_{k}=f_{k},
$$

where

$$
f_{k}=\left(\left.\begin{array}{c}
0 \\
0 \\
\vdots \\
0 \\
f_{k, k}
\end{array} \right\rvert\,, \quad f_{k, k}=b_{k}-\sum_{\ell=1}^{k} a_{-\ell} y_{k-\ell, k-1}\right.
$$

Thus, the vector $z_{k}$ is the same as the last column of the matrix $A_{k}^{-1}$ multiplied by $\mathbf{f}_{k, k}$. Hence, for recurrent calculation of the vectors $y_{k}$ it is sufficient to evaluate recurrently the last column of the matrix $A_{k}^{-1}$, or as done here for further economy an appropriately chosen multiple of this column. It is here that advantage is takcn of the Toeplitz structure of $A$.

Let us denote by $g_{k}$ and $h_{k}$ the first and last columns, respectively, each scaled by the as yet unspecified factor $q_{k}$, of the matrix $A_{k}^{-1}$ :

$$
g_{k}=\left(\begin{array}{c}
g_{0, k} \\
g_{1, k} \\
\vdots \\
g_{k-1, k} \\
g_{k, k}
\end{array}\right), \quad h_{k}=\left(\begin{array}{c}
h_{0, k} \\
h_{1, k} \\
\vdots \\
h_{k-1, k} \\
h_{k, k}
\end{array}\right), \quad \text { and } \quad A_{k} g_{k}=\left(\begin{array}{c}
q_{k} \\
0 \\
\vdots \\
n \\
0
\end{array} \left\lvert\,, \quad A_{k} h_{k}=\left(\begin{array}{c}
0 \\
0 \\
\vdots \\
0 \\
q_{k}
\end{array}\right)\right.\right. \text {. }
$$

It is clear that when $k=0$ the unscaled vectors colncide and contain the single element $1 / a_{0}$; we choose $q_{0}=a_{0}$ so that $g_{0}=h_{0}=1$. We will determine $g_{k}, h_{k}$, and $q_{k}$ from $g_{k-1}, h_{k-1}$, and $q_{k-1}$ using the following two sums:

$$
\begin{aligned}
& g_{k}=\left(\begin{array}{c}
g_{0, k-1} \\
g_{1, k-1} \\
\vdots \\
g_{k-1, k-1} \\
0
\end{array} \left\lvert\,+v\left(\left.\begin{array}{c}
0 \\
h_{0, k-1} \\
\vdots \\
h_{k-2, k-1} \\
h_{k-1, k-1}
\end{array} \right\rvert\,,\right.\right.\right. \\
& h_{k}=r\left|\begin{array}{c}
g_{0, k-1} \\
g_{1, k-1} \\
\vdots \\
g_{k-1, k-1} \\
0
\end{array}\right|+\left(\begin{array}{c}
0 \\
h_{0, k-1} \\
\vdots \\
h_{k-2, k-1} \\
h_{k-1, k-1}
\end{array}\right)
\end{aligned}
$$

where $v$ and $r$ are unknown scalars which we are going to derive.
Since $g_{k}$ and $h_{k}$ are columns of the matrix $A_{k}^{-1}$ scaled by $q_{k}$, then

$$
\begin{aligned}
& \left.A_{k} g_{k}=A_{k} \left\lvert\, \begin{array}{c}
g_{0, k-1} \\
g_{1, k-1} \\
\vdots \\
g_{k-1, k-1} \\
0
\end{array}\right.\right)+v A_{k}\left(\begin{array}{c}
0 \\
h_{0, k-1} \\
\vdots \\
h_{k-2, k-1} \\
h_{k-1, k-1}
\end{array}\right)=\left(\begin{array}{c}
q_{k} \\
0 \\
\vdots \\
0 \\
0
\end{array}\right) \text {, } \\
& A_{k} h_{k}=r A_{k}\left(\begin{array}{c}
g_{0, k-1} \\
g_{1, k-1} \\
\vdots \\
g_{k-1, k-1} \\
0
\end{array}\right)+A_{k}\left(\begin{array}{c}
0 \\
h_{0, k-1} \\
\vdots \\
h_{k-2, k-1} \\
h_{k-1, k-1}
\end{array}\right)=\left(\begin{array}{c}
0 \\
0 \\
\vdots \\
0 \\
q_{k}
\end{array}\right) .
\end{aligned}
$$

These relationships reduce to the following equations for determining the unknown scalars:

$$
\left\{\begin{array} { l } 
{ q _ { k - 1 } + f _ { 2 } v = q _ { k } }  \tag{4}\\
{ f _ { 1 } + q _ { k - 1 } v = 0 }
\end{array} \quad \left\{\begin{array}{l}
q_{k-1} r+f_{2}=0 \\
f_{1} r+q_{k-1}=q_{k}
\end{array}\right.\right.
$$

where

$$
\mathrm{f}_{1}=\sum_{\ell=1}^{k} \mathrm{a}_{-\ell} \mathrm{g}_{\mathrm{k}-\ell, \mathrm{k}-1}, \quad \mathrm{f}_{2}=\sum_{\ell=1}^{\mathbf{k}} \mathrm{a}_{\ell}^{\mathrm{h}}{ }_{\ell-1, k-1} .
$$

Solving equations (4) we find

$$
v=-f_{1} / q_{k-1}, \quad r=-f_{2} / q_{k-1}, \quad q_{k}=q_{k-1}-f_{1} f_{2} / q_{k-1}
$$

Note that this algorithm for solving linear systems with $T$-matrices requires that $A_{k}$ be non-singular for all $k$.

### 2.5. Prograning details - subroutine TSLSL

The calling sequence of subroutine TSLS is consistent with those of the other TOEPLITZ subroutines. However, it proves convenient in the implementation to consider the input matrix as two arrays and to partition the work space. Therefore, subroutine TSLSl was produced to directiy implement the algorithm, and subroutine TSLS merely acts as a user interface that calls TSLSl. TSLSl may be called directly by the user, if desired.

The calling sequence of subroutine TSLSl is
CALL TSLSI(Al, A2, B, X,Cl, C2, M) .
On entry,

Al is a singly subscripted array of $M$ elements which contains the first row of the $T$-matrix. Al is unaltered by TSLSI.

A2 is a singly subscripted array of $M-1$ elements which contains the first column of the $T$-matrix beginning with the second element. A? is unaltered by TSLSl.
$B$ is a singly sutscripted array of $M$ elements which contains the right hand side of the system. B ls unaltered by TSLSL.

Cl, C2 are singly subscripted arrays of $M-1$ elements used for work space.
$M$ is the order of the $T$-matrix and the number of elements in $B$ and $X$.

On return,
$X$ is a singly subscripted array of $M$ elements which contains the solution of the systel. $X$ may coincide with $B$.

Subroutine TSLS has double precision, complex, and double precision complex versions with names TSLDl, TSLCl, and TSLZl, respectively, whose calling sequences are the same as that of TSLS1 with $A 1, A 2, B, C l, C 2$, and $X$ variables of the corresponding type.

Towards timing estimation, nete that the algorithm for solving linear systems with $T$-matrices requires approximateiy $3 M^{2}$ multiplications.

### 2.6. Additional inforeation

The calling sequences of subroutines TSLS and TSLS for the solution of linear systems with $T$-matrices limit the right hand sides to single column vectors. There may be situations where the solutions of two or more such systems with the same coefficient matrix are desired. In these situations, modifications of the subroutines that would permit all solutions to be obtained in a single step could markedly improve efficiency. Fortunately, the algorithm organization for $T$-matrices enables such modifications to be made with little effort.

Three changes need to be made: 1) The parameter list mist be extended to include the column order of $X$ and $B$, and the leading dinension for these newly created two-dimensional arrays; 2) References to $X$ and $B$ must be rendered twodimensional; and 3) DO loops must be introduced for cycling over the columns of $X$ and $B$. Resulting form of TSLS and TSLSl are given below and can be compared with the official versions listed in Appendix $B$; to facilitate the comparison, the changes are indicated in lower case. The identical changes could be made to the double precision, complex, and double precision complex versions of these subroutines.

SUBROUTINE TSLS (A, X,R,M,mcol, 1dx)
INTEGER M,mcol, 1 dx
REAL A(1), X(ldx,mcol),R(1)
TSLS CALLS TSLSI TO SOLVE THE REAL LINEAR SYSTEM
A * $\mathrm{X}=\mathrm{B}$
WITH THE T - MATRIX A .
ON ENTRY
A REAL (2*M - 1)
THE FIRST ROW. OF THE T - MATRIX FOLLOWED BY ITS
FIRST COLUMN BEUINNING WITH THE SECOND ELEMENT .
ON RETURN A IS UNALTERED .
X REAL(M,mcol)
THE RIGHT HAND SIDE matrix B .
REAL (2*M - 2)
A WORK VECTOR .
INTEGER
THE ORDER OF THE MATRIX A.
mcol irteger
the number of columns of the matrices $x$ and $b$.
ldx integer
the leading dimension of the array x .
ON RETURN
X THE SOLUTION matrix .
SUBROLTINES AND FUNCTIONS
TOEPLITZ PACKAGE ... TSLS1
CALL SUBROUTINE TSLS1
CALL TSLS $1(A, A(M+1), X, X, R, R(M), M, m c o l, 1 d x)$
RETUKN
END
SUBROUTINE TSLS1(A1, A2, $\mathrm{B}, \mathrm{X}, \mathrm{C} 1, \mathrm{C} 2, \mathrm{M}, \mathrm{mcol}, \mathrm{ldx}$ )
INTEGER M,mcol,ldx
REAL A1(M), A2(1), B(1dx acol), X(ldx,mcol),C1(1),C2(1)
TSLS 1 SOLVES THE REAL LINEAR SYSTEM
$\mathrm{A} * \mathrm{X}=\mathrm{B}$
WITH THE T - MATRIX A .
ON ENTRY
A1 REAL(M)
THE FIRST R:'! ! F THE T - MATRIX A.
ON RETURN A1 IS UNALTERED .
A2 $\operatorname{REAL}(M-1)$
THE FIRST COLUMN OF THE T - MATRIX A
BEGINNING WITH THE SECOND ELEMENT.
ON RETURN A2 IS UNALTERED .
B $\quad \operatorname{REAL}(\mathrm{M}, \mathrm{mcol})$
THE RIGHT HAND SIDE matrix .
ON RETURN B IS UNALTERED.
C1 $\operatorname{REAL}(M-1)$
A WORK VECTOR .
C2 REAL (M - 1)
A WORK VECTOR .
M INTEGER
THE ORDER OF THE MATRI: A .
mcol integer
the number of columns of the matrices $x$ and $b$.
ldx integer
the leading dimension of the arrays $x$ and $b$.
ON RETURN
$\mathrm{X} \quad$ REAL(M,mcol)
THE SOLLTION matrix. X MAY COINCIDE WITH B .

INTERNAL VARIABLES
INTEGER I1, I2, $\mathrm{j}, \mathrm{N}, \mathrm{N} 1, \mathrm{~N} 2$
REAT, R1,R2,R3,R5,R6
SOLVE THE SYSTEM WITH THE PRINCIPAL MINOR OF ORDER 1 .
$R 1=A 1(1)$
do $5 \mathrm{j}=1$, mcol
$X(1, j)=B(1, j) / R 1$
5 continue
IF (M.EQ. 1) GO TO 80
RECURRENT PROCESS FOR SOLVING THE SYSTEM
WITH THE $T$ - MATRIX FOR $N=2, M$.
DO $70 \mathrm{~N}=2, \mathrm{M}$
COMPUTE MULTIPLES OF THE FIRST AND LAST COLUNNS OF THE INVERSE OF THE PRINCIPAL MINOR OF ORDER N.
$\mathrm{N} 1=\mathrm{N}-1$

```
        N2 = N - 2
        R5 = A2(N1)
        R6 = A1(N)
        IF (N .EQ. 2) GO TO 20
        C1(N1) = R2
        DO 10 I1 = 1, N2
            I2 = N - I1
            R5 = R5 + A2(I1)*C1(I2)
            R6 = R6 + A1(I1+1)*C2(I1)
        CONTINUE
    10
    20
    CONTINUE
        R2 = -R5/R1
        R3 = -R6/R1
        R1 = R1 + R5:*R3
        IF (N .EQ. 2) GO TO 40
        R6 = C2(1)
        C2(N1) = 0.0E0
        DO 30 Ii = 2, N1
            R5 =. C2(I1)
            C2(I1) = C1(I1)*R3 + R6
            Cl(I1) = Cl(I1) + R6*R2
            R6 = R5
        CONTINUE
    30
    40 CONTINUE
        C2(1) = R3
C
C COMPUTE THE SOLUTION OF THE SYSTEM WITH THE
C PRINCIPAL MINOR OF ORDER N .
C
    do 65 j = 1, mcol
        R5 = 0.0E0
        DO 50 I1 = 1, N1
            I2 = N - I1
            R5 = R5 + A2(I1)*X(I2,j)
        CONTINUE
        R6 = (B(N,j) - R5)/R1
        DO 60 I1 = 1, N1
        X(I1,j) = X(I1,j) + C2(I1)*R6
        CONTINUE
        X(N,j) = R6
    continue
    CONTINUE
    CONTINUE
        RETURN
    END
```


## 3. Solution of Linear Equations with C-Matrices

### 3.1. Purpose

The TOEPLITZ subroutine in this section is designed to solve linear algebraic equations with $C$-matrices; it is limited to complex matrices because the algorithm employs complex arithmetic. Users with real circulant matrices can either declare them complex or consider them simply T-matrices and employ the subroutines of section 2 . Running times as real T-matrices are shorter, but the unitary transformations employed in the algorithm described below for C-matrices offer greater stability. A double precision version of the subroutine is also available.

### 3.2. Usage

Single precision C-matrices. CSLC solves a linear system with a conplex circulant matrix. The calling sequence is
$\operatorname{CALL} \operatorname{CSLC}(A, X, R, M) \cdot$
On entry,

A is a singly subscripted array of $M$ elements which contains the first row of the $C$-matrix. A is unaltered by CSLC.
$X \quad$ is a singly subscripted array of $M$ elements which contains the right hand side of the system.
$R \quad$ is a singly subscripted array of $M$ elements used for work space.
$M$ is the order of $A$ and the number of elements in $X$.

On return,
$X$ contains the solution of the system.

Double precision C-matrices. In those computing systems where it is available, the ca:ling sequence of the double precision C-matrix subroutine CSLZ is the same as that of CSLC with $A, X$, and $R$ DOUBLE COMPLEX variables.

### 3.3. Example

The following program segment illustrates the use of the single precision subroutine CSLC for C-matrices. An example of the use of CSLZ could be obtained by changing the subroutine name and type declaration. The system is of order 4 with seefficients as follows.

$$
A=\left(\begin{array}{cccc}
1+1 & 2+2 i & 3+31 & 4+4 i \\
4+4 i & 1+i & 2+2 i & 3+3 i \\
3+3 i & 4+4 i & 1+1 & 2+2 i \\
2+2 i & 3+3 i & 4+4 i & 1+1
\end{array}\right) \quad X=\left(\begin{array}{c}
10+101 \\
10+101 \\
10+101 \\
10+101
\end{array}\right)
$$

```
    COMPLEX A(4),X(4),R(4)
    INTEGER M,I
    DATA A(1)/(1.0,1.0)/,A(2)/(2.0,2.0)/,A(3)/(3.0,3.0)/,
* A(4)/(4.0,4.0)/
    DATA X(1)/(10.0,10.0)/,X(2)/(10.0,10.0)/,X(3)/(10.0,10.0)/,
    * X(4)/(10.0,10.0)/
    M=4
    CALL CSLC(A,X,R,M)
    D) 10 I = 1,M
        WRITE(.......) X(I)
    10 CONTINUE
        STOP
        END
```

The solution of the system is

$$
X=((1.0,0.0),(1.0,0.0),(1.0,0.0),(1.0,0.0))
$$

### 3.4. Algorith

The algurithm for solving a system of linear algebraic equations

$$
\begin{equation*}
A x=\mathbf{b} \tag{1}
\end{equation*}
$$

with a C-matrix $A$ of order $M$ proceeds from a similarity transformation of $A$ to a diagonal matrix

$$
D=Q^{\star} A Q \text {, }
$$

* 

where $Q$ is unitary. (The symbol * denotes conjugate transpose.) The elements of $Q$ are inverse discrete Fourier transformations defined as

$$
q_{i j}=E^{(1 .-1) \cdot(j-1)} / \sqrt{M},
$$

where $E=\exp (2 \pi \sqrt{-1} / M)$. The solucion $x$ of the system (1) is then determined as

$$
\begin{equation*}
\mathrm{x}=\mathrm{QD}{ }^{-1} \mathrm{Q}^{\star} \mathrm{b} \tag{2}
\end{equation*}
$$

The diagonal elements $d_{i i}$ of $D$ can be calculated as simply

$$
d_{i 1}=\sqrt{M} \sum_{j=1}^{M} q_{1 j} a_{j-1}, \quad i=1,2, \ldots, M
$$

In other words, if $d$ is a column vector composed of the diagonal elemente $d_{11}, d_{22}, \ldots, d_{M M}$ of $D$, and $a$ is a column vector composed of the elements $a_{0}, a_{1}, \ldots, a_{M-1}$ of the first row of $A$, then these vectors are related by $\mathrm{d}=\sqrt{\mathrm{M}} \mathrm{Qa}$.

### 3.5. Prograning details

In the implementation of subroutine CSLC, ins'ead of $Q$ the matrix $\bar{Q}=\sqrt{M} Q$ is used, and formula (2) of subsection 3.4 becomes

$$
\mathrm{x}=\overline{\mathrm{Q}}^{-1} \overline{\mathrm{Q}}^{\star} \mathrm{b} / \mathrm{M}
$$

The vector $d$ composed of the diagonal elements $d_{i i}$ of $D$ is then calculated more simply as

$$
d=\bar{Q} \mathbf{a}
$$

Towards timing estimation, note that the algorithm for solving linear systems with $C$-matrices requires approximately $3 M^{2}$ multiplications.

### 3.6. Additional information

The calling sequence of subroutine CSLC for the solution of linear systems with C-matrices limits the right hand side to a single column vector. There may be situations where the solutions of two or more such systems with the same coefficient matrix are desired. In these situations, modifications of the subroutine that would permit all solutions to be obtained in a single step could markedly improve efficiency. Unlike TSLS and TSLSI discussed in subsection 2.6 , CSLC admits no simple modification for this purpose; however, subroutine SALWC could be used instead.

Subroutine SALWC is discussed in subsection 3.5 of Chapter ? -- it is called as a service subroutine in the solution of second- and third-level matrix systems. SALWC is similar to CSLC; its different organization, however, enables it to be separately useful, although somewhat awkward, for the solution of C-matrix systems with multiple right hand sides. Its use requires three calls with some arithmetic in-between, the presentation of the transpose of the right hand side matrix, and additional work space; also, unlike CSLC, it overwrites the coefficient array.

The folilowing program segment illustrates the use of SALWC for C-matrix systems of drder $M$ with MROW right hand sides (refer to subsection 3.5 of Chapter 2 for a description of the SALWC calling sequence).

```
COMPLEX \(A(M), X(L D X, M), R 1(M), R 2(M)\)
:
    RM 7 FLOAT (M)
    CALL \(\operatorname{SALWC}(A, R 1, R 2,1, M, 1,-1)\)
    CALL SALWC(X,R1,R2,MROW,M,LDX,1)
    DO \(10 \mathrm{~J}=1, \mathrm{M}\)
        DO \(5 \mathrm{I}=1\), MROW
            \(X(I, J)=X(I, J) / A(J) / R M\)
5 CONTINUE
10 CONTINUE
CALL SALWC(X,R1,R2,MROW,M,LDX,-1)
```

The dominant term in the multiplication count for the above segment is $M^{2} \cdot(2 \cdot M R O W+1)$, while for MROW calls of CSLC it is $3 M^{2} \cdot M R O W$. Comparing these quantities leads to the expectation that when MROW is 1 the two algorithms should be about equally fast, and as MROW increases a savings of up to $1 / 3$ should be possible with the above segment. For double precision, substitute SALWZ.

## 4. Solution of Linear Equations with G-Matrices

4.1. Purpose

Capability to solve linear algebraic equations with G-matrices is required for processing second- and third-level Toeplitz- and circulari-type matrices described in Chapters 2 and 3. The availability of the LINPACK
package makes it unnecessary to duplicate effort to provide this capability; the TOEPLITZ package simply invokes that subset of LINPACK which treats general square matrices. Usage will be briefly described for the single precision real version; double precision, complex, and double precision complex versions are aln available. p-ferral to the LINPACK Users' Guide [15] is recommended for fuller discussion than will be given here, including algorithm descriptions and programming details.

### 4.2. Usage

Single precision real G-matrices. SGEFA and SGESL together solve a linear system $A x=b$ with a real general matrix $A$; SGEFA computes the LU factorization of $A$ and SGESL uses the factorization to solve the linear system.

The calling sequence for SGEFA is
CALL SGEFA (A,LDA, M, PVT, INFO) •
On entry,

A is a doubly subscripted $M$ by $M$ array which contains the G-matrix.

LDA is the leading dimension of the array $A$.
$M$ is the order of $A$ and the number of elements in PVT.

On return,

A contains information from the LU factorization.

PVT is a singly subscripted array of $M$ elements which contains information to be transmitted to SGESL about the pivoting strategy used in the factorization. Note: In the LINPACK package PVT is specified as an integer array. For use in the TOEPLITZ package, PVT has the variable type of $A$; this simplifies the partition of the work space.

INrO is an integer which if nonzero warns of singularity of A. Note: Nonsingularity of $A$ and indeed all its principal minors is fundamental for use of the TOEPLITZ package; no interrogation of INFO is made anywhere.

The calling sequence for SGESL is
CALL SGESL(A,LDA, M, PVT, X,JOB) -
On entry,

A is a doubly subscripted $M$ by $M$ array which contains the information from the factorization stored by SGEFA.

LDA is the leading dimension of the array $A$.
$M$ is the order of $A$ and the number of elements in $X$ and PVT.

PVT is a singly subscripted array of $M$ elements which contains the pivot information stored by SGEFA.
$X$ is a singly subscripted array of $M$ elements which contains the right hand side of the system.

JOB is an integer which specifies the system to be solved. If JOB is zero, the system $A x=b$ is solved. If $J O B$ is nonzero, the system $A^{T} x=b$ is solved. Note: In its use with the TOEPLITZ package, JOB is always zero.

On return,
$x$ contains the solution of the system.

Double precision real G-matrices. The calling sequences of the double precision real G-matrix subroutines DGEFA and DGESL are the same as those of SGEFA and SGESL with A, X, and PVT DOUBLE PRECISION variables.

Single precision complex G-matrices. The calling sequences of the single precision complex G-matrix subroutines CGEFA and CGESL are the same as those of SGEFA and SGESL with A, X, and PVT COMPLEX variables.

Double precision complex G-matrices. In those computing systems where they are available, the calling sequences of the double precision complex G-matrix subroutines ZGEFA and ZGESL are the same as those of SGEFA and SGESL with $A, X$, and PVT DOUBLE COMPLEX variables.

## 5. Orthogonal Factorization of Colum-Circulant Matrices

### 5.1. Purpose

Given an $M$ by $L$ column-circulant matrix $A$, the TOEPLI'TZ subroutines in this section determine an $M$ by $L$ matrix $Q$ with orthonormal columns and an upper triangular matrix $S$ of order $L$ such that $A S=Q$. The $A S=Q$ factorization can be transformed to the more familiar $A=Q R$ factorization by inverting $S$, i.e., $R=S^{-1}$. Usage will be de. ribed here for the single precision real version. Double precision, complex, and double precision complex versions are also available.
5.2. Osage

Single precision real column-circulant matrices. CQRS performs the orthogonal factorization $A S=Q$ of a real column-circulant matrix $A$. The calling sequence is

CALL CQRS (A,Q,S,M,L,LDQ,LDS) •
On entry,

A is a singly subscripted array of $M$ elements which contains the first column of the column-circulant matrix. A is unaltered by CQRS.
$M$ is the number of rows of the matrices $A$ and $Q$. $M$ must be at ${ }^{\top}$ east equal to $L$.
$L$ is the number of columns of the matrices $A$ and $Q$ and the order of the upper triangular matrix $S$.

LDQ is the leading dimension of the array $Q$.

LDS is the leading dimension of the array $S$.

On return,
$Q \quad i s$ a doubly subscripted $M$ by $L$ array which contains the factor with orthonormal columns.
is a doubly subscripted $L$ by $L$ array which contains the upper triangular factor. Elements below the main diagonal of $S$ are not accessed.

Double precision real column-circulant matrices. The calling sequence of the double precision real column-circulant orthogonal factorization subroutine CQRD is the same as that of CQRS with $A, Q$, and $S$ DOUBLE PRECISION variablas.

Single precision complex column-circulant matrices. The calling sequence of the single precision complex column-circulant orthogonal factorization subroutine CQRC is the same as that of CQRS with $A, Q$, and $S$ COMPLEX variables.

Double prectsion complex column-circulant matrices. In those computing systems where it is available, the calling sequence of the double precision complex column-circulant orthogonal factorization subroutine CQRZ is the same as that of CQRS with $A, Q$, and $S$ COUBLE COMPLEX variables.

### 5.3. Example

The following program segment illustrates the use of the single precision subroutine for orthogonal factorization of real column-circulant matrices; factors $Q$ and $S$ are returned satisfying $A S=Q$. Examples of the use of CQRD, CQRC, and CQRZ could be obtained by changing the subroutine name and type declaration. The matrix is 4 by 3 with coefficients as follows.

$$
A=\left|\begin{array}{lll}
1 & 4 & 3 \\
2 & 1 & 4 \\
3 & 2 & 1 \\
4 & 3 & 2
\end{array}\right|
$$

REAL $A(4), Q(4,3), S(3,3)$
INTEGER M,L,LDQ,LDS,I,J
DATA $A(1) / 1.0 /, A(2) / 2.0 /, A(3) / 3.0 /, A(4) / 4.0 /$
$M=4$
L = 3
$L D Q=4$
LDS $=3$
CALL CQRS(A,Q,S,M,L,LDQ,LDS)

```
    DO 10 I = l, M
    WRITE(...,...) (Q(I,J),J=1,L)
10 CONTINUE
    DO 20 I = 1, L
        WRITE(...,...) (S(I,J),J=I,L)
20 continue
    STOP
    END
```

The factors $Q$ and $S$ are
$Q=\left|\begin{array}{ccc}1 / \sqrt{30} & 16 / \sqrt{270} & 10 / \sqrt{7344} \\ 2 / \sqrt{30} & -3 / \sqrt{270} & 78 / \sqrt{7344} \\ 3 / \sqrt{30} & -2 / \sqrt{270} & -26 / \sqrt{7344} \\ 4 / \sqrt{30} & -1 / \sqrt{270} & -22 / \sqrt{7344}\end{array}\right| \quad S=\left\{\left.\begin{array}{ccc}1 / \sqrt{30} & -4 / \sqrt{270} & -7 / \sqrt{7344} \\ 0 & 5 / \sqrt{270} & -16 / \sqrt{7344} \\ 0 & 0 & 27 / \sqrt{7344}\end{array} \right\rvert\,\right.$

### 5.4. Algorithn

The algorithm description can be found in [12]. Note that usage of this algorithm for orthogonal factorization of column-circulant matrices requires that the matrix have full rank $L$.

### 5.5. Programing details

The algorithm for the orthogonal factorization of an $M$ by $L$ columncirculant matrix requires approximately $6 \mathrm{ML}+\mathrm{L}^{2}$ multiplications.

## CHAPTER 2: TORPLITZ- AND CIRCIIANT-TYPE MATRICES OF THE SBCOND LEVEL

## 1. Structure and Representation

1.1. Overview

A matrix

$$
A=\left|\begin{array}{llllll}
A_{11} & A_{12} & A_{13} & \cdots & \cdot & A_{1 L}  \tag{1}\\
A_{21} & A_{22} & A_{23} & \cdots & \cdots & A_{2 L} \\
A_{31} & A_{32} & A_{33} & \cdots & A_{3 L} & A_{3 L} \\
\cdots & \cdot & \cdots & \cdots & \cdot & \cdot \\
A_{L 1} & A_{L 2} & A_{L 3} & \cdots & \cdot & A_{L L}
\end{array}\right|
$$

with $L$ elements in a row (or column) where the elements $A_{1 j}$ are blocks of order $M$ is called a two-level matrix. L is called the first-level order and $M$ becomes the second-level order of the matrix $A$. The order $N$ of $A$ is then the product of the orders of its levels: $N=L \star M$.

We will call the two-level matrix (1) an XY-type if A considered as a block matrix is an $X$-type and each of its blocks $A_{i j}$ is a $Y$-type. As $X$ - and Y-types in the TOEPLITZ package we consider $T$-, $C$-, and Grmairices defined in section 1 of Chapter 1 . Examples of two-level matrices can be found below and in subsections $2.3,3.3,4.3$, and 5.3 of this chapter.

By permuting corresponding rows and columns, we can transform any XY-type to YX-type (see Tyrtyshnikov [25]). For example, the TC-matrix

$$
\left(\left.\begin{array}{lll|lll}
a & b & c & d & e & f \\
c & a & b & f & d & e \\
b & c & a & e & f & d \\
\hline g & h & 1 & a & b & c \\
1 & g & h & c & a & b \\
h & 1 & g & b & c & a
\end{array} \right\rvert\,\right.
$$

with $L=2, M=3$ can be permuted to the CT-matrix
$\left(\left.\begin{array}{ll|ll|ll}a & 1 & b & g & c & h \\ e & a & f & b & d & c \\ \hline c & h & a & 1 & b & g \\ d & c & e & a & f & b \\ \hline b & g & c & h & a & i \\ f & b & d & c & e & a\end{array} \right\rvert\,:\right.$
with $L=3, M=2$ by interchanging $r o w$ and column pairs (1,6) and (3,4). This circumstance allows us to limit consideration to one of each XY- YX-type pair.

The scheme for compact representation of two-level matrices is the following. Let $A$ be of $X Y$-type with first-level orler $L$ and second-level order M. Furthermore, let $\tilde{L}$ be the number of elements required in the compact representation of $\ddot{A}$ and $\tilde{M}$ be the number of elements required in the compact representstion of $Y$. Recall that for $T-C$, and G-matrices of order $M$ as described in section 1 of Chapter 1 the values of $\tilde{M}$ are, respectively, $2 \star M-1$, $M$, and $M \star * 2$. In the TOEPLITZ package such a two-level matrix is represented by a doubly subscripted $\tilde{M}$ by $\tilde{L}$ array. The blocks in the array are indexed by the second subscript and ordered in accordance with the $x$-type compact representation. In turn, the elements in a block are indexed by the first subscript and ordered in accordance with the block's Y-type compact representation.

### 1.2. TG-matrices

A matrix
is called a TG-matrix if $A_{1}$ and $A_{-1}, i=0,1,2, \ldots, L-1$, are $G$-matrices of order M (see subsection 1.3 of Chapter 1 ).

In the TOEPLITZ package this TG-matrix is represented by a doubly subscripted $M * * 2$ by $2 * L-1$ array in which the blocks are ordered in the following way:

$$
A_{0}, A_{1}, A_{2}, \ldots, A_{L-1}, A_{-1}, A_{-2}, \ldots, A_{-L+1}
$$

### 1.3. CT-antrices

A complex matrix

$$
A=\left|\begin{array}{lllll}
A_{0} & A_{1} & A_{2} & \cdots & A_{L-1}  \tag{2}\\
A_{L-1} & A_{0} & A_{1} & \cdots & A_{L-2} \\
A_{L-2} & A_{L-1} & A_{0} & \cdots & A_{L-3} \\
\cdot & \cdot & \cdots & \cdots & \cdot \\
A_{1} & A_{2} & A_{3} & \cdots & A_{0}
\end{array}\right| .
$$

is called a CT-matrix if $A_{i}, i=0,1,2, \ldots, L-1$, are $T$-matrices of order $M$ (see subsection 1.1 of Chapter 1).

In the TOEPLITZ package this CT-matrix is represented by a doubly subscriptei $2 \star M-1$ by $L$ array in which the blocks are ordered in the following way:

$$
A_{0}, A_{1}, A_{2}, \ldots, A_{L-1}
$$

### 1.4. CC-matrices

A matrix of form (2) is called a CC-matrix if $A_{1}, 1=0,1,2, \ldots, L-1$, are $C$ matrices of order $M$ (see subsection 1.2 of Chapter 1 ).

In the TOEPLITZ package this CC-matrix is represented by a doubly subscripted $M$ by $L$ array in which the blocks are ordered in the following way:

$$
A_{0}, A_{1}, A_{2}, \ldots, A_{L-1}
$$

### 1.5. CG-matrices

A matrix of form (2) is called a CG-matrix if $A_{1}, i=0,1,2, \ldots, L-1$, are G-matrices of order $M$ (see subsection 1.3 of Chapter 1 ).

In the TOEPLITZ package this CG-matrix is represented by a doubly subscripted $M * * 2$ by $L$ array in which the blocks are ordered in the following way:

$$
A_{0}, A_{1}, A_{2}, \ldots, A_{L-1} .
$$

### 1.6. Other types of two-level matrices

GT-, TC-, and GC-matrices, defined in analogous ways, can be permuted, respectively, to TG-, CT-, and CG-matrices (see example in subsection l.1). Therefore, the TOEPLITZ package does not include subroutines for solving linear systems with two-level matrices of these types. At the present time no algorithm is known that capitalizes effectively on the structure of TT-matrices, so TT-matrices should be treated as TG-matrices.

## 2. Solution of Linear Equations with TG-matrices

### 2.1. Purpose

The TOEPLITZ subroutines in this section are designed to solve linear algebraic equations with TG-matrices, that is, block-Toeplitz matrices whose blocks are G-matrices. Usage will be described for the single precision real version. Double precision, complex, and double precision complex versions are also avallable. Indeed, the complex version is called in solving three-level CTG-matrix systems (see subsection 2.5 of Chapter 3 ).

### 2.2. Usage

Single precision real TG-matrices. TGSLS solves a linear system with a real block-Toeplitz matrix whose blocks are G-matrices. The calling sequence 18

CALL TCSLS (A, X,R,M,L,LDA) .

On entry,

A is a doubly subscripted $M * * 2$ by $2 * L-1$ array which contains the TG-matrix in the form described in subsection l.2. A is unaltered by TGSLS.
$X$ is a singly subscripted array of $M * L$ elements which contains the right hand side of the system.
$R$ is a singly subscripted array of $2 \star M * * 2 * L+3 * M * * 2+M$ elements used for work space.
$M$ is the order of each G-matrix block of $A$.

L is the number of blocks in each row or column of $A$.

LDA is the leading dimension of the array $A$.

On return,
$X$ contains the solution of the system.

Double precision real TG-matrices. The calling sequence of the double precision real TG-matrix subroutine TGSLD is the same as that of TGSLS with $A$, $X$, and $R$ DOUBLE PRECISION variables.

Single precision complex TG-matrices. The calling sequence of the single precision complex TG-matrix subroutine TGSLC is the same as that of TGSLS with $A, X$, and $R$ COMPLEX variables.

Double precision complex TG-matrices. In wose computing systems where it is available, the calling sequence of the double precision complex TG-matrix subroutine TGSLZ is the same as that of TGSLS with $A, X$, and $R$ DOUBLE COMPLEX variables.

### 2.3. Buaple

The following program segment illustrates the use of the single precision subroutine TGSLS for real TG-matrices. Examples of the use of TGSLD, TGSLC,
and TGSLZ could be obtained by changing the subroutine name and type declaration. The system is of order 4 with coefficients as follows.


REAL $A(4,3), X(4), R(30)$
INTEGER M,L,LDA,I,J
DATA $\mathrm{A}(1,1) / 1.0 /, \mathrm{A}(2,1) / 2.0 / \mathrm{A}(3,1) / 3.0 /, \mathrm{A}(4,1) / 4.0 /$,

* $\mathrm{A}(1,2) / 5.0 /, \mathrm{A}(2,2) / 6.0 /, \mathrm{A}(3,2) / 7.0 /, \mathrm{A}(4,2) / 8.0 /$,
* $\quad \mathrm{A}(1,3) / 9.0 /, \mathrm{A}(2,3) / 10.0 /, \mathrm{A}(3,3) / 11.0 /, \mathrm{A}(4,3) / 12.0 /$

DATA $\mathrm{X}(1) / 16.0 /, \mathrm{X}(2) / 20.0 /, \mathrm{X}(3) / 24.0 /, \mathrm{X}(4) / 28.0 /$
$M=2$
$\mathrm{L}=2$
LDA $=4$
CALL TGSLS(A,X,R,M,L,LDA)
$\mathrm{J}=\mathrm{M} * \mathrm{~L}$
DO $10 \mathrm{I}=1$, J WRITE(...,....) X(I)
10 CONTINUE
STOP
END
The solutinn of the system is

$$
x=(1.0,1.0,1.0,1.0) .
$$

### 2.4. Algorithe

The algorithm for solving a linear system

$$
\begin{equation*}
A x=b \tag{1}
\end{equation*}
$$

with the TG-matrix

$$
A=\left(\begin{array}{llllll}
A_{0} & A_{1} & A_{2} & \cdot & \cdot & A_{L-1} \\
A_{-1} & A_{0} & A_{1} & \cdot & \cdot & A_{L-2} \\
A_{-2} & A_{-1} & A_{0} & \cdot & \cdot & A_{1-3} \\
\cdot & \cdot & \cdot & \cdot & \cdot & \cdot
\end{array}\right),
$$

where $A_{i}$ and $A_{-1}, i=0,1,2, \ldots, L-1$, are G-matrices of order $M$, is the block analogue of the algorithm for solving linear systems with T-matrices (see subsection 2.4 of Chapter 1).

Let us introduce the following notation:

$$
C_{k}=\left(\left.\begin{array}{ccccc}
A_{0} & A_{1} & \cdot & \cdot & A_{k} \\
A_{-1} & A_{0} & \cdot & \cdot & \cdot \\
\cdot A_{k-1} \\
A_{-k} & A_{-k+1} & \cdot & \cdot & A_{0}
\end{array} \right\rvert\,, \quad y_{k}=\left(\begin{array}{c}
y_{0, k} \\
y_{1, k} \\
\vdots \\
y_{k, k}
\end{array}\right), \quad d_{k}=\left(\begin{array}{c}
b_{0} \\
b_{1} \\
\vdots \\
b_{k * M+M-1}
\end{array}\right)\right.
$$

where $y_{1, k}, 1=0,1, \ldots, k$, are vectors of $M$ elements. The algorithm consists of step-by-step recurrent solution of systems

$$
\begin{equation*}
c_{k} y_{k}=d_{k} \tag{2}
\end{equation*}
$$

for $k=0,1,2, \ldots, 1-1$. The final result of the recurrent process is the solution of the given system (1):

$$
\mathrm{x}=\mathrm{y}_{\mathrm{L}-1}
$$

At step $1, y_{0}=A_{0}^{-1} d_{0}$. At step $k+1$, the vector $y_{k}$ is calculated from $y_{k-1}$ as follows. Let us consider the vector $y_{k}$ to be the sum of two vectors, one of which, augmented by a zero vector of $M$ elements, was determined at the $k-t h$ step:

$$
\left(\begin{array}{c}
y_{0, k}  \tag{3}\\
y_{1, k} \\
\vdots \\
y_{k-1, k} \\
y_{k, k}
\end{array}\right)=\left(\begin{array}{c}
y_{0, k-1} \\
y_{1, k-1} \\
\vdots \\
y_{k \cdots-1, k-1} \\
0
\end{array}\right)+\left(\begin{array}{c}
z_{0, k} \\
z_{1, k} \\
\vdots \\
z_{k-1, k} \\
z_{k, k}
\end{array}\right)
$$

Substituting this sum into equation (2) and taking into account that the vector $y_{k-1}$ satisfies the equation

$$
c_{k-1} y_{k-1}=d_{k-1}
$$

we see that the unknown vector $z_{k}$ from (3) consisting of component vectors $z_{0, k}, z_{1, k}, \ldots, z_{k, k}$ each $:=M$ elements is the solution of the system

$$
C_{k} z_{k}=f_{k}
$$

where

Thus, the vector $z_{k}$ is a linear combination of the last $M$ columns of the matrix $C_{k}^{-1}$, and the elements of the vector $f_{k, k}$ are the coefficients of that linear combination. Hence, for recurrent calculation of the vectors $y_{k}$ it is sufficient to evaluate recurrently the last block column of the matrix $C_{k}^{-1}$, or as done here for further economy an appropriately chosen block multiple of this block column. It is here that advantage is taken of the block-Toeplitz structure of $A$.

Let us denote by $G_{k}$ and $H_{k}$ the first and last block columns, respectively scaled by M-order matrices $P_{k}$ and $Q_{k}$, of the matrix $C_{k}^{-1}$ :

$$
G_{k}=\left\{\begin{array}{c}
G_{0, k} \\
G_{1, k} \\
\vdots \\
G_{k-1, k} \\
G_{k, k}
\end{array}\right\} \quad, \left.\quad \begin{gathered}
H_{0, k} \\
H_{1, k} \\
\vdots \\
H_{k}= \\
H_{k-1, k} \\
H_{k, k}
\end{gathered} \right\rvert\,
$$

and

$$
C_{k} G_{k}=\left(\begin{array}{c}
P_{k} \\
0 \\
\vdots \\
0 \\
0
\end{array} \left\lvert\,, \quad C_{k} H_{k}=\left(\left.\begin{array}{c}
0 \\
0 \\
\vdots \\
0 \\
\\
Q_{k}
\end{array} \right\rvert\,\right.\right.\right.
$$

It is clear that when $k=0$ the unscaled block columns coincide and contain the single block $A_{0}^{-1}$; we shoose $P_{0}=Q_{0}=A_{0}$ so that $G_{0}=H_{0}=I$. we will determine $G_{k}, H_{k}, P_{k}$, and $G_{k}$ from $G_{k-1}, H_{k-1}, P_{k-1}$, and $Q_{k-1}$ using the following two sums:

$$
\begin{aligned}
& G_{k}=\left|\begin{array}{c}
G_{0, k-1} \\
G_{1, k-1} \\
\vdots \\
G_{k-1, k-1} \\
0
\end{array}\right|+\left(\left.\begin{array}{c}
0 \\
H_{G, k-1} \\
\vdots \\
H_{k-2, k-1} \\
H_{k-1, k-1}
\end{array} \right\rvert\, V,\right. \\
& H_{k}=\left|\begin{array}{c}
G_{0, k-1} \\
G_{1, k-1} \\
\vdots \\
G_{k-1, k-1} \\
0
\end{array}\right| R+\left|\begin{array}{c}
0 \\
H_{0, k-1} \\
\vdots \\
H_{k-2, k-1} \\
H_{k-1, k-1}
\end{array}\right|,
\end{aligned}
$$

where $V$ and $R$ are unknown $M$ by $M$ matrices which we are going to derive.
Since $G_{k}$ and $H_{k}$ are block columns of the matrix $C_{k}^{-1}$ scaled by $P_{k}$ and $Q_{k}$, respectively, then

$$
C_{k} G_{k}=C_{k}\left(\begin{array}{c}
G_{0, k-1} \\
G_{1, k-1} \\
\vdots \\
G_{k-1, k-1} \\
0
\end{array}\right)+C_{k}\left(\begin{array}{c}
0 \\
H_{0, k-1} \\
\vdots \\
H_{k-2, k-1} \\
H_{k-1, k-1}
\end{array}\right) v=\left(\begin{array}{c}
P_{k} \\
0 \\
\vdots \\
0 \\
0
\end{array}\right),
$$

$$
C_{k} H_{k}=C_{k}\left(\begin{array}{c}
G_{0, k-1} \\
G_{1, k-1} \\
\vdots \\
G_{k-1, k-1} \\
0
\end{array}\left|R+C_{k}\right| \begin{array}{c}
0 \\
H_{0, k-1} \\
\vdots \\
H_{k-2, k-1} \\
{ }_{H},{ }_{k-1, k-1}
\end{array}\right)=\left(\left.\begin{array}{c}
0 \\
0 \\
\vdots \\
0 \\
Q_{k}
\end{array} \right\rvert\,\right.
$$

These relationships reduce to the following equations for determining the unknown matrices:

$$
\left\{\begin{array} { l } 
{ P _ { k - 1 } + F _ { 2 } V = P _ { k } }  \tag{4}\\
{ F _ { 1 } + Q _ { k - 1 } V = 0 }
\end{array} \quad \left\{\begin{array}{l}
P_{k-1} R+F_{2}=0 \\
F_{1} R+Q_{k-1}=Q_{k}
\end{array}\right.\right.
$$

where

$$
F_{1}=\sum_{\ell=1}^{k} A_{-\ell} G_{k-\ell, k-1}, \quad F_{2}=\sum_{\ell=1}^{k} A_{\ell}{ }_{\ell-1, k-1} .
$$

Solving systems (4) we find

$$
\begin{aligned}
& V=-\left(Q_{k-1}\right)^{-1} F_{1}, \quad R=-\left(P_{k-1}\right)^{-1} F_{2}, \\
& P_{k}=P_{k-1}-F_{2}\left(Q_{k-1}\right)^{-1} F_{1}, \\
& Q_{k}=Q_{k-1}-F_{1}\left(P_{k-1}\right)^{-1} F_{2} .
\end{aligned}
$$

Note that this algorithm for solving linear systems with TG-matrices requires that $C_{k}$ be non-singular for all $k$.

### 2.5. Prograning details - subroutine TGSLSI

Subroutine TGSLS merely acts as an interface to subroutine TGSLSI, in the manner of TSLS and TSLSI for T-matrices as explained in subsection 2.5 of Chapter 1.

The calling sequence of subroutine TGSLSI is
CALL TGSLS1(A1,A2,B,X,Cl,C2,R1,R2,R3,R5,R6,R,M,L,LDA) •

On entry,

Al is a doubly subscripted $M * * 2$ by $L$ array which contains the first row of blocks of the TG-matrix. Al is unaltered by TGSLSl.

A2 is a doubly subscripted $M * * 2$ by $L-1$ array which contains the first column of blocks of the TG-matrix beginning with the second block. A2 is unaltered by TGSLS1.
$B \quad$ is a singly subscripted array of $M * L$ elements which contains the right hand side of the system. $B$ is unaltered by TGSLS1.
$C 1, C 2$ are triply subscripted arrays with dimension ( $M, M, L-1$ ) used for work space.
$R 1, R 2, R 3, R 5, R 6$ are doubly subscripted arrays with dimension ( $M, M$ ) used for work space.
$R \quad$ is a singly subscripted array of $M$ elements used for work space.
$M$ is the order of each G-matrix block of the TG-matrix.

L is the number of blocks in each row or column of the TG-matrix.

LDA is the leading dimension of the arrays Al and A 2 .

On return,
$X \quad$ is a singly subscripted array of $M * L$ elements which contains the solution of the system. $X$ may coincide with $B$.

For solving G-matrix systems in accordance with the algorithm described in subsection 2.4, TGSLS calls the LINPACK subroutines SGEFA and SGESL (see section 4 of Chapter 1 ).

Vector operations are facilitated by calls to the LINPACK BLA subroutine SAXPY. This subroutine is coind fficiently but there is a cust associated
with communcation to it; this cost can become relatively large when computation within SAXPY itself is small and the further computations of TGSLSl are highly optimized by the compller. Therefore, when the number of vector components ( $M$ for two-level TG-matrices) is small and the compiler is capable of a high level of optimization, it may be more efficient to perform the vector computations in-line instead of repeatedly calling SAXPY. (It is of Interest to note that in TGSLCl, overhead associated with the use of the corresponding LINPACK BLA subroutine CAXFy is much less significant in the presence of the slower complex arithmetí.) To facilitate possible change to in-line computation, directions are provided through code comments in subroutine TGSLSl (and also TGSLD1, TGSLCl, and TGSLZ1).

For solving systems with double precision, complex, and double precision complex TG-matrices, versions corresponding to TGSLSI are available with names TGSLDl, TGSLC1, and TGSLZl, respectively. These in turn call the corresponding versions of the LINPACK subroutines.

The algorithm implemented in subroutine TGSLS requires approximately $2 M^{3} L^{2}$ multiplications.

## 3. Solution of Linear Equations with CT-matrices

### 3.1. Purpose

The TOEPLITZ subroutine in this section is designed to solve linear algebiaic equations with CT-matrices, that is, complex block-circulant matrices whose biocks are T-matrices. $\quad$ double precision version of the subroutine is also avallable.

### 3.2. Usage

Single precision CT-matrices. CTSLC solves a linear system with a complex block-circulant matrix whose blocks are T-matrices. The calling sequence is

CALL CTSLC(A,X,R,M,L,LDA) •

On entry,

A is a doubly subscripted $2 \star M-1$ by $L$ array which contains the CT-matrix in the form described in subsection 1.3. A is destroyed by CTSLC.
$X$ is a singly subscripted array of $M * L$ elements which contains the right hand side of the syst.m.
$R$ is a singly subscripted array of $\max (2 * M-2,2 * L)$ elements used for work space.
$M$ is the order of each T-matrix block of $A$.

L is the number of blocks in each row or column of $A$.

LDA is the leading dimension of the array $A$.

On return,
$X$ contains the solution of the system.

Double precision CT-matrices. In those computing systems where it is available, the calling sequence of the double precision CT-matrix subroutine CTSLZ is the same as that of CTSLC with $A, X$, and $R$ DOURLE COMPLEX variables.

### 3.3. Bratple

The following program segment illustrates the use of the single precision subroutine CTSLC for CT-matrices. An example of the use of CTSLZ could be obtained by changing the subroutine name and type declaration. The system is of order 4 with coefficients as follows.
$A=\left|\begin{array}{ll|ll}1+1 & 2+21 & 2+21 & 3+31 \\ 3+31 & 1+1 & 4+41 & 2+21 \\ \hline 2+21 & 3+31 & 1+1 & 2+21 \\ 4+41 & 2+21 & 3+31 & 1+1\end{array}\right|$

$$
x=\left(\left.\begin{array}{c}
8+81 \\
10+101 \\
8+81 \\
10+101
\end{array} \right\rvert\,\right.
$$

```
    COMPLEX A(3,2),X(4),R(4)
    INTEGER M,L,LDA,I,J
    DATA A(1,1)/(1.0,1.0)/,A(2,1)/(2.0,2.0)/,A(3,i)/(3.0.3.0)/,
* A(1,2)/(2.0,2.0)/,A(2,2)/(3.0,3.0)/,A(3,2)/(4.0,4.0)/
    DATA X(1)/(8.0,8.0)/,X(2)/(10.0,10.0)/,X(3)/(8.0,8.0)/,
* X(4)/(10.0,10.0)/
    M=2
    L = 2
    LDA = 3
    CALL CTSLC(A,X,R,M,L,LDA)
    J = M*L
    DO 10 I = 1, J
        WRITE(....,...) X(I)
10 CONTINUE
    STOP
END
```

The solution of the system is

$$
X=((1.0,0.0),(1.0,0.0),(1.0,0.0),(1.0,0.0))
$$

### 3.4. Algorithm

The algorithm for solving a linear system

$$
\begin{equation*}
A x=b \tag{1}
\end{equation*}
$$

with the $C T-m a t r i x$

$$
A=\left|\begin{array}{cccccc}
A_{0} & A_{1} & A_{2} & \cdot & \cdot & A_{L-1} \\
A_{L-1} & A_{0} & A_{1} & \cdot & \cdot & A_{L-2} \\
A_{L-2} & A_{L-1} & A_{0} & \cdot & \cdot & A_{L-3} \\
\cdot & \cdot & \cdot & \cdot & \cdot & \cdot
\end{array}\right| \cdot \cdot\left|\begin{array}{lllll} 
\\
A_{1} & A_{2} & A_{3} & \cdot & A_{0}
\end{array}\right|
$$

where $A_{1}, i=0,1,2, \ldots, L-1$, are $T$-matrices of order $M$, proceeds from a similarity transformation of $A$ to a block-diagonal matrix

$$
D=Q^{\star} A Q
$$

in which each diagonal block is a T-matrix. (The symbol * denotes confugate transpose.) $Q$ is a two-level matrix with first-level order $L$ and second-level order M whose blocks are scalar matrices; the matrix of the scalars themselves is unitary.

Block $Q_{i j}$ of $Q$ is defined as

$$
Q_{i j}=E^{(1-1) \cdot(j-1)_{\star I} / \sqrt{L}},
$$

where $E=\exp (2 \pi \sqrt{-1} / L)$ and $I$ is the identity matrix of order $M$. However, as for C-matrices (see section 3 of Chapter l), it is more efficient to use instead the matrix

$$
\bar{Q}=\sqrt{\mathrm{L}} \mathrm{Q} .
$$

Thus the solution $x$ of the system (1) can be found by the following steps:
a) Transform the matrix $A$ to the block-diagonal matrix

$$
D=\vec{Q}_{\hat{A}}^{\hat{A}} \bar{Q} / L
$$

b) Transform the right hand side

$$
y=\bar{Q}^{*} b
$$

c) Solve the system

$$
\mathrm{Dz}=\mathrm{y}
$$

d) Transform the vector $z$ back to

$$
x=\bar{Q} z / L
$$

Note that since A is a CT-matrix, its transformation to becomes simply

$$
D_{1 i}=\sum_{j=1}^{L} \bar{Q}_{1 j} A_{j-1},
$$

where $D_{1 i}$ is the i-th diagonal block of $D$ and $A_{j-1}$ i. the block with index j-l at the tof of A. Furthermore, since D 18 block-diagonal each block of which is a T-matrix, the s;istem (1) reduces to $L$ systems with T-matrices.

### 3.5. Prograning detalis - aubroutine SALNC

The implementation of subroutine CTSLC corresponds to the algorithm described in subsection 3.4. All needed operations with matrices $\bar{Q}$ and $\bar{Q}$ are implemented by the service subroutine SALWC. The structure of these matrices
and the compact form of input representation are such that, from the point of view of programming, these operations (or more properly $Q$ and $Q^{*}$ ) can be considered respectively as inverse and direct discrete Fourier transformations upon a set of row vectors in a certain rectangular matrix.

The calling sequence of subroutine SALWC is
CALL SALWC (A,R1,R2,M,L,LDA, YOB) •
On entry,

A is a doubly subscripted $M$ by $L$ array which contains the matrix upon whose rows the Fourler transformation will be performed.

R1,R2 are singly subscripted arrays of $L$ elements used for work space.
$M$ is the number of rows of $A$.

L is the number of columns of $A$.

LDA is the leading dimension of the array $A$.
$J O B$ indicates what is to be computed. If $J O B$ is 1 , the direct Fourier transformation will be performed and if $J O B$ is -1 , the inverse Fourier transformation will be performed.

On return,

A contains the transformed rows of the matrix.

For solving the $L$ systems with T-matrices, first-level subroutines TSLC and TSLCI are called. For solving systems with double precision CT-matrices (using CTSLZ), the double precision subroutine SALWZ is called, as well as TSLZ and TSLZ1.

The overall algorithm implemented in subroutine CTSLC requires approxiwately $4 M L^{2}+3 M^{2} L$ multiplications $-4 M^{2}$ in SALWC and $3 M^{2} L$ in TSLCI.

## 4. Solution of Linear Bquations with CC-matrices

### 4.1. Purpose

The TOEPLITZ subroutine in this section is designed to solve linear algebraic equations with CC-matrices, that is, complex block-circulant matrices whose blocks themselves are circulant matrices. A double precision version of the subroutine is also available.

### 4.2. Usage

Single precision CC-matrices. CCSLC solves a linear system with a complex block-circulant matrix whose blocks are C-matrices. The calling sequence is

CALL CCSLC(A,X,R,M,L,LDA) •
On ntry,

A is a doubly subscripted $M$ by $L$ array which contains the CC-matrix of the system in the form described in subsection l.4. A is destroyed by CCSLC.
$X$ is a singly subscripted array of $M * L$ elements which contains the right hand side of the system.
$R \quad$ is a singly subscripted array of $\max (M, 2 \star L)$ elements used for work space.
$M$ is the order of each $C$-matrix block of $A$.

L is the number of blocks in each row or column of $A$.

LDA is the leading dimension of the array $A$.

On return,
$\chi$
contains the solution of the system.

Double precision CC-matrices. In those computing systems where it is available, the calling sequence of the double precision CC-matrix subroutine $\operatorname{CCSLZ}$ is the same as that of CCSLC with $A, X$, and $R$ DOUBLE COMPLEX variables.

### 4.3. Example

The following program segment illustrates the use of the single precision subroutine CCSLC for CC-matrices. An example of the use of CCSLZ could be obtained by changing the subroutine name and type declaration. The system is of order 4 with coefficients as follows.

$$
\begin{aligned}
& A=\left(\begin{array}{ll|ll}
1+1 & 2+2 i \\
2+2 i & 1+1 & 2+2 i & 4+4 i \\
\hline 2+2 i & 4+4 i & 1+4 i & 2+2 i \\
4+4 i & 2+2 i & 2+2 i & 1+1
\end{array}\left|\quad X=\begin{array}{l}
9+21 \\
9+91 \\
9+9 i \\
9+9 i
\end{array}\right|\right. \\
& \text { COMPLEX A(2,2),X(4),R(4) } \\
& \text { INTEGER M,L,LDA,I,J } \\
& \text { DATA } A(1,1) /(1.0,1.0) /, A(2,1) /(2.0,2.0) / \text {, } \\
& \text { * } \quad \mathrm{A}(1, ?) /(2.0,2.0) /, \mathrm{A}(2,2) /(4.0,4.0) / \\
& \text { DATA X(1)/(9.0,9.0)/,X(2)/(9.0,9.0)/, } \\
& \text { * } \mathrm{X}(3) /(9.0,9.0) /, \mathrm{X}(4) /(9.0,9.0) / \\
& M=2 \\
& \mathrm{~L}=2 \\
& \text { LDA }=2 \\
& \text { CALL CCSLC (A,X,R,M,L,LDA) } \\
& J=M \star L \\
& \text { DO } 10 \mathrm{I}=1 \text {, J } \\
& \text { WRITE (........) X(I) } \\
& 10 \text { CONTINUE } \\
& \text { STOP } \\
& \text { END }
\end{aligned}
$$

The solution of the system is

$$
X=((1.0,0.0),(1.0,0.0),(1.0,0.0),(1.0,0.0))
$$

### 4.4 Algorith

The algorithm used in subroutine CCSLC is the same as that described in subsection 3.4 for CT-matrices except that the solution of C-matrix rather than T-matrix systems is involved.

### 4.5. Programing details

Programming details of subroutine CCSLC are as for CTSLC (see subsection 3.5) except that subroutine CSLC is called instead of subroutines TSLC and TSLCl; the number of multiplications is approximately $3 \mathrm{LL}^{2}+3 M^{2} \mathrm{~L}-3 \mathrm{LL}^{2}$ in SALWC and $3 M^{2} \mathrm{~L}$ in CSLC.

## 5. Solution of Linear Equations with CG-matrices

### 5.1. Purpose

The TOEPLITZ subroutine in this section is designed to solve linear algebraic equations with CG-matrices, that is, complex block-circulant matrices whose blocks are general matrices. A double precision version of the subroutine is also available.
5.2. Usage

Single precision CG-matrices. CGSLC solves a linear system with a complex block-circulant matrix whose blocks are G-matrices. The calling sequence is

CALL CGSLC(A,X,R,M,L,LDA) .
On entry,
A is a doubly subscripted $M \star{ }^{2} 2$ by $L$ array which contains the CG matrix of the system in the form described in subsection l.5. A is destroyed by CGSLC.
$X$ is a singly subscripted array of $\mathrm{M}_{\mathrm{L}} \mathrm{L}$ elements which contains the right hand side of the system.
$R \quad$ is a singly subscripted array of $\max (M, 2 \star L)$ elements used for work space.
$M$ is the order of each G-matrix block of $A$.

L is the number of blocks in each row or column of $A$.

LDA is the leading dimension of the array $A$.

On return,
$X$ contains the solution of the system.

Double precision CG-matrices. In those computing systems where it is availiable, the calling sequence of the double precision CG-matrix subroutine CGSLZ is the same as that of CGSLC with $A, X$, and $R$ DOUBLE COMFLEX variables.

### 5.3. Example

The following program segment illustrates the use of the singie precision subroutine CGSLC for CG-matrices. An example of the use of CGSLZ could be obtained by changing the subroutine name and type declaration. The system is of order 4 with coefficients as follows.

$$
\begin{aligned}
& A=\left|\begin{array}{cc|cc}
1+1 & 3+3 i & 5+5 i & 7+7 i \\
2+2 i & 4+4 i & 6+6 i & 8+81 \\
\hline 5+5 i & 7+7 i & 1+i & 3+3 i \\
6+6 i & 8+8 i & 2+2 i & 4+4 i
\end{array}\right| \\
& X=\left|\begin{array}{l}
16+16 i \\
20+20 i \\
16+16 i \\
20+20 i
\end{array}\right| \\
& \text { COMPLEX } A(4,2), X(4), R(4) \\
& \text { INTEGER M,L,LDA,I,J } \\
& \text { DATA A }(1,1) /(1.0,1.0) /, \mathrm{A}(2,1) /(2.0,2.0) /, \mathrm{A}(3,1) /(3.0,3.0) / \text {, } \\
& \star \quad \mathrm{A}(4,1) /(4.0,4.0) /, \mathrm{A}(1,2) /(5.0,5.0) /, \mathrm{A}(2,2) /(6.0,6.0) / \text {, } \\
& \text { * } \mathrm{A}(3,2) /(7.0,7.0) /, \mathrm{A}(4,2) /(8.0,8.0) / \\
& \text { DATA } \mathrm{X}(1) /(16.0,16.0) /, \mathrm{X}(2) /(20.0,20.0) /, \mathrm{X}(3) /(16.0,16.0) / \text {, } \\
& \text { * } X(4) /(20.0,20.0) / \\
& M=2 \\
& L=2 \\
& \text { LDA }=4 \\
& \text { CALL } \operatorname{CGSLC}(A, X, R, M, L, L D A) \\
& J=M * L \\
& \text { DO } 10 \mathrm{I}=1 \text {, } \mathrm{J} \\
& \text { WRITE (........) X(I) } \\
& 10 \text { CONTINUE } \\
& \text { STOP } \\
& \text { END }
\end{aligned}
$$

The solution of the system is

$$
X=((1.0,0.0),(1.0,0.0),(1.0,0.0),(1.0,0.0))
$$

### 5.4. Algorithe

The algorithm used in subroutine CGSLC is the same as that described in subsection 3.4 for CT-matrices except that the solution of G-matrix rather than $T$-matrix systems is involved.

### 5.5. Programing detatis

Programming detatle $o^{\prime}$. subroutine CGSLC are as for CTSLC (see subsection 3.5) except that LINPACK subroutines CGEFA and CGESL (see section 4 of Chapter i) are called instead of subroutines TSLC and TSLCl; the number of multiplications is $M^{2} L^{2}+M^{3} L / 3$ plus terms of lesser degree $-M^{2} L^{2}$ in SALWC (first call), $M^{3}{ }^{2} / 3$ in CGEFA, and lesser amounts in CGESL and further calls of SALWC.

## CHAPTER 3: TORPLITZ- AND CIRCULANT-TYPE MATRICES OF THE THIRD LEVEL

## 1. Structure and Representation

### 1.1. Overview

A matrix

$$
A=\left|\begin{array}{cccccc}
A_{11} & A_{12} & A_{13} & \cdots & \cdot & A_{1 K}  \tag{1}\\
A_{21} & A_{22} & A_{23} & \cdots & \cdot & A_{2 K} \\
A_{31} & A_{32} & A_{33} & \cdots & \cdot & A_{3 K} \\
\cdot & \cdot & \cdot & \cdots & \cdot & \cdot \\
A_{K 1} & A_{K 2} & A_{K 3} & \cdots & \cdot & A_{K K}
\end{array}\right|
$$

with $K$ elements in a row (or column) where the elements $A_{1 j}$ are two-level matrices (see Chapter 2) with rirst-level order $L$ and second-level order $M$ is called a three-level matrix. $K$ is called the first-level order, $L$ becomes the second-level order, and $M$ becomes the third-level order of the matrix $A$. The order N of A is then the product of the orders of its levels: $\mathrm{N}=\mathrm{K} \mathrm{K}_{\mathrm{L}} \mathrm{*} \mathrm{M}$.

We will call the three-level matrix (1) an XYZ-type if A considered as a block matrix is an $X$-type and each of its blocks $A_{i j}$ is a YZ-type (see section 1 of Chapter 2). As $\mathrm{X}-, \mathrm{Y}$-, and Z-types in the TOEPLITZ package we consider T-, C-, and G-matrices defined in section 1 of Chapter 1 . Examples of threelevel matrices can be found in subsections $2.3,3.3,4.3$, and 5.3 of this chapter.

By permuting corresponding rows and columns, we can transform an XYZ-type to any of types XZY, YXZ, YZX, ZXY, or ZYX (see Tyrtyshnikov [25]). This circumstance allows us to limit consideration to a few among the possibie three-level types.

The scheme for compact representation of three-level matrices is the following. Let $A$ be of XYZ-type with level orders $K$, $L$, and $M$, respectively. Furthermcre, let $\tilde{K}$ be the number of elements required in the compact representation of $X$, and $\tilde{M} \star \tilde{L}$ be the number of elements required in the compact representation of a YZ-type with level orders $L$ and $M$. Recail that for TG-,

CT-, CC-, and CG-matrices described in section 1 of Chapter 2 the values of $\tilde{M} * \tilde{L}$ are, respectively, $M * * 2 *(2 * L-1)$, ( $2 * M-1$ ) $L L, M * L$, and $M * * 2 * L$. Ir the TOEPLITZ package such a three-level matrix is represented by a doubly subscripted $\tilde{M} * \tilde{L}$ by $\tilde{K}$ array. The blocks in the array are indexed by the second subscript and ordered in accordance with the $X$-type compact representation. In turn, the elements in a block are indexed by the first subscript and ordered in accordance with the block's YZ-type compact representation packed linearly by columns.

### 1.2. GTG-matrices

A complex matrix
is called a CTG-matrix if $A_{i}, 1=0,1,2, \ldots, K-1$, are TG-matrices of first-level order $L$ and second-level order $M$ (see subsection 1.2 of Chapter 2).
11. the TOEPLITZ package this CTG-matrix is represented by a doubly subscripted $M \star * 2 *(2 * L-1)$ by $K$ array in which the first-level blocks are ordered in the following way:

$$
A_{0}, A_{1}, A_{2}, \ldots, A_{K-1}
$$

Each block $A_{i}$ is a TG-matrix packed linearly by columns.

### 1.3. CCT-mattices

A matrix of form (2) is called a CCT-matrix if $A_{1}, i=0,1,2, \ldots, K-1$, are CT-matrices of first-level order $L$ and second-level order $M$ (see subsection 1.3 of Chapter 2).

In the TOEPLITZ package this CCT-matrix is represented by a doubly subscripted $(2 * M-1) \star L$ by $K$ array. The storage arrangement for the CCT-matrix is as for the CTG-matrix except that each block $A_{i}$ is a CT-matrix.

### 1.4. CCC-matrices

A matrix of form (2) is called a CCC-matrix if $A_{1}, 1=0,1,2, \ldots, K-1$, are CC-matrices of first-level order $L$ and second-level order $M$ (see subsection 1.4 of Chapter 2).

In the TOEPLITZ package this CCC-matrix is ropresented by a doubly subscripted $M * L$ by $K$ array. The storage arrangement for the CCC-matrix is as for the CTG-matrix except that each block $A_{f}$ is a CC-matrix.

### 1.5. COC-matrices

A matrix of form (2) is called a CCG-matrix if $A_{1}, i=0,1,2, \ldots, K-1$, are CG-matrices of first-level order $L$ and second-level order $M$ (see sutsection 1.5 of Chapter 2).

In the TOEPLITZ package this CCG-matrix is represented by a doubly subscripted $M * * 2 \star \mathrm{~L}$ by K array. The storage arrangement for the CCG-matrix is as for the CTG-matrix except that each block $A_{1}$ is a CG-matrix.

### 1.6. Other types of three-1evel matrices

CGT-, TCG-, TGC-, GCT-, GTC-, CTC-, TCC-, CGC-, and GCC-matrices, defined in analogous ways, can be transformed to the types discussed in subsections 1.2-1.5 by permuting corresponding levels. Therefore, the TOEPLITZ package does not include subroutines for solving linear systems with three-level matrices of these types. At the present time no algorithm is known that capitalizes eifectively on the structure of linear systems with three-level matrices more than one of whose levels is of $T$ - or G-type.

## 2. Solution of Linear Equations with CTG-matrices

### 2.1. Purpose

The TOEPLITZ subroutine in this section is designed to solve linear algebraic equations with CTG-watrices, that is, complex block-circulant matrices whose blocks are TG-matrices. A double precision version is also avaflable.
2.2. Usage

Single precision CTG-matrices. CTGSLC solves a linear system with a CTG-matrix. The calling sequence is

CALL CTGSLC(A,X,R,M,L,K,LDA) •
On entry,

A is a doubly subscripted $M * * 2 *(2 * L-1)$ by $K$ array which contains the CTG-matrix in the form described in subsection l.2. A is destroyed by CTGSLC.
$X \quad$ is a singly subscripted array of $M * L * K$ elements which contains the right hand side of the system.
$R \quad$ is a singly subscripted array of $\max (2 * M * * 2 * L+3 * M * * 2+M, 2 * K)$ elements used for work space.

M is thr order of each inner G-matrix block of $A$.

L is the number of inner blocks in each row or column of the TG-matrices which comprise the outer blocks of $A$.
$K$ is the number of outer blocks in each row or column of $A$.

LDA is the leading dimension of the array $A$.

On return,

X contains the solution of the system.

Double precision CTG-matrices. In those computing systems where it is available, the calling sequence of the double precision CTG-matrix subroutine CTGSLZ is the same as that of CTGSLC with $A, X$, and $R$ DOUBLE COMPLEX variables.

### 2.3. Example

The following program segment illustrates the use of the single precision subroutine CTGSLC for CTG-matrices. An example of the use of CTGSIZ could be obtained by changing the subroutine name and type declaration. The system is of order 8 with coefficients as fullows.
$A=\left(\begin{array}{cc|cc|cc|cc}1+i & 3+3 i & 5+5 i & 7+7 i & 13+13 i & 15+15 i & 17+17 i & 19+19 i \\ 2+2 i & 4+4 i & 6+6 i & 8+8 i & 14+14 i & 16+16 i & 18+18 i & 20+20 i \\ \hline 9+9 i & 11+11 i & 1+i & 3+3 i & 21+21 i & 23+23 i & 13+13 i & 15+15 i \\ 10+10 i & 12+12 i & 2+2 i & 4+4 i & 22+22 i & 24+24 i & 14+14 i & 16+16 i \\ \hline 13+13 i & 15+15 i & 17+17 i & 19+19 i & 1+i & 3+3 i & 5+5 i & 7+7 i \\ 14+14 i & 16+16 i & 18+18 i & 20+20 i & 2+2 i & 4+4 i & 6+6 i & 8+8 i \\ \hline 21+21 i & 23+23 i & 13+13 i & 15+15 i & 9+9 i & 11+11 i & 1+i & 3+3 i \\ 22+22 i & 24+24 i & 14+14 i & 16+16 i & 10+10 i & 12+12 i & 2+2 i & 4+4 i\end{array}\right) \quad x=\left(\begin{array}{c}80+80 i \\ 88+88 i \\ 96+96 i \\ 104+104 i \\ 80+80 i \\ 88+88 i \\ 96+96 i \\ 104+104 i\end{array}\right)$

```
    COMPLEX A(12,2),X(8),R(30)
    INTEGER M,L,K,LDA,I,J
    DATA A(1,1)/(1.0,1.0)/,A(2,1)/(2.0,2.0)/,A(3,1)/(3.0,3.0)/,
* A(4,1)/(4.0,4.0)/,A(5,1)/(5.0,5.0)/,A(6,1)/(6.0,6.0)/,
* A(7,1)/(7.0,7.0)/,A(8,1)/(8.0,8.0)/,A(9,1)/(9.0,9.0)/,
* A(10,1)/(10.0,10.0)/,A(11,1)/(11.0,11.0)/,A(12,1)/(12.0,12.0)/,
* A(1,2)/(13.0,13.0)/,A(2,2)/(14.0,14.0)/,A(3,2)/(15.0,15.0)/,
* A(4,2)/(16.0,16.0)/,A(5,2)/(17.0,17.0)/,A(6,2)/(18.0,18.0)/,
* A(7,2)/(19.0,19.0)/,A(8,2)/(20.0,20.0)/,A(9,2)/(21.0,21.0)/,
* A(10,2)/(22.0,22.0)/,A(11,2)/(23.0,23.0)/,A(12,2)/(24.0,24.0)/
    DATA X(1)/(80.0,80.0)/,X(2)/(88.0,88.0)/,X(3)/(96.0,96.0)/,
* X(4)/(104.0,104.0)/,X(5)/(80.0,80.0)/,X(6)/(88.0,88.0)/,
* X(7)/(96.0,96.0)/,X(8)/(104.0,104.0)/
    M = 2
    L =2
    K - 2
    LDA = 12
    CALL CTGSLC(A,X,R,M,L,K,LDA)
    J = M*L*K
    D0 10 I = 1, J
        WRITE(...,...) X(I)
10 CONTINUE
    S'TOP
    END
```

The solution of the system is

$$
\begin{aligned}
X= & ((1.0,0.0),(1.0,0.0),(1.0,0.0),(1.0,0.0) \\
& (1.0,0.0),(1.0,0.0),(1.0,0.0),(1.0,0.0))
\end{aligned}
$$

### 2.4. Algorithm

The algorithm for solving a linear system

$$
\begin{equation*}
A x=b \tag{1}
\end{equation*}
$$

with the CTG-matrix

$$
A=\left(\begin{array}{lllllll}
A_{0} & A_{1} & A_{2} & \cdots & \cdot & A_{K-1} \\
A_{K-1} & A_{0} & A_{1} & \cdots & \cdot & \cdot & A_{K-2} \\
A_{K-2} & A_{K-1} & A_{0} & \cdots & \cdot & \cdot & A_{K-3} \\
\cdot & \cdot & \cdot & \cdot & \cdot & \cdot & \cdot \\
A_{1} & A_{2} & A_{3} & \cdots & \cdot & \cdot & A_{0}
\end{array}\right)
$$

where $A_{i}, i=0,1,2, \ldots, K-1$, are $T G$ matrices of first-level order $L$ and secondlevel order $M$, is analogous to that described in subsection 3.4 of Chapter 2 for CT-matrices. It proceeds from a similarity transformation of $A$ to a block-diagonal matrix

$$
D=Q^{\star} A Q
$$

in which each diagonal block is a TG-matrix. (The symbol * denotes conjugate transpose.) $Q$ is a two-level matrix with first-level order $K$ and second-level order L*M whose blocks are scalar matrices; the matrix of the scalars themselves is unitary.

B1cek $Q_{i j}$ of $Q$ is defined as

$$
Q_{1 j}=E^{(1-1) \cdot(j-1)_{\star I / \sqrt{K}}},
$$

where $E=\exp (2 \pi \sqrt{-1} / K)$ and $I$ is the identity matrix of order L*M. However, as for C-matrices (see section 3 of Chapter 1), it is more efficient to use instead the matrix

$$
\bar{Q}=\sqrt{K} Q .
$$

Thus the solution $x$ of the system (1) can be found by the following steps:
a) Transform the matrix A to the block-diagonal matrix

$$
D=\bar{Q}^{\star} A \bar{Q} / K .
$$

b) Transform the right hand side

$$
y=\overparen{Q}^{-t} b .
$$

c) Solve the system

$$
\mathrm{Dz}=\mathrm{y} .
$$

d) Transform the vector $z$ back to

$$
x=\bar{Q} z / K
$$

Note that since $A$ is a CTG-matrix, its transformation to $D$ becomes simply

$$
D_{i i}=\sum_{j=1}^{K} \bar{Q}_{i j}^{A_{j-1}}
$$

where $D_{i i}$ is the i-th diagonal block of $D$ and $A_{j-1}$ is the outer block with index j-1 at the top of A. Furthermore, since $D$ is block-diagonal each block of which is a TG-matrix, the system (1) reduces to $K$ systems with TG-matrices.

### 2.5. Prograning details

The implementation of subroutine CTGSLC corresponds to the algorithm described in subsection 2.4. All needed operations with matrices $\bar{Q}$ and $\vec{Q}^{*}$ are implemented by the service subrouting SALWC described in subsection 3.5 of Chapter 2. For solving the $K$ systems with TG-matrices, second-level subroutines TGSLC and TGSLCl are called.

For solving systems with double precision CTG-matrices (using CTGSLZ), corresponding versions of subroutines TGSLC, TGSLC1, and SALWC are called, namely, TGSLZ, TGSLZ1, and SALWZ.

The number of multiplications in executing subroutine CTGSLC is $2 M^{3} L^{2} K+$ $2 M^{2} L^{2}$ plus terms of lesser degree.

## 3. Solution of Linear Equations with CCT-matrices

### 3.1. Purpose

The TOEPLITZ subroutine in this section is designed to solve linear algebraic equations with CCT-matrices, that is, complex block-circulant matrices whose blocks are CT-matrices. A double precision version is also available.
3.2. Jsage

Single precision CCT-matrices. CCTSLC solves a linear system with a CCT-matrix. The calling sequence is

CALL CCTSLC(A,X,R,M,L,K,LDA) •
On entry,

A is a doubly subscripted ( $2 * M-1$ ) *L by $K$ array which contains the CCT-matrix in the form described in subsection 1.3. A is destroyed by CCTSLC.
$X \quad$ is a singly subscripted array of $M * L * K$ elements which contains the right hand side of the system.
$R$ is a singly subscripted array of $\max (2 * M-2,2 * L, 2 * K)$ elements used for work space.

M is the order of each inner T-matrix block of $A$.

L is the number of inner blocks in each row or column of the CT-matrices which comprise the outer blocks of $A$.
$K$ is the number of outer blocks in each row or column of $A$.

LDA is the leading dimension of the array $A$.

On return,

Double precision CCI-matrices. In those computing systems where it is available, the calling sequence of the double precision CCT-matrix subroutine CCTSLZ is the same as that of CCTSLC with $A, X$, and $R$ DOUBLE COMPLEX variables.

### 3.3. Example

The following program segment illustrates the use of the single precision subroutine CCTSLC for CCT-matrices. An example of the use of CCTSLZ could be obtained by changing the subroutine name and type declaration. The system is of order 8 with coefficients as follows.


```
    COMPLEX A(6,2),X(8),R(4)
    INTEGER M,L,K,LDA,I,J
    DATA A (1,1)/(1.0,1.0)/,A(2,1)/(2.0,2.0)/,A(3,1)/(3.0,3.0)/,
* A(4,1)/(4.0,4.0)/,A(5,1)/(5.0,5.0)/,A(6,1)/(6.0,6.0)/,
* A(1,2)/(7.0,7.0)/,A(2,2)/(8.0,8.0)/,A(3,2)/(9.0,9.0)/,
* A(4,2)/(10.0,10.0)/,A(5,2)/(11.0,11.0)/,A(6,2)/(12.0,12.0)/
    DATA X(1)/(48.0,48.0)/,X(2)/(52.0,52.0)/,
* X(3)/(48.0,48.0)/,X(4)/(52.0,52.0)/,
* X(5)/(48.0,48.0)/,X(6)/(52.0,52.0)/,
* X(7)/(48.0,48.0)/,X(8)/(52.0,52.0)/
    M=2
    L=2
    K=2
    LDA=6
    CALL CCTSLC(A,X,R,M,L,K,LDA)
    J = M*L*K
    DO 10 I = 1, J
        WRITE(....,...) X(I)
        10 CONTINUE
            STOP
    END
```

The solution of the system is

$$
\begin{aligned}
\mathrm{X}=( & (1.0,0.0),(1.0,0.0),(1.0,0.0),(1.0,0.0) \\
& (1.0,0.0),(1.0,0.0),(1.0,0.0),(1.0,0.0))
\end{aligned}
$$

### 3.4. Algoritha

The algorithm used in subroutine CCTSLC is the same as that described in subsection 2.4 for CTG-matrices except that the solution of CT-matrix rather than TG-matrix systems is involved.

### 3.5. Programing details

Programming details of subroutine CCTSLC are as for CTGSLC (see subsection 2.5) except that subroutine CTSLC is called instead of subroutines TGSLC and TGSLCl; the number of multiplications is approximately $4 M L K^{2}+4 M L{ }^{2} K+3 M^{2} L K$.

## 4. Solution of Linear Equations with CCC-matrices

### 4.1. Purpose

The TOEPLITZ subroutine in this section is designed to solve linear algebraic equations with CCC-matrices, that is, complex block-circulant matrices whose blocks are CC-matrices. A double precision version is also available.

### 4.2. Usage

Single precision CCC-matrices. CCSLC solves a linear system with a CCC-matrix. The calling sequence is

CALL CCCSLC(A, X,R,M,L,K,LDA) •
On entry,

A is a doubly subscripted $M * L$ by $K$ array which contains the CCC-matrix in the form descrited in subsection 1.4. A is destroyed by CCCSLC.
$X \quad$ is a singly subscripted array of $M * L * K$ elements which contains the right hand side of the system.
$R$ is a siugly subscripted array of max $(M, 2 * L, 2 \star K)$ elements used for work space.

M 15 the order of each inner C-matrix block of $A$.

L is the number of inner blocks in each row or column of the CC-matrices which comprise the outer blocks of $A$.
$K$ is the number of outer blocks in each row or coluinn of $A$.

LDA is the leading dimension of the array $A$.

On return,
$X$ contains the solution of the system.

Double precision CCC-matrices. In those computing systems whre it is available, the calling sequence of the double precision CCC-matrix subroutine CCCSLZ is the same as that of CCCSLC with $A$, $X$, and $R$ DOUBLE COMPLEX variables.

### 4.3. Exarple

The following program segment fllustrates the use of the single precision subroutine CCCSLC for CCC-matrices. An example of the use of CCCSLZ could be obtained by changing the subroutine name and type declaration. The system is of order 8 with coefficients as follows.


```
    COMPLEX A(4,2), X(8),R(4)
    INTEGER M,L,K,LDA,I,J
    DATA \(\mathrm{A}(1,1) /(1.0,1.0) /, \mathrm{A}(2,1) /(2.0,2.0) /, \mathrm{A}(3,1) /(3.0,3.0) /\),
* \(\mathrm{A}(4,1) /(4.0,4.0) / \mathrm{A}(1,2) /(5.0,5.0) /, \mathrm{A}(2,2) /(6.0,6.0) /\),
* \(\quad \mathrm{A}(3,2) /(7.0,7.0) /, \mathrm{A}(4,2) /(8.0,8.0) /\)
    DATA \(\mathrm{X}(1) /(36.0,36.0) /, \mathrm{X}(2) /(36.0,36.0), \mathrm{X}(3) /(36.0,36.0) /\),
* \(\mathrm{X}(4) /(36.0,36.0) /, \mathrm{X}(5) /(36.0,36.0), \mathrm{x}(6) /(36.0,36.0) /\),
* \(\mathrm{X}(7) /(36.0,36.0) /, \mathrm{x}(8) /(36.0,36.0) /\)
    \(M=2\)
    \(\mathrm{L}=2\)
    \(K=2\)
    \(\mathrm{LDA}=4\)
    CALL CCCSLC(A, X,R,M,L,K,LDA)
    \(\mathrm{J}=\mathrm{M} \mathrm{L}_{\mathrm{L}} \mathrm{*}_{\mathrm{K}}\)
    Do \(10 \mathrm{I}=1, \mathrm{~J}\)
        WRITE(...,....) X(I)
10 Continue
    STOP
    END
```

The solution of the system is

$$
\begin{aligned}
\mathrm{x}= & ((1.0,0.0),(1.0,0.0),(1.0,0.0),(1.0,0.0), \\
& (1.0,0.0),(1.0,0.0),(1.0,0.0),(1.0,0.0)) .
\end{aligned}
$$

### 4.4. Algorithm

The algorithm used in subroutine CCCSLC is the same as that described in subsection 2.4 for CTG-matrices except that the solution of CC-matrix rather than TG-matrix systems is involved.

### 4.5. Prograting details

Programming details of subroutine CCCSLC are as for CTGSLC (see subsection 2.5) except that subroutine CCSLC is called itistead of subroutines TGSLC and TGSLCl; the number of multiplications is approximately $3 \mathrm{MLK}^{2}+3 \mathrm{ML}^{2} \mathrm{~K}+3 \mathrm{M}^{2} \mathrm{LK}$.

## 5. Solution of Linear Equations with CCG-matrices

### 5.1. Purpose

The TOEPLITZ subroutine in this section is designed to solve linear algebraic equations with CCG-matrices, that is, complex block-circulant matrices whose blocks are CG-matrices. A double precision version is also available.
5.2. Usage

Single precision CCG-matrices. CCGSLC solves a linear system with a CCG-matrix. The calling sequence is

CALL CCGSLC(A,X,R,M,L,K,LDA) .
On entry,

A is a doubly subscripted $M * * 2 \star L$ by $K$ array which contains the CCG-matrix in the form described in subsection 1.5. A is destroyed by CCGSLC.
$X \quad$ is a singly subscripted array of $M * L * K$ elements which contains the right hand side of the system.
$R \quad$ is a singly subscripted array of $\max (M, 2 * L, 2 * K)$ elements used for work space.
$M$ is the order of each inner G-matrix block of $A$.

L is the number of inner blocks in each row or column of the CG-matrices which comprise the outer blocks of A.

K is the number of outer blocks in each row or column of $A$.

LDA is the leading dimension of the array $A$.

On return,
$X$ contains the solution of the system.

Double precision CCG-matrices. In those computing systems where it is available, the calling sequence of the double precision CCG-matrix subroutine CCGSLZ is the same as that of CCGSLC with $A, X$, and $R$ DOUBLE COMPLEX variables.

### 5.3. Example

The following program segment illustrates the use of the single precision subroutine CCGSLC for CCG-matrices. An example of the use of CCGSLZ could be obtained by changing the subroutine name and type declaration. The system is of order 8 with coefficients as follows.
$A=\left(\begin{array}{cc|cc|cc|cc}1+i & 3+3 i & 5+5 i & 7+7 i & 9+9 i & 11+11 i & 13+13 i & 15+15 i \\ 2+2 i & 4+4 i & 6+6 i & 8+8 i & 10+10 i & 12+12 i & 14+14 i & 16+16 i \\ \hline 5+5 i & 7+7 i & 1+i & 3+3 i & 13+13 i & 15+15 i & 9+9 i & 11+11 i \\ 6+6 i & 8+8 i & 2+2 i & 4+4 i & 14+14 i & 16+16 i & 10+10 i & 12+12 i \\ \hline 9+91 & 11+11 i & 13+13 i & 15+151 & 1+1 & 3+3 i & 5+5 i & 7+7 i \\ 10+10 i & 12+12 i & 14+14 i & 16+16 i & 2+2 i & 4+4 i & 6+6 i & 8+8 i \\ \hline 13+13 i & 15+15 i & 9+9 i & 11+11 i & 5+5 i & 7+7 i & 1+i & 3+3 i \\ 14+14 i & 16+16 i & 10+10 i & 12+12 i & 6+6 i & 8+8 i & 2+2 i & 4+4 i\end{array}\right) \quad x=\left(\begin{array}{l}64+64 i \\ 72+72 i \\ 64+64 i \\ 72+72 i \\ 64+64 i \\ 72+721 \\ 64+64 i \\ 72+72 i\end{array}\right)$

```
COMPLEX A(8,2),X(8),R(4)
    INTEGER M,L,K,LDA,I,J
    DATA A(1,1)/(1.0,1.0)/,A(2,1)/(2.0,2.0)/,A(3,1)/(3.0,3.0)/,
    * A(4,1)/(4.0,4.0)/,A(5,1)/(5.0,5.0)/,A(6,1)/(6.0,6.0)/,
    * A(7,1)/(7.0,7.0)/,A(8,1)/(8.0,8.0)',}\textrm{A}(1,2)/(9.0,9.0)/, 
    * A(2,2)/(10.0,10.0)/,A(3,2)/(11.0,11.0)/,A(4,2)/(12.0,12.0)/,
    * A(5,2)/(13.0,13.0)/,A(6,2)/(14.0,14.0)/,A(7,2)/(15.0,15.0)/,
    * A(8,2)/(16.0,1E 0)/
    DATA X(1)/(64.0,64.0)/,X(2)/(72.0,72.0)/,
    * X(3)/(64.0,64.0)/,X(4)/(72.0,72.0)/,
    * X(5)/(64.0,64.0)/,X(6)/(72.0,72.0)/,
    * X(7)/(64.0,64.0)i,X(8)/(72.0,72.0)/
    M=2
    L=2
    K=2
    LDA = 8
    CALL CCGSLC(A,X,R,M,L,K,LDA)
    J = M*L*K
    DO 10 I = 1, J
    WRITE(.......) X(I)
10 CONTINUE
    STOP
    END
```

The solution of the system is

$$
\begin{aligned}
\mathrm{x}=( & (1.0,0.0),(1.0,0.0),(1.0,0.0),(1.0,0.0) \\
& (1.0,0.0),(1.0,0.0),(1.0,0.0),(1.0,0.0)) .
\end{aligned}
$$

### 5.4 Algoritha

The algorithm used in subroutine CCGSLC is the same as that described in subsection 2.4 for CTG-matrices except that the solution of CG-matrix rather i: An TG-matrix systems is involved.

### 5.5. Programing details

Programming details of subroutine CCGSLC are as for CTGSLC (ser subsection 2.5) except that subroutine CGSLC is called instead of subroutines TGSLC and TGSLCl; the number of maltiplications is $M^{2} L^{2}+M^{2} L^{2} K+M^{3} L K / 3$ plus terms of lesser degree.

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## APPEADIX A. TABLES OF EXECUTION TIES


#### Abstract

We provide here three tables of sample execution times for the TOEPLITZ package subroutines. The first two tables report times for the single precision and double precision versions, respectively, on the VAX 11/780; the third table reports times for the single precision version on the IBM 3033. The VAX compilations were made with the Foriran 77 compiler running under UNIX; the IBM compilations were made with the Fortran H Extended (Enhanced) compiler running under MVS. Using these tables and the approximate multiplication counts given in the discussions of the algorithms in the previous chapters, it should be possible to extrapolate execution times for problems of different dimensions.


## SUMAARY OF EXECUTION TIMES POR THE <br> SINGLE PRECISION TOEFLITZ SUBROUTINES ON THE VAX 11/780

| SUBROUTINE | $3^{\text {rd }}$ LEVEL | $2^{\text {nd }}$ Level | $1^{\text {st }}$ LEVEL | TIME (sec.) |
| :---: | :---: | :---: | :---: | :---: |
| TSLS (TSLS ${ }^{\text {) }}$ |  |  | 100 | 0.35 |
| TSLC (TSLCl) |  |  | 100 | 1.6 |
| CSLC |  |  | 100 | 1.2 |
| CQRS |  | 100(rows) | 20(columns) | 0.25 |
| CQRC |  | 100(rows) | 20(columns) | 0.83 |
| TGSLS (TGSLS 1 ) |  | 10 | 10 | 5.9 |
| 'GGSLS (IN-LINE SAXPY) |  | 10 | 10 | 6.7 |
| TGSLC (TGSLC1) |  | 10 | 10 | 14.5 |
| TGSLC (IN-LINE CAXPY) |  | 10 | 10 | 16.6 |
| CTSLC |  | 20 | 20 | 2.7 |
| CCSLC |  | 20 | 20 | 2.2 |
| CGSLC |  | 20 | 20 | 13.5 |
| CTGSLC | 6 | 6 | 6 | 9.2 |
| CCTSLC | 8 | 8 | 8 | 2.5 |
| CCCSLC | 8 | 8 | 8 | 2.0 |
| CCGSLC | 8 | 8 | 8 | 5.7 |

SUMMARY OF EXBCUTION TIMES FOR THE
dOUBLE PRECISION TOEPLITZ SUBROUTINES ON THE VAX 11/780

| SUBROUTINE | $3^{\text {rd }}$ LEVEL | $2^{\text {nd }}$ Level | $1^{\text {st }}$ Level | TIMF (sec.) |
| :---: | :---: | :---: | :---: | :---: |
| TSLD (TSLDI) |  |  | 100 | 0.52 |
| TSLZ (TSLZ ${ }^{\text {P }}$ |  |  | 100 | 2.5 |
| CSLZ |  |  | 100 | 1.9 |
| CQRD |  | 100(rows) | 20(columns) | 0.38 |
| CQRZ |  | 100(rows) | 20(columns) | 1.6 |
| TGSLD (TGSLDI) |  | 10 | 10 | 8.5 |
| TGSLZ (TGSLZ ${ }^{\text {( }}$ ) |  | 10 | 10 | 27.3 |
| CTSLZ |  | 20 | 20 | 4.0 |
| CCSLZ |  | 20 | 20 | 3.3 |
| CGSLZ |  | 20 | 20 | 22.2 |
| CTGSLZ | 6 | 6 | 6 | 14.8 |
| CCTSLZ | 8 | 8 | 8 | 3.7 |
| CCCSLZ | 8 | 8 | 8 | 3.2 |
| CCGSLZ | 8 | 8 | 8 | 9.4 |

## SUMIARY OF EXECUTION TIAES POR THE SINGLE PRECISION TORPLITZ SUBROUTINES ON TEE IBH 3033

| SUBROUTINE | $2^{\text {nd }}$ Level | $1{ }^{\text {st }}$ Level | TIME (sec.) |
| :---: | :---: | :---: | :---: |
| TSLS (TSLS1) |  | 100 | . 028 |
|  |  | 200 | . 11 |
|  |  | 300 | . 25 |
|  |  | 400 | . 44 |
|  |  | 500 | . 69 |
| TSLC (TSLC1) |  | 100 | . 19 |
| CSLC |  | 100 | . 18 |
|  |  | 200 | . 70 |
|  |  | 300 | 1.6 |
|  |  | 400 | 2.8 |
|  |  | 500 | 4.3 |
| TGSLS (TGSLSI) | 10 | 10 | . 47 |
| TGSLS (IN-LINE SAXPY) | 10 | 10 | . 27 |
| TGSLC (TGSLCl) | 10 | 10 | 1.5 |
| TGSLC (IN-LINE CAXPY) | 10 | 10 | 1.4 |
| CTSLC | 20 | 20 | . 35 |
| CCSLC | 20 | 20 | . 31 |
| CGSLC | 20 | 20 | 1.5 |

## APPENDIX B. PROGRAH LISTINGS

There follows the single precision version of the TOEPLITZ package program listings; both single precision and double precision versions of the subprograms are available with the TOEPLITZ package. The listings appear in the following order:

TSLS, TSLS1, TSLC, TSLC1, CSLC, CQRS, CQRC, TGSLS, TGSLS 1, TGSLC, TGSLCl, CTSLC, CCSLC, CGSLC, SALWC, CTGSLC, CCTSLC, CCCSLC, CCGSLC.

```
SUBROUTINE TSLS(A,X,R,M)
INTEGER M
REAL A(1),X(M),R(1)
TSLS CALLS TSLS1 TO SOLVE THE REAL LINEAR SYSTEM
A}%\textrm{X}=\textrm{B
WITH THE T - MATRIX A .
ON ENTRY
    A REAL(2*M - 1)
            THE FIRST ROW OF THE T - MATRIX FOLLOWED BY ITS
            FIRST COLUMN BEGINNING WITH THE SECOND ELEMENT .
            ON RETURN A IS UNALTERED .
    X RENL(M)
            IFIL RIGHT HAND SIDE VECTOR B .
            REAL(2*M - 2)
            A WORK VECTOR .
            INTEGER
            THE ORDER OF THE MATRIX A.
ON RETURN
    X THE SOLUTION VECTOR .
TOEPLITZ PACKAGE. THIS VERSION DATED 07/23/82.
SUBROUTINES AND FUNCTIONS
    'IOEPLITZ PACKAGE ... TSLS1
CALL SUBROUTINE TSLSI
CALL TSLS1(A,A(M+1),X,X,R,R(M),M)
RETURN
END
```

```
SUBROUTINE TSLS1(A1,A2,B,X,C1,C2,M)
INTEGER M
REAL A1(M),A2(1), B(M),X(M),C1(1),C2(1)
```

INTEGER I1,I2,N,N1,N2
REAL R1,R2,R3,R5,R6
$A * X=B$
WITH THE T - MATRIX A.
ON ENTRY
A1 REAL(M)

A2 $\quad \operatorname{REAL}(M-1)$

B $\quad \operatorname{REAL}(M)$

C1 REAL(M - 1) A WORK VECTOR .

C2 $\operatorname{REAL}(M-1)$ A WORK VECTOR .

M INTEGER

ON RETURN
$X \quad$ REAL (M)

INTERNAL VARIABLES
$\mathrm{R1}=\mathrm{A} 1(1)$
$X(1)=B(1) / R 1$
IF (M .EQ. 1) GO TO 80

WITH THE T - MATRIX FOR $\mathrm{N}=2, \mathrm{M}$.

TSLS1 SOLVES THE REAL LINEAR SYSTEM
THE FIRST ROW OF THE T - MATRIX A .
ON RETURN A1 IS UNALTERED .

THE FIRST COLUMN OF THE T - MATRIX A beginning with the second element . ON RETURN A2 IS UNALTERED .

THE RIGHT HAND SIDE VECTOR . ON RETURN B IS UNALTERED . THE ORDER OF THE MATRIX A.

THE SOLUTION VECTOR. X MAY COINCIDE WITH B .
TOEPLITZ PACKAGE. THIS VERSION DATED 07/23/82.

SOLVE THE SYSTEM WITH THE PRINCIPAL MINOR OF ORDER 1 .

RECURRENT PROCESS FOR SOLVING THE SYSTEM

C

```
C COMPUTE MULTIPLES OF THE FIRST AND LAST COLUMNS OF
C THE INVERSE OF THE PRINCIPAL MINOR OF ORDER N .
C
    N1 = N - 1
    N2 = N - 2
    R5 = A2(N1)
    R6 = Al(N)
    IF (N .EQ. 2) GO TO 20
        C1(N1) = R2
        DO 10 I1 = 1, N2
            I2 = N - I1
            R5 = R5 + A2(I1)*C1(I2)
            R6 = R6 + A1(I1+1)*C2(I1)
        CONTINUE
    10
    20
    CONTINUE
    R2 = -R5/R1
        R3 = -R6/R1
        R1 = R1 + R5*R3
        IF (N .EQ. 2) GO TO 40
        R6 = C2(1)
        C2(N1) = 0.0E0
        DO 30 I1 = 2, N1
            R5 = C2(I1)
            C2(I1) = C1(I1)*R3 + R6
            C1(I1) = C1(I1) + R6*R2
            R6 = R5
        CONTINUE
    30
    40 CONTINUE
        C2(1) = R3
            COMPUTE THE SOLUTION OF THE SYSTEM WITH THE
            PRINCIPAL MINOR OF ORDER N .
        R5 = 0.0E0
        DO 50 I1 = 1, N1
        I2 = N - I1
        R5 = R5 + A2(I1)*X(I2)
    50 CONTINUE
        R6 = (B(N) - R5)/R1
        D0 60 I1 = 1, N1
        X(I1) = X(I1) + C2(I1)*R6
    60 CONTINUE
        X(N) = R6
    70 CONTINUE
    80 CONTINUE
        RETURN
    END
```

```
SUBROUTINE TSLC(A,X,R,M)
INTEGER M
COMPLEX A(1),X(M),R(1)
TSLC CALLS TSLC1 TO SOIVE THE COMPLEX LINEAR SYSTEM
\(A * X=B\)
WITH THE T - MATRIX A.
ON ENTRY
A \(\quad \operatorname{COMPLEX}(2 * M-1)\)
THE FIRST ROW OF THE T - MATRIX FOLLOWED BY ITS FIRST COLUMN BEGINNING WITH THE SECOND ELEMENT . ON RETURN A IS UNALTERED .
X COMPLEX (M)
THE RIGHT HAND SIDE VECTOR B .
R COMPLEX(2*M - 2) A WORK VECTOR .
M INiEGER THE ORDER OF THE MATRIX A.
ON RETURN
X THE SOLUTION VECTOR .
TOEPLITZ PACKAGE. THIS VERSION DATED 07/23/82.
SUBROUTINES AND FUNCTIONS
TOEPLITZ PACKAGE ... TSLC1
CALL SUBROUTINE TSLC1
CALL T'SLC1 (A, \(A(M+1), X, X, R, R(M), M)\)
RETURN
END
```

```
SUBROUTINE TSLC1(A1,A2,B,X,C1,C2,M)
1NTEGER M
COMPLEX A1(M),A2(1),B(M),X(M),C1(1),C2(1)
```

COMPUTE MULTIPLES OF THE FIRST AND LAST COLUNNS OF THE INVERSE OF THE PRINCIPAL MINOR OF ORDER $N$.
$\mathrm{N} 1=\mathrm{N}-1$
$\mathrm{N} 2=\mathrm{N}-2$
$\mathrm{R} 5=\mathrm{A} 2(\mathrm{~N} 1)$
$\mathrm{R} 6=\mathrm{A} 1(\mathrm{~N})$
IF (N . EQ. 2) GO TO 20
$\mathrm{C} 1(\mathrm{~N} 1)=\mathrm{R} 2$
DO 10 I1 = 1 , N2
$\mathrm{I} 2=\mathrm{N}-\mathrm{I} 1$
$\mathrm{R} 5=\mathrm{R} 5+\mathrm{A} 2(\mathrm{I} 1) * \mathrm{C} 1$ (I2)
$R 6=R 6+A 1(I 1+1) * C 2(I 1)$
CONTINUE
CONTINUE
$\mathrm{R} 2=-\mathrm{R} 5 / \mathrm{R} 1$
R3 $=-R 6 / R 1$
$\mathrm{R} 1=\mathrm{R} 1+\mathrm{R} 5 * \mathrm{R} 3$
IF (N .EQ. 2) GO TO 40
R6 = C2 (1)
$\mathrm{C} 2(\mathrm{~N} 1)=(0.0 \mathrm{E} 0,0.0 \mathrm{E} 0)$
DO $30 \mathrm{I} 1=2, \mathrm{~N} 1$
$\mathrm{R} 5=\mathrm{C} 2(\mathrm{I} 1)$
$\mathrm{C} 2(\mathrm{I} 1)=\mathrm{C} 1(\mathrm{I} 1) * \mathrm{R} 3+\mathrm{R} 6$
$\mathrm{Cl}(\mathrm{I} 1)=\mathrm{C} 1(\mathrm{I} 1)+\mathrm{R} 6 * \mathrm{R} 2$
$R 6=R 5$
CONTINUE
CONTINUE
$\mathrm{C} 2(1)=\mathrm{R} 3$
COMPUTE THE SOLUTION OF THE SYSTEM WITH THE
PRINCIPAL MINOR OF ORDER N.
$\mathrm{R} 5=(0.0 \mathrm{E} 0,0.0 \mathrm{E} 0)$
DO $50 \mathrm{I} 1=1, \mathrm{~N} 1$
$\mathrm{I} 2=\mathrm{N}-\mathrm{I} 1$
$\mathrm{R} 5=\mathrm{R} 5+\mathrm{A} 2(\mathrm{I} 1) \div \mathrm{X}(\mathrm{I} 2)$
50
CONTINUE
$\mathrm{R} 6=(\mathrm{B}(\mathrm{N})-\mathrm{R} 5) / \mathrm{R} 1$
D0 $60 \mathrm{I} 1=1, \mathrm{~N} 1$
$X(I 1)=X(I 1)+C 2(I 1) * R 6$
CONTINUE
$X(N)=R 6$
70 CONTINUE
80 CONTINUE
RETURN
END

```
SUBROUTINE CSLC(A,X,R,M)
INTEGER M
COMPLEX A(M),X(M),R(M)
```

CSLC SOLVES THE COMPLEX LINEAR SYSTEM
$A * X=B$
WITH THE C - MATRIX A .
ON ENTRY
A COMPLEX(M)
THE FIRST ROW OF THE C - MATRIX . ON RETURN A IS UNALTERED .
$\mathrm{X} \quad \operatorname{COMPLEX}(\mathrm{M})$
THE RIGHT HAND SIDE VECTOR B .
R COMPLEX(M)
A WORK VECTOR .
M INTEGER
THE ORDER OF THE MATRIX A .
ON RETURN
X THE SOLUTION VECTOR .
TOEPI,ITZ PACKAGE. THIS VERSION DATED 07/23/82.
SUBROUTINES AND FUNCTIONS
FORTRAN ... CMPLK, COS, FLOAT,SIN
INTERNAL VARIABLES
INTEGER I1,I2
REAL P,RI,RM,V1,V2
COMPLEX E,E1,F,F1,T,T1
$\mathrm{T} 1=\mathrm{X}(1)$
$X(1)=T 1 / A(1)$
IF (M .EQ. 1) GO TO 50
RM $=$ FLOAT ( $M$ )
COMPUTE THE INVERSE DISCRETE FOURIER TRANSFORMATION
OF THE FIRST ROW OF THE MATRIX AND THE DISCRETE FOURIER TRANSFORMATION OF THE RIGHT HAND SIDE VECTOR .
$\mathrm{T}=(0.0 \mathrm{E} 0,0.0 \mathrm{E} 0)$
$\mathrm{RI}=-1.0 \mathrm{E} 0$
DO $20 \mathrm{I} 1=1, \mathrm{M}$
$\mathrm{RI}=\mathrm{RI}+1.0 \mathrm{E} 0$
MINIMIZE ERROR IN FORMING MULTIPLES OF 2*PI

$$
P=((201 . \mathrm{E} 0 / 32 . \mathrm{E} 0) * \mathrm{RI}+1.93530717958647692528 \mathrm{E}-3 * \mathrm{RI}) / \mathrm{RM}
$$

C

$$
V 1=\operatorname{Cos}(P)
$$

$$
\mathrm{V} 2=\operatorname{SIN}(P)
$$

$$
E=\operatorname{CMPLX}(V 1,-V 2)
$$

$$
E 1=\operatorname{CMPLX}(V 1, V 2)
$$

$$
F=A(1)
$$

$\mathrm{F} 1=\mathrm{T} 1$
DO $10 \mathrm{I} 2=2, \mathrm{M}$

$$
F=E \div F+A(I 2)
$$

$$
\mathrm{F} 1=\mathrm{E} 1 * \mathrm{~F} 1+\mathrm{X}(\mathrm{I} 2)
$$

10 CONTINUE
$R(I 1)=(E 1 * F 1) /(E * F)$
$\mathrm{T}=\mathrm{T}+\mathrm{R}(\mathrm{I} 1)$
20 CONTINUE

C
C
C
C

C
C

COMPUTE THE SOLUTION OF THE GIVEN SYSTEM BY
THE INVERSE DISCRETE FOURIER TRANSFORMATION .
$X(1)=T / R M$
$\mathrm{RI}=0.0 \mathrm{EO}$
DO $40 \mathrm{I} 1=2, \mathrm{M}$
$\mathrm{RI}=\mathrm{RI}+1.0 \mathrm{E} 0$
MINIMIZE ERROR IN FORMING MULTIPLES OF $2 * P I$.
$P=((201 . E 0 / 32 . E O) * R I+1.93530717958647692528 E-3 * R I) / R M$
$\mathrm{V} 1=\cos (\mathrm{P})$
$\mathrm{V} 2=\operatorname{SIN}(\mathrm{P})$
$\mathrm{E}=\mathrm{CMPLX}(\mathrm{V} 1,-\mathrm{V} 2)$
$\mathrm{F}=\mathrm{R}(1)$
DO $30 \mathrm{I} 2=2, \mathrm{M}$ $\mathrm{F}=\mathrm{E} * \mathrm{~F}+\mathrm{R}(\mathrm{I} 2)$
30 CONTINUE
$X(\mathrm{I} 1)=\mathrm{E} * \mathrm{~F} / \mathrm{RM}$
40 CONTINUE
50 CONTINUE
RETURN
END

```
SUBROUTINE CQRS(A,Q,S,M,L,LDQ,LDS)
INTEGER M,L,LDQ,LDS
REAL A(M),Q(LDQ,L),S(LDS,L)
```

INTEGER I, J, J1, JI
REAL SCALE,SNRM2
REAL C,SDOT
INITIALIZATION (LAST COLUMN OF Q USED AS WORK VECTOR) .
DO $10 \mathrm{I}=1, \mathrm{M}$
$Q(I, 1)=A(I)$

```
        Q(I,L) = A(I)
    10 CONTINUE
C
C RECURRENT PROCESS FOR THE LATTICE ALGORITHM WITH NORMALIZATION .
C
    DO 70 J1 = 1, L
        J = J1 + 1
        SCALE = 1.0EO/SNRM2 (M,Q(1,J1),1)
        IF (J1 .EQ. L) GO TO 60
        C = -SCALE*(Q(M,J1)*Q(1,L) +
    *
        Q(1,J) = Q(M,J1) + C%Q(1,L)
        DO 20 I = 2,M
            Q(I,J) = Q(I-1,J1) + C%Q(I,L)
        CONTINUE
        IF (J .EQ. L) GO TO 30
            Q(1,L) = Q(1, L) + CrQ (N,J1)
            CALL SAXPY(M-1,C,Q(1,J1),1,Q(2,L),1)
            CONTINUE
            S(1,J) = C
            IF (J .EQ. 2) GO TO 50
                DO 40 I = 2, J1
                    JI = J - I
                        S(I,J) = S(I-1,J1) + C*S(JI,J1)
            CONTINUE
            CONTINUE
        CONTINUE
        CALL SSCAL(M,SCALE,Q(1,J1),1)
        S(J1,J1) = 1.OE0
        CALL SSCAL(J1,SCALE,S (1,J1),1)
7 0 ~ C O N T I N U E ~
    RETURN
    END
```

```
SUBROUTINE CQRC(A,Q,S,M,L,LDQ,LDS)
INTEGER M,L,LDQ,LDS
COMPLEX A(M),Q(LDQ,L),S(LDS,L)
CQRC COMPUTES THE QR FACTORIZATION IN THE FORM \(A * R(\) INVERSE \()=Q\)
OF THE COMPLEX COLUMN-CIRCULANT MATRIX A .
ON ENTRY
A COMPLEX(M)
THE FIRST COLUMN OF THE COLUMN-CIRCULANT MATRIX . ON RETURN A IS UNALTERED .
M INTEGER
THE NUMBER OF ROWS OF THE MATRICES A AND Q .
M MUST bE AT LEAST AS LARGE AS L .
L INTEGER
THE NUMBER OF COLUMNS OF THE MATRICES A AND Q AND THE ORDER OF THE UPPER TRIANGULAR MATRIX S .
LDQ INTEGER
THE LEADING DIMENSION OF THE ARRAY Q .
LDS INTEGER
THE LEADING DIMENSION OF THE ARRAY S .
ON RETURN
Q COMPLEX (M,L)
THE Q MATRIX OF THE FACTORIZATION .
THE COLUMNS OF Q ARE ORTHONORMAL .
S COMPLEX(L,L)
THE INVERSE OF THE R MATRIX OF THE FACTORIZATION. ELEMENTS BELOW THE MAIN DIAGONAL ARE NOT ACCESSED .
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SUBROUTINES AND FUNCTIONS
```

```
LINPACK ... CAXPY,CDOTC,CSSCAL,SCNRM2
```

```
FORTRAN ... CONJG
INTERNAL VARIABLES
INTEGER I,J,J1,JI
REAL SCALE,SCNRM2
COMPLEX C,CDOTC
INITIALIZATION (LAST COLUMN OF Q USED AS WORK VECTOR) .
DO \(10 \mathrm{I}=1, \mathrm{M}\)
```

```
        Q(I,1) = A(I)
        Q(I,L) = A(I)
    10 CONTINUE
C
C RECURRENT PROCESS FOR THE LATTICE ALGORITHM WITH NORMALIZATION .
C
    DO 70 J1 = 1, L
        J = J1 + 1
        SCALE = 1.OE0/SCNRM2(M,Q(1,J1),1)
        IF (J1 .EQ. L) GO TO 60
            C = -SCALE*(CONJG(Q(M,J1))*Q(1,L) +
                CDOTC(M-1,Q(1,J1),1,Q(2,L),1))/SCNRM2(M,Q(1,L),1)
            Q(1,J) = Q(M,J1) + C*Q(1,L)
            DO 20 I = 2, M
                Q(I,J) = Q(I-1,J1) + C*Q(I,L)
            CONTINUE
            IF (J .EQ. L) GO TO 30
                Q(1,L) = Q(1,L) + C%Q(M,J1)
                    CALL CAXPY(M-1,C,Q(1,J1),1,Q(2,L),1)
            CONTINUE
            S(1,J) = C
            IF (J .EQ. 2) GO TO 50
                    DO 40 I = 2, J1
                    JI = J - I
                    S(I,J) = S(I-1,J1) + C*S(JI,J1)
                    CONTINUE
            CONTINUE
        CONTINUE
        CALL CSSCAL(M,SCALE,Q(1,J1),1)
        S(J1,J1) = (1.0E0,0.0E0)
        CALL CSSCAL(J1,SC.LEE,S(1,J1),1)
70 CONTINUE
    RETURN
    END
```

SUBROUTINE TGSLS (A, X,R,M,L,LDA)
INTEGER M,L,LDA
REAL A(LDA, 1$), X(M, L), R(1)$

TGSLS CALLS TGSLS1 TO SOLVE THE REAL LINEAR SYSTEM
$\mathrm{A} \div \mathrm{X}=\mathrm{B}$
WITH THE TG - MATRIX A .
ON ENTRY
A $\quad \operatorname{REAL}(M * * 2,2 * L-1)$
THE FIRST ROW OF BLOCKS OF THE TG - MATRIX
FOLLOWED BY ITS FIRST COLUMN OF BLOCKS BEGINNING WITH THE SECOND BLOCK. EACH BLOCK IS REPRESENTED BY COLUMNS. ON RETURN A IS UNALTERED .
$\mathrm{X} \quad \operatorname{REAL}(\mathrm{M} \div \mathrm{L})$
THE RIGHT HAND SIDE VECTOR B .
R $\quad \operatorname{RFAL}(M * * 2 *(2 * L+3)+M)$
A WORK VECTOR .
M INTEGER
THE ORDER OF THE BLOCKS OF THE MATRIX A.
L INTEGER
THE NUMBER OF BLOCKS IN A ROW OR COLUMN OF THE MATRIX A .

LDA INTEGER
THE LEADING DIMENSION OF THE ARRAY A .
ON RETURN
X THE SOLUTION VECTOR .
TOEPLITZ PACKAGE. THIS VERSION DATED 07/23/82.
SUBROUTINES AND FUNCTIONS
TOEPLITZ PACKAGE ... TGSLS1
INTERNAL VARIABLES
INTEGER MM,MML,MML1,MML2,MML3,MML4,MML5,MNE $\kappa$
CALL SUBROUTINE TGSLS 1
$M M=M \div * 2$
MML $=$ MM* $(\mathrm{L}-1)+1$
MML1 $=2 *$ MML -1
MML2 $=$ MML1 + MM
MML3 $=$ MML2 + MM
MML4 $=$ MML3 + MM

```
    MML5 = MML4 + MM
    MML6 = MML5 + MM
C
        CALL TGSLS1(A,A(1,L+1),X,X,R,R(MML),R(MML1),R(MML2),
        * R(MML3),R(MML4),R(MML5),R(MML6),M,L,LDA)
    C
        RETURN
        END
```

```
    SUBROUTINE TGSLS1(A1,A2,B,X,C1,C2,R1,R2,R3,R5,R6,R,M,L,LDA)
    INTEGER M,L,LDA
    REAL A1(LDA,L),A2(LDA,1), B(M,L),X(M,L),C1(M,M,1),
* C2(M,M,1),R1(M,M),R2(M,M),R3(M,M),R5(M,M),R6(M,M),R(M)
```

TGSLSI SOLVES THE REAL LINEAR SYSTEM
$A * X=B$
WITH THE TG - MATRIX A .
ON ENTRY
A1 REAL (M**2, L)
THE FIRST ROW OF BLOCKS OF THE TG - MATRIX A .
EACH BLOCK IS REPRESENTED BY COLUMNS .
ON RETURN AI IS UNALTERED .
A2 $\operatorname{REAL}(M \div * 2, \mathrm{~L}-1)$
THE FIRST COLUMN OF BLOCKS OF THE TG - MATRIX A
beginning with The second block. Each block is
REPRESENTED BY COLUMNS. ON RETURN A2 IS UNALTERED .
B $\quad \operatorname{REAL}(M * \mathrm{~L})$
THE RIGHT HAND SIDE VECTOR .
ON RETURN B IS UNALTERED .
C1 $\operatorname{REAL}(M, M, L-1)$
A WORK ARRAY .
C2 $\operatorname{REAL}(M, M, L-1)$
A WORK ARRAY .
R1 $\operatorname{REAL}(M, M)$
A WORK ARRAY .
R2 REAL(M,M)
A WORK ARRAY .
R3 $\operatorname{REAL}(M, M)$
A WORK ARRAY .
R5 REAL (M, M)
A WORK ARRAY .
R6 REAL(M,M)
A WORK ARRAY.
R REAL(M)
A WORK VECTOR .
M INTEGER
THE ORDER OF THE BLOCKS OF THE MATRIX A .
L INTEGER
THE NUMBER OF BLOCKS IN A ROW OR COLUMN
OF THE MATRIX A.

```
C
    I3 = 1
    DO 20 J = 1, M
    DO 10 I = 1,M
        C1(I,J,1) = A1(I3,1)
        R1(I,J) = A1(I3,1)
        R3(I,J) = R1(I,J)
        I3 = I3 + 1
10 CONTINUE
    X(J,1) = B (J,1)
20 CONTINUE
    CALL SGEFA(R3,M,M,R,II)
    CALL SGESL(R3,M,M,R,X(1,1),0)
    IF (L .EQ. 1) GO TO 420
    RECURRENT PROCESS FOR SOLVING THE SYSTEM
    WITH THE TG - MATRIX FOR N = 2, L .
    D0 410 N = 2, L
        COMPUTE MULTIPLES OF THE FIRST AND LAS'T BLOCK COIUMNS OF
        THE INVERSE OF THE PRINCIPAL MINOR OF ORDER M*N .
        N1 = N - 1
        N2 = N - 2
        I3 = 1
        DO 40 J = 1, M
        DO 30 I = 1,M
            R5(I,J) = A2 (I3,N1)
            R6(I,J) = A1(I3,N).
            I3 = I3 + 1
        CONTINUE
40 CONTINUE
```

```
    IF (N .EQ. 2) GO TO 100
    DO 60 J = 1,M
        DO 50 I = 1, M
            C1(I,J,N1) = R2(I,J)
        CONTINUE
        CONTINUE
        DO 90 I1 = 1, N2
        I2 = N - Il
        DO 80 J = 1, M
            I3 = 1
            DO 70 I = 1,M
C FOR IN-LINE SAXPY, ACTIVATE NEXT 5 LINES AND DEACTIVATE FOLLOWING 3 .
C DO 65 II = 1, M
C R5(II,J) = R5(II,J) + C1(I,J,I2)*A2(I3,I1)
    R6(II,J) = R6(II,J) + C2(I,J,I1)*A1(I3,I1+1)
    I3 = I3 + 1
        CONTINUE
                        CALL SAXPY(M,C1(I,J,I2),A2(I3,I1),1,R5(1,J),1)
                        CALL SAXPY(M,C2(I,J,I1),A1(I3,I1+1),1,R6(1,J),1)
                        I3 = I3 + M
                            CONTINUE
            7 0
            80
            90
    100
    110 CONTINUE
        CALL SGESL(R3,M,M,R,R2(1,J),0)
        CONTINUE
        DO 140 J = 1, M
            DO 130 I = 1, M
            R3(I,J) = R6(I,J)
            R6(I,J) = -C1(I,J,1)
            CONTINUE
        CONTINUE
        nO 160 J = 1, M
            DO 150 I = 1, M
C FOR IN-LINE SAXPY, ACTIVATE NEXT 3 LINES AND DEACTIVATE FOLLOWING 1 .
C DO 145 II = 1, M
C C1(II,J,1) = C1(II, J,1) + R2(I,J)*R3(II,I)
C 145
    1 5 0
    160
            CONTINUE
            CALL SAXPY(M,R2(I,J),R3(1,I),1,C1(1,J,1),1)
        CONTINUE
    CONTINUE
    CALL SGEFA(R6,M,M,R,II)
    DO 180 J = 1, M
        CALL SGESL(R6,M,M,R,R3(1,J),0)
        DO 170 I = 1,M
C FOR IN-LINE SAXPY, ACTIVATE NEXT 3 LINES AND DEACTIVATE FOLLOWING 1.
C DO 165 II = 1, M
C R1(II,J)=R1(II,J) + R3(I,J)*R5(II,I)
C }16
    CONTINUE
    CALL SAXPY(M,R3(I,J),R5(1,I),1,R1(1,J),1)
```

```
170
1 8 0
190
200
210
220
230
CONTINUE
    CONTINUE
    IF (N .EQ. 2) GO TO 320
    DO 200 J = 1, M
            DO 190 I = 1,M
                        R6(I,J) = C2(I,J,1)
            cONTINUE
    CONTINUE
    DO 310 I1 = 2, N1
        IF (I1 .EQ. N1) GO TO 230
            DO 220 J = 1, M
                DO 210 I = 1, M
                    R5(I,J) = C2(I,J,II)
                CONTINUE
            CONTINUE
        CONTINUE
        DO 260 J = 1, M
        DO 240 I = 1, M
                C2(I,J,I1) = R6(I,J)
            CONTINUE
            DO 250 I = 1, M
C FOR IN-LINE SAXPY, ACTIVATE NEXT 3 LINES AND DEACTIVATE FOLLOWING 1 .
C
C
C 245
250
260
                                    DO 245 II = 1, M
                                    C2(II,J,I1) = C2(II,J,I1) + R3(I,J)*C1(II,I,I1)
                                    CONTINUE
                            CALL SAXPY(M,R3(I,J),C1(1,I,I1),1,C2(1,J,I1),1)
                            CONTINUE
        conTINUE
        DO 280 J = 1, M
            DO 270 I = 1,M
C FOR IN-LINE SAXPY, ACTIVATE NEXT 3 LINES AND DEACTIVATE FOLLOWING 1 .
C
            DO 265 II = 1, M
            C1(II,J,I1) = C1(II,J,I1) + R2(I,J)*R6(II,I)
C 265
                    CONTINUE
                    CALL SAXPY(M,R2(I,J),R6(1,I),1,C1(1,J,I1),1)

\section*{CONTINUE}
```

DO $340 \mathrm{~J}=1$, M DO 330 I = 1, M
C2(I,J,1) = R3(I,J)
CONTINUE
330
340
C
C COMPUTE THE SOLUTION OF THE SYSTEM WITH THE
C PRINCIPAL MINOR OF ORDER M*N .
C
DO 360 J = 1, M

```
```

            DO 350 I = 1, M
            R3(I,J) = R1(I,J)
    350
        CUNTINUE
        X(J,N)=B(J,N)
        CONTINUE
        DO 380 I1 = 1, N1
        I2 = N - I1
        I3 = 1
        DO 370 I = 1, M
    C FOR IN-LINE SAXPY, ACTIVATE NEXT 4 LINES AND DEACTIVATE FOLLOWING 2 .
C DO 365 II = 1, M
C X(II,N)=X(II,N) - X(I,I2)*A2(I3,I1)
C I3 = I3 + 1
C 365
370
380
CALL SGEFA(R3,M,M,R,II)
CALL SGESL(R3,M,M,R,X(1,N),0)
DO 400 I1 = 1, N1
DO 390 I = 1, M
C FOR IN-LINE SAXPY, ACTIVATE NEXT 3 LINES AND DEACTIVATE FOLLOWING 1.
C DO 385 II = 1, M
C
X(II,I1) = X(II,I1) + X(I,N)*C2(II, I,I1)
C 385 CONTINUE
CALL SAXPY(M,X(I,N),C2(1,I,I1),1,X(1,I1),1)
CONTINUE
4 0 0 ~ C O N T I N U E ~
410 CONTINUE
420 CONTINUE
RETURN
END

```
```

SUBROUTJNE TGSLC(A,X,R,M,L,LDA)
INTEGER M,L,LDA
COMPLEX A(LDA, 1),X(M,L),R(1)

```
TGSLC CALLS TGSLCI TO SOLVE THE COMPLEX LINEAR SYSTEM
\(A * X=B\)
WITH THE TG - MATRIX A.

ON ENTRY

A COMPLEX (M* \(\because 2,2 \div L\) - 1)
THE FIRST ROW OF bLOCKS OF THE TG - MATRIX FOLLOWED BY ITS FIRST COLUMN OF BLOCKS BEGINNING WITH THE SELOND BLOCK. EACH BLOCK IS REPRESENTED BY COLUMNS. ON RETURN A IS UNALTERED .
\(X \quad\) COMPLEX (M*L)
THE RIGFT HAND SIDE VECTOR B .
\(\mathrm{R} \quad \operatorname{COMPLEX}(M * * 2 *(2 * L+3)+M)\) A WORK VECTOR .

M INTEGER
THE ORDER OF THE BLOCKS OF THE MATRIX A .
L INTEGER
THE NUMBER OF BLOCKS IN A ROW OR COLUMN OF THE MATRIX A.

LDA INTEGER
THE LEADING DIMENSION OF THE ARRAY A.

ON RETURN

X THE SOIUTION VECTOR .

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SUBROUTINES AND FUNCTIONS

TOEPLITZ PACKAGE ... TGSLC1

INTERNAL VARIABLES

INTEGER MM, MML, MML1,MML2,MML3,MML4, MML5, MML6
CALL SUBROUTINE TGSLC1
\(M M=M \% 2\)
\(M M L=M M^{*}(L-1)+1\)
MML1 \(=2\) MMML -1
MML2 \(=\) MML1 + MM
MML3 \(=\) MML2 + MM
MML4 \(=\) MML3 + MM
```

        MML5 \(=\) MML4 + MM
        MML6 \(=\) MML5 + MM
    C
        CALL TGSLC1 (A, A (1, L+1) , X,X,R,R(MML),R(MML1),R(MML2),
        * \(\quad\) (MML3) \(, R(M M L 4), R(M M L 5), R(M M L 6), M, L, L D A)\)
    C
        RETURN
        END
    ```
```

    SUBROUTINE TGSLC1(A1,A2,B,X,C1,C2,P1,R2,R3,R5,R6,R,M,L,LDA)
    INTEGER M,L,LDA
    COMPLEX A1(LDA,L),A2(LDA, 1), B(M,L),X(M,L),C1(M,M,1),
    * C2(M,M,1),R1(M,M),R2(M,M),R3(M,M),R5(M,M),R6(M,M),R(M)

```
TGSLC1 SOLVES THE COMPLEX LINEAR SYSTEM
A * \(\mathrm{X}=\mathrm{B}\)
WITh THE TG - MATRIX A .
ON ENTRY
A1 COMPLEX ( \(\mathrm{M} * \% 2, \mathrm{~L}\) )
    THE FIRST ROW OF BLOCKS OF THE TG - MATRIX A .
    EACH BLOCK IS REPRESENTED BY COLUMNS .
    ON RETURN A1 IS UNALTERED .
A2 \(\operatorname{COMPLEX}\left(\mathrm{M}^{*} * 2, \mathrm{~L}-1\right)\)
    THE FIRST COLUMN OF BLOCKS OF THE TG - MATRIX A
    BEGINNING WITH THE SECOND BLOCK. EACH BLOCK IS
    REPRESENTED BY COLUMNS. ON RETURN A2 IS UNALTERED .
B COMPLEX (M*L)
    THE RIGHT HAND SIDE VECTOR .
    ON RETURN B IS UNALTERED .
    C1 \(\operatorname{COMPLEX}(\mathrm{M}, \mathrm{M}, \mathrm{L}-1)\)
        A WORK ARRAY .
    C2 \(\operatorname{COMPLEX}(\mathrm{M}, \mathrm{M}, \mathrm{L}-1)\)
        A WORK ARRAY .
    R1 \(\operatorname{COMPLEX}(M, M)\)
        A WORK ARRAY .
    R2 COMPLEX (M,M)
    A WORK ARRAY .
    R3 COMPLEX(M,M)
        A WORK ARRAY .
    R5 COMPI.FX(M,M)
        A WORK ARRAY .
    R6 \(\operatorname{COMPLEX}(M, M)\)
    A WORK ARRAY .
    R COMPLEX (M)
        A WORK VECTOR .
    M INTEGER
    THE ORDER OF THE BLOCKS OF TIE MATRIX A .
    L INTEGER
        THE NUMBER OF BLOCKS IN A ROW OR COLUMN
        OF THE MATRIX A.

C
C C C C C C C C

LDA INTEGER THE LEADING DIMENSION OF THE ARRAY A .

ON RETURN
X COMPLEX (M*L)
THE SOLUTION VECTOR. X MAY COINCIDE WITH B .
TOEPLITZ PACKAGE. THIS VERSION DATED 07/23/82.
SUBROUTINES AND FUNCTIONS
LINPACK ... CAXPY,CGEFA,CGESL ... (FOR IN-LINE CAXPY, SEE DIRECTIONS IN COMMENTS)

INTERNAL VARIABLES
INTEGER I, I1,I2, I3, II , J,N,N1,N2
SOLVE THE SYSTEM WITH THE PRINCIPAL MINOR OF ORDER M .
I3 \(=1\)
DO \(20 \mathrm{~J}=1, \mathrm{M}\)
DO \(10 \mathrm{I}=1, \mathrm{M}\)
\(\mathrm{C} 1(\mathrm{I}, \mathrm{J}, 1)=\mathrm{A} 1(13,1)\)
\(\operatorname{R1}(I, J)=A 1(13,1)\)
R3(I, J) \(=\) R1 (I, J)
\(13=13+1\)
10
CONTINUE
\(X(J, 1)=B(J, 1)\)
20 CONTINUE
CALL CGEFA (R3,M,M,R,II)
CALL CGESL (R3, M, M, R,X \((1,1), 0)\)
IF (L .EQ. 1) GO TO 420
RECURRENT PROCESS FOR SOLVING THE SYSTEM
WITH THE TG - MATRIX FOR \(N=2\), L .
DO \(410 \mathrm{~N}=2\), L
COMPUTE MULTIPLES OF THE FIRST AND LAST BLOCK COLUMNS OF THE INVERSE OF THE PRINCIPAL MINOR OF ORDER \(M * N\).
\(\mathrm{N} 1=\mathrm{N}-1\)
\(\mathrm{N} 2=\mathrm{N}\) - 2
I3 \(=1\)
DO \(40 \mathrm{~J}=1, \mathrm{M}\)
DO 30 I = 1, M
R5 \((\mathrm{I}, \mathrm{J})=\mathrm{A} 2(\mathrm{I} 3, \mathrm{~N} 1)\)
R6(I, J) \(=\) A1 \((13, N)\)
\(13=13+1\)
CONTINUE
```

            IF (N .EQ. 2) GO TO 100
            DO 60 J = 1, M
        DO 50 I = 1, M
            C1(I,J,N1) = R2(I, J)
        CONTINUE
        CONTINUF,
        DO 90 I1 = 1, N2
            I2 = N - I1
            DO 80 J = 1, M
            I3 = 1
            DO 70 I = 1, M
    C FOR IN-LINE CAXPY, ACTIVATE NEXT 5 LINES AND DEACTIVATE FOLLOWING 3 .
C
C R5(II,J) = R5(II,J) + C1(I,J,I2)*A2(I3,I1)
C R6(II,J) = R6(II,J) + C2(I,J,I1)*A1(I3,I1+1)
C I3 = I3 + 1
C }6

```
DO 65 II = 1, M
                                    CONTINUE
                                    CALL CAXPY(M, C1 (I,J,I2),A2(I3,I1),1,R5(1,J),1)
                                    CALL CAXPY(M,C2(I,J,I1),A1(I3,I1+1),1,R6(1,J),1)
                    I3 = I3 + M
                    CONTINUE
                        CONTINUE
            CONTINUE
    CONTINUE
    DO 120 J = 1, M
            DO 110 I = 1,M
            R2(I,J) = -R5(I,J)
    110 CONTINUE
            CALL CGESL(R3,M,M,R,R2(1,J),0)
    120
    130
    140
        DO 140 J = 1, M
            DO 130 I = 1, M
            R3(I,J) = R6(I, J)
            R6(I,J) = -C1(I,J,1)
            CONTINUE
            CONTINUE
        DO 160 J = 1, M
            DO 150 I = 1, M
C FOR IN-LINE CAXPY, ACTIVATE NEXT 3 LINES AND DEACTIVATE FOLLOWING 1 .
C
C
C }14
    150
    160 CONTINUE
        DO 145 II = 1, M
            C1(II,J,1)=C1(II,J,1) + R2(I,J)`R3(II,I)
            CONTINUE
            CALL CAXPY(M,R2(I , J),R3(1,I),1,C1(1,J,1),1)
            CONTINUE
        CALL CGEFA(R6,M,M,R,II)
        DO 180 J = 1, M
            CALL CGESL(R6,M,M,R,R3(1,J),0)
            DO 170 I = 1, M
C FOR IN-LINE CAXPY, ACTIVATE NEXT 3 LINES AND DEACTIVATE FOLLOWING 1.
C DO 165 II = 1,M
C
C }16
        R1(II,J)=R1(II, J) + R3(I,J)*R5(II,I)
        CONTINUE
        CALL CAXPY(M,R3(I ,J),R5(1,I),1,R1(1,J),1)
```

CONTINUE
CONTINUE
IF (N .EQ. 2) GO TO 320
DO $200 \mathrm{~J}=1$, M
DO $190 \mathrm{I}=1$, M
$\mathrm{R} 6(\mathrm{I}, \mathrm{J})=\mathrm{C} 2(\mathrm{I}, \mathrm{J}, 1)$
CONTINUE
CONTINUE
DO $310 \mathrm{I} 1=2$, N 1
IF (I1 .EQ. N1) GO TO 230
DO $220 \mathrm{~J}=1$, M
DO $210 \mathrm{I}=1$, M
$R 5(I, J)=C 2(I, J, I 1)$
CONTINUE
CONTINUE
CONTINUE
DO $260 \mathrm{~J}=1, \mathrm{M}$
DO $240 \mathrm{I}=1, \mathrm{M}$ $\mathrm{C} 2(\mathrm{I}, \mathrm{J}, \mathrm{I} 1)=\mathrm{R} 6(\mathrm{I}, \mathrm{J})$
CONTINUE
DO $250 \mathrm{I}=1$, M
C FOR IN-LINE CAXPY, ACTIVATE NEXT 3 LINES AND DEACTIVATE FOLLOWING 1 .
C DO 245 II $=1, \mathrm{M}$
C
C 245

250
260

$$
\mathrm{C} 2(\mathrm{II}, \mathrm{~J}, \mathrm{I} 1)=\mathrm{C} 2(\mathrm{II}, \mathrm{~J}, \mathrm{I} 1)+\mathrm{R} 3(\mathrm{I}, \mathrm{~J}) * \mathrm{C} 1(\mathrm{II}, \mathrm{I}, \mathrm{I} 1)
$$ CONTINUE

CALL CAXPY(M,R3(I,J),C1(1,I,I1),1,C2(1,J,I1),1)
CONTINUE
CONTINUE
DO $280 \mathrm{~J}=1, \mathrm{M}$
DO $270 \mathrm{I}=1$, M
C FOR IN-LINE CAXPY, ACTIVATE NEXT 3 LINES AND DEACTIVATE FOLLOWING 1 .
C
DO 265 II $=1, M$
$\mathrm{C} 1(\mathrm{II}, \mathrm{J}, \mathrm{I} 1)=\mathrm{C} 1(\mathrm{II}, \mathrm{J}, \mathrm{I} 1)+\mathrm{R} 2(\mathrm{I}, \mathrm{J}) * \mathrm{R} 6(\mathrm{II}, \mathrm{I})$
C 265
CONTINUE
CALL CAXPY(M,R2(I,J),R6(1,I),1,C1(1,J,I1),1)
CONTINUE
CONTINUE
DO $300 \mathrm{~J}=1$, M
DO $290 \mathrm{I}=1$, M
$\operatorname{R6}(\mathrm{I}, \mathrm{J})=\mathrm{R} 5(\mathrm{I}, \mathrm{J})$
CONTINUE
CONTINUE
CONTINUE
CONTINUE
DO $340 \mathrm{~J}=1, \mathrm{M}$
DO $330 \mathrm{I}=1, \mathrm{M}$
$\mathrm{C} 2(\mathrm{I}, \mathrm{J}, 1)=\mathrm{R} 3(\mathrm{I}, \mathrm{J})$
CONTINUE
CONTINUE
C
C
C
COMPUTE THE SOLUTION OF THE SYSTEM WITH THE
PRINCIPAL MINOR OF ORDER M ${ }^{*} N$.

DO $360 \mathrm{~J}=1, \mathrm{M}$

```
        DO 350 I = 1, M
    R3(I,J) = R1(I,J)
        CONTINUE
        X(J,N) = B(J,N)
    CONTINUE
        DO 380 I1 = 1, N1
        I2 = N - I1
        I3 = 1
        DO 370 I = 1, M
C FOR IN-LINE CAXPY, ACTIVATE NEXT 4 LINES AND DEACTIVATE FOLLOWING 2 .
C DO 365 II = 1, M
C X(II,N) = X(II,N) - X(I,I2)*A2(I3,I1)
    I3 = I3 + 1
C, 365
            CONTINUE
            CALL CAXPY(M, -X(I,I2),A2(I3,I1),1,X(1,N),1)
            I3 = I3 + M
            CONTINUE
            CONTINUE
            CALL. CGEFA(R3,M,M,R,II)
            CALL CGESL(R3,M,M,R,X(1,N),0)
            DO 400 I1 = 1, N1
            D0 390 I = 1, M
C FOR IN-LINE CAXPY, ACTIVATE NEXT 3 LINES AND DEACTIVATE FOLLOWING 1 .
C DO 385 II = 1, M
C X(II,I1) = X(II,I1) + X(I,N)*C2(II,I,I1)
C 385
    390
    4 0 0
    410 CONTINUE
    420 CONTINUE
    RETURN
    END
```

SUBROUTINE CTSLC(A, X, R, M, L, LDA)
INTEGER M, L, LDA
COMPLEX A(LDA, $\left.\mathrm{I}_{\mathrm{I}}\right), \mathrm{X}(\mathrm{M}, \mathrm{L}), \mathrm{R}(1)$

CTSLC SOLVES THE COMPLEX LINEAR SYSTEM
$\mathrm{A} * \mathrm{X}=\mathrm{B}$
WITH THE C'T - MATRIX A .
ON ENTRY
A $\quad \operatorname{OMPLEX}(2 \div M-1, L)$ THE FIRST ROW OF BLOCKS OF THE CT - MATRIX . EACH BLOCK IS REPRESENTED BY ITS FIRST ROW FOLLOWED BY ITS FIRST COLUMN BEGINNING WITH THE SECOND ELEMENT. ON RETURN A HAS BEEN DESTROYED .
$\mathrm{X} \quad \operatorname{COMPLEX}(\mathrm{M} * \mathrm{~L})$ THE RIGHT HAND SIDE VECTOR B .
$R \quad \operatorname{COMPLEX}(\operatorname{MAX}(2 * M-2,2 * L))$ A WORK VECTOR .

M INTEGER
THE ORDER OF THE BLOCKS OF THE MATRIX A.
L INTEGER
THE NUMBER OF BLOCKS IN A ROW OR COLUMN OF THE MATRIX A.

LDA INTEGER
THE LEADING DIMENSION OF THE ARRAY A .
ON RETURN
X THE SOLUTION VECTOR .
TOEPLITZ PACKAGE. THIS VERSION DATED 07/23/82.
SUBROUTINES AND FUNCTIONS

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TOEPLITZ PACKAGE ... SALWC,TSLC
```

FORTRAN ... FLOAT
INTERNAL VARIABLES
INTEGER I1,I2
REAL RL
RL $=$ FLOAT(L)
REDUCE THE CT - MATRIX TO A BLOCK-DIAGONAL MATRIX
by THE INVERSE DISCRETE FOURIER TRANSFORMATION.
CALL SALWC(A,R,R(L+1),2*M - 1,L,LDA,-1)

C COMPUTE THE DISCRETE FOURIER TRANSFORMATION OF C THE RIGHT HAND SIDE VECTOR .

C
CALL $\operatorname{SALWC}(X, R, R(L+1), M, L, M, 1)$
C
C SOLVE THE BLOCK-DIAGONAL SYSTEM, BLOCKS OF WHICH C ARE T - MATRICES .

C
DO $10 \mathrm{I} 2=1, \mathrm{~L}$ CALI. $\operatorname{TSLC}(A(1, I 2), X(1, I 2), R, M)$
10 CONTINUE
C
C COMPUTE THE SOLUTION OF THE GIVEN SYSTEM BY
C THE INVERSE DISCRETE FOURIER TRANSFORMATION .
C
CALI SALWC $(X, R, R(L+1), M, L, M,-1)$
C
DO $30 \mathrm{I} 2=1, \mathrm{~L}$ DO 20 I1 = 1, M $X(I 1, I 2)=X(I 1, I 2) / R L$
20 CONTINUE
30 CONTINUE
RETURN
END

SUBROUTINE CCSLC (A, X, R, M, L, LDA $)$
INTEGER M,L,LDA
COMPLEX A(LDA,L), X(M,L),R(1)

CCSLC SOLVES THE COMPLEX LINEAR SYSTEM
$A * X=B$
WITH THE CC - MATRIX A .
ON ENTRY
A COMPLEX (M,L)
THE FIRST ROW OF BLOCKS OF THE CC - MATRIX .
EACH BLOCK IS REPRESENTED BY ITS FIRST ROW . ON RETURN A YAS BEEN DESTROYED .
$\mathrm{X} \quad \operatorname{COMPLEX}(\mathrm{M} * \mathrm{~L})$
THE RIGHT HAND SIDE VECTOR B .
R $\operatorname{COMPLEX}(\operatorname{MAX}(M, 2 * L))$
A WORK VECTOR .
INTEGER
THE ORDER OF THE BLOCKS OF THE MATRIX A .
L INTEGER
THE NUMBER OF BLOCKS IN A ROW OR COLUMN OF THE MATRIX A .

LDA INTEGER
THE LEADING DIMENSION OF THE ARRAY A .
ON RETUKN
X THE SOLUTION VECTOR .
TOEPLITZ PACKAGE. THIS VERSION DATED 07/23/82.
SUBROUTINES AND FUNCTIONS

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TOEPLITZ PACKAGE ... CSLC,SALWC
```

    FORTRAN ... FLOAT
    INTERNAL VARIABLES
INTEGER I1,I2
REAL RL
RL $=$ FLOAT( L )
REDUCE THE CC - MATRIX TO A BLOCK-DIAGONAL MATRIX
BY THE INVERSE DISCRETE FOURIER TRANSFORMATION .
CALL SALWC (A,R,R(L+1),M,L,LDA,-1)
COMPUTE THE DISCRETE FOURIER TRANSFORMATION OF

C THE RIGHT HAND SIDE VECTOR .
C
CALL SALWC(X,R,R(L+1),M,L,M,1)
C
C SOLVE THE BLOCK-DIAGONAL SYSTEM, BLOCKS OF WHICH
C ARE C - MATRICES .
C
DO $10 \mathrm{I} 2=1$, L
CALL $\operatorname{CSLC}(A(1, I 2), X(1, I 2), R, M)$
10 CONTINUE
C
C COMPUTE THE SOLUTION OF THE GIVEN SYSTEM BY
C THE INVERSE DISCRETE FOURIER TRANSFORMATION .
C
CALL $\operatorname{SALWC}(X, R, R(L+1), M, L, M,-1)$
C
DO 30 I2 $=1$, $L$
DO $20 \mathrm{I} 1=1, \mathrm{M}$
$X(I 1, I 2)=X(I 1, I 2) / R L$
20 CONTINUE
30 CONTINUE
RETURN
END

```
SUBROUTINE CGSLC(A,X,R,M,L,LDA)
INTEGER M,L,LDA
COMPLEX A(LDA,L) ,X(M,L),R(1)
```

CGSLC SOLVES THE COMPLEX LINEAR SYSTEM
$A * X=B$
WITH THE CG - MATRIX A .
ON ENTRY
A $\quad \operatorname{COMPLEX}(M \div \div \cdot 2, L)$
THE FIRST ROW OF BLOCKS OF THE CG - MATRIX .
EACH BLOCK IS REPRESENTED BY COLUMNS .
ON RETURN A HAS BEEN DESTROYED .
$\mathrm{X} \quad$ COMPLEX (M:L)
THE RIGHT HAND SIDE VECTOR B .
$R \quad \omega \operatorname{MPLEX}(\operatorname{MAX}(M, 2 * L))$
A WORK VECTOR .
M INTEGER
THE ORDER OF THE BLOCKS OF THE MATRIX A .
L INTEGER
THE NUMBER OF BLOCKS IN A ROW OR COLUMN
OF THE MATRIX A .
LDA INTEGER
THE LEADING DIMENSION OF THE ARRAY A.
ON RETURN
X THE SOLUTION VFCTOR .
TOEPLITZ PACKAGE. THIS VERSION DATED 07/23/82 .
SUBROUTINES AND FUNCTIONS
TOEPLITZ PACKAGE ... SALWC
LINPACK ... CGEFA,CGESL
FORTRAN ... FLOA $\Gamma$

INTERNAL VARIABLES
INTEGER I1,I2,II
REAL RL
$\mathrm{RL}=\mathrm{FLOAT}(\mathrm{L})$
REDUCE THE CG - MATRIX TO A BLOCK-DIAGONAL MATRIX BY THE INVERSE DISCRETE FOURIER TRANSFORMATION .

CALL SALWC (A,R,R(L+1), M**2,L,LDA,-1)

C COMPUTE THE DISCRETE FOURIER TRANSFORMATION OF
C THE RIGHT HAND SIDE VECTOR .
C
CALL SALWC(X,R,R(L+1),M,L,M,1)
C
C SOLVE THE BLOCK-DIAGONAL SYSTEM, BLOCKS OF WHICH
C ARE G - MATRICES .
C
DO $10 \mathrm{I} 2=1$,
CALL CGEFA(A(1,I2), M,M,R,II)
CALL CGESL(A(1, 12), M, M,R,X(1,I2), 0)
10 CONTINUE
C
C COMPUTE THE SOLUTION OF THE GIVEN SYSTEM BY
C THE INVERSE DISCRETE FOURIER TRANSFORMATION
C
CALL $\operatorname{SALWC}(X, R, R(L+1), M, L, M,-1)$
C
Do $30 \mathrm{I} 2=1$, L
DO $20 \mathrm{I} 1=1, \mathrm{M}$
$X(\mathrm{I} 1, \mathrm{I} 2)=\mathrm{X}(\mathrm{I} 1, \mathrm{I} 2) / \mathrm{RL}$
20 CONTINUE
30 CONTINUE
RETURN
END

```
SURROUTINE SALWC(A,R1,R2,M,L,LDA,JOB)
INTEGER M, L, LDA, JOB
COMPLEX &(LLA, L),R1(L),R2(L)
IF (L .EQ. 1) GO TO 60
R1(1) = (1.OE0,0.0E0)
RI = 0.0EO
DO 10 I1 = 2, L
    RI = RI + 1.0EO
```

```
RL = FIMOAT(L)
```

RL = FIMOAT(L)
RL $=$ FI,OAT(L)

```
C

ON ENTRY
A COMPLEX (M,L)
THE INPUT MATRIX .
R1 COMPLEX(L)
A WORK VECTOR .
R2 COMPLEX(L)
A WORK VECTOR .
M INTEGER

L INTEGER

INTEGER

JOB INTEGER

ON RETURN

SUBROUTINES AND FUNCTIONS

INTERNAL VARIABLES
INTEGER I,I1,I2
REAL P,Rj,RL,V1,V2
COMPLEX E,F
IF (L .EQ. 1) GO TO 60
\(R 1(1)=(1.0 E 0,0.0 \mathrm{E} 0)\)
RI \(=0.0 \mathrm{EO}\)
\(\mathrm{RI}=\mathrm{RI}+1.0 \mathrm{E} 0\)

SALWC COMPUTES THE DIRECT OR INVERSE DISCRETE FOURIER TRANSFORMATION FOR ROWS OF A COMPLEX RECTANGULAR MATRIX .

THE NUMBER OF ROWS OF THE MATRIX A .
the number of columins of the matrix a .

THE LEADING DIMENSION OF THE ARRAY A .
= 1 FOR DIRECT FOURIER TRANSFORMATION .
= -1 FOR INVERSE FOJRIER TRANSFORMATION.

A THE TRANSFORMED ROWS OF THE MATRIX .
TOEPLITZ PACKAGE. THIS VERSION DATED 07/23/82.

FORTRAN ... CMPLX,COS,FLOAT,SIN

C MINIMIZE ERROR IN FORMING MULTIPLES OF 2\%PI .
        \(\mathrm{P}=((201 . \mathrm{E} 0 / 32 . \mathrm{E} 0) * \mathrm{RI}+1.93530717958647692528 \mathrm{E}-3 * \mathrm{RI}) / \mathrm{RL}\)
C
        \(\mathrm{V} 1=\cos (\mathrm{P})\)
        \(\mathrm{V} 2=\operatorname{SIN}(\mathrm{P})\)
        IF (JOB .EQ. (-1)) V2 = -V2
        R1(I1) \(=\operatorname{CMPLX}(V 1, V 2)\)
10 CONTINUE
    DO \(50 \mathrm{I}=1, \mathrm{M}\)
        DO 30 I1 \(=1\), \(L\)
        \(\mathrm{E}=\mathrm{R} 1(\mathrm{II})\)
        \(\mathrm{F}=\mathrm{A}(\mathrm{I}, 1)\)
        DO \(20 \mathrm{I} 2=2\), L
                \(\mathrm{F}=\mathrm{E} * \mathrm{~F}+\mathrm{A}(\mathrm{I}, \mathrm{I} 2)\)
            CONTINUE
            R2(I1) \(=\mathrm{E} \star\) F
        CONTINUE
        DO 40 I1 = 1, L
            \(A(I, I 1)=R 2(I 1)\)
40 CONTINUE
50 CONTINUE
60 CONTINUE
    RETURN
    END
```

SUBROUTINE CTGSLC(A,X,R,M,L,K,LDA)
INTEGER M,L,K,LDA
COMPLEX A(LDA,K),X(M,L,K),R(1)
CTGSLC SOLVES THE COMPLEX LINEAR SYSTEM
A* X = B
WITH THE CTG - MATRIX A .
ON ENTRY

```
A \(\quad \operatorname{COMPLEX}\left(\mathrm{M}^{*} * 2 *(2 * \mathrm{~L}-1), \mathrm{K}\right)\)
THE FIRST ROW OF OUTER BLOCKS OF THE CTG - MATRIX .
EACH OUTER BLOCK IS REPRESENTED BY ITS FIRST ROW
OF INNER BLOCKS FOLLOWED BY ITS FIRST COLUMN
OF INNER BLOCKS BEGINNING WITH THE SECOND BLOCK .
EACH INNER BLOCK IS REPRESENTED BY COLUMNS .
ON RETURN A HAS BEEN DESTROYED .
\(\mathrm{X} \quad \operatorname{COMPLEX}(\mathrm{M} \div \mathrm{L} * \mathrm{~K})\)
THE RIGHT HAND SIDE VECTOR B .
\(R \quad \operatorname{COMPLEX}(\operatorname{MAX}(N * * 2 *(2 * L+3)+M, 2 * K))\)
A WORK VECTOR .
M INTEGER
THE ORDER OF THE INNER BLOCKS OF THE MATRIX A .
L INTEGER
    THE NUMBER OF INNER BLOCKS IN A ROW OR COLUMN
    OF AN OUTER BLOCK OF THE MATRIX A .
K INTEGER
    THE NUMBER OF OUTER BLOCKS IN A ROW OR COLUMN
    OF THE MATRIX A.
LDA INTEGER
THE LEADING DIMENSION OF THE ARRAY A.
ON RETURN
    X THE SOLUTION VECTOR .
    TOEPLITZ PACKAGE. THIS VERSION DATED 07/23/82 .
    SUBROUTINES AND FUNCTIONS
        TOEPLITZ PACKAGE ... SALWC,TGSLC
        FORTRAN ... FLOAT
INTERNAL VARIABLES
INTEGER I1,I2,I3,ML,MM
REAL RK
\[
\mathrm{RK}=\mathrm{FLOAT}(\mathrm{~K})
\]
\[
M M=M * 2
\]
\[
\mathrm{ML}=\mathrm{M} \div \mathrm{L}
\]

C REDUCE THE CTG - MATRIX TO A BLOCK-DIAGONAL MATRIX
C SOLVE THE BLOCK-DIAGONAL SYSTEM, BLOCKS OF WHICH
C ARE TG - MA'iRICES .

DO 10 I3 \(=1, \mathrm{~K}\) CALL TGSLC \((A(1, I 3), X(1,1, I 3), R, M, L, M M)\)
10 CONTINUE
COMPUTE THE SOLUTION OF THE GIVEN SYSTEM BY
THE INVERSE DISCRETE FOURIET TRANSFORMATION .
CALL SALWC(X,R,R(K+1),ML,K,ML,-1)
DO \(40 \mathrm{I} 3=1\), K
DO \(30 \mathrm{I} 2=1\), L
DO \(20 \mathrm{II}=1, \mathrm{M}\)
\(\mathrm{X}(\mathrm{I} 1, \mathrm{I} 2, \mathrm{I} 3)=\mathrm{X}(\mathrm{I} 1, \mathrm{I} 2, \mathrm{I} 3) / \mathrm{RK}\)
CONTINUE
30 CONTINUE
40 CONTINUE
RETURN
END
```

SUBROUTINE CCTSLC(A,X,R,M,L,K,LDA)
INTEGER M,L,K,LDA
COMPLEX A(LDA,K),X(M,L,K),R(1)
A : X = B
WITH THE CCT - MATRIX A .
ON ENTRY
A COMPLEX((2*M - 1)*L,K)
THE FIRST ROW OF OUTER BLOCKS OF THE CCT - MATRIX .
EACH OUTER BLOCK IS REPRESENTED BY ITS FIRST ROW
OF INNER BLOCKS. EACH INNER BLOCK IS REPRESENTED
BY ITS FIRST ROW FOLLOWED BY ITS FIRST COLUMN
BEGINNING WITH THE SECOND ELEMENT .
ON RETURN A HAS BEEN LESTROYED .
X COMPLEX(M*L*K)
THF RIGHT HAND SIDE VECTOR B .
R COMPLEX(MAX(2*M - 2,2*L,2*K))
A WORK VECTOR .
M INTEGER
THE ORDER OF THE INNER BLOCKS OF THE MATRIX A .
L INTEGER
THE NJMBER OF INNER BLOCKS IN A ROW OR COLUMN
OF AN OUTER BLOCK OF THE MATRIX A.
K INTEGER
THE NUMBER OF OUTER BLOCKS IN A ROW OR COLUMN
OF THE MATRIX A .
LDA INTEGER
THE LEADING DIMENSION OF THE ARRAY A .
ON RETURN
X THE SOLUTION VECTOR
TOEPLITZ PACKAGE. THIS VERSION DATED 07/23/82 .
SUBROUTINES AND FUNCTIONS
TOEPLITZ PACKAGE ... CTSLC,SALWC
FORTRAN ... FLOAT
INTERNAL VARIABLES
INTEGER I1,I2,I3,M2,ML
REAL RK

```

> RK \(=\) FLOAT \((K)\)
> M2 \(=2 \div M-1\)
\(\mathrm{ML}=\mathrm{M} * \mathrm{~L}\)
C
C
C
C
C
C COMPUTE THE DISCRETE FOURIER TRANSFORMATION OF
C

C SOLVE THE BLOCK-DIAGONAL SYSTEM, BLOCKS OF WHICH
C
DO \(10 \mathrm{I} 3=1\), K
CALL CTSLC(A(1, I3) \(\mathrm{X}(1,1, \mathrm{I} 3), \mathrm{R}, \mathrm{M}, \mathrm{L}, \mathrm{M} 2)\)
10 CONTINUE
    THE INVERSE DISCRETE FOURIER TRANSFORMATION .

CALL SALWC (X,R,R(K+1),ML,K,ML,-1)
DO \(40 \mathrm{I} 3=1\), K DO \(30 \mathrm{I} 2=1\), L DO 20 I1 = 1, M \(X(\mathrm{I} 1, \mathrm{I} 2, \mathrm{I} 3)=\mathrm{X}(\mathrm{I} 1, \mathrm{I} 2, \mathrm{I} 3) / \mathrm{RK}\) CONTINUE CONTINUE

\section*{40 CONTINUE}

RETURN
END

SUBROUTINE CCCSLC(A, X,R,M,L,K,LDA)
INTEGER M,L,K,LDA
COMPLEX A(LDA,K),X(M,L,K),R(1)
```

RK = FLOAT(K)

```
\(M L=M \star L\)

C REDUCE THE CCC - MATRIX TO A BLOCK-DIAGONAL MATRIX C BY THE INVERSE DISCRETE FOURIER TRANSFORMATION .

CALL SALWC(A,R,R(K+1),ML,K,LDA,-1)
C
C COMPUTE THE DISCRETE FOURIER TRANSFORMATION OF
C THE RIGHT HAND SIDE VECTOR .
C
CALL SALWC (X,R,R(K+1),ML,K,ML,1)
C
C SOLVE THE BLOCK-DIAGONAL SYSTEM, BLOCKS OF WHICH
C ARE CC - MATRICES .
C
DO 10 I3 = 1, K
\(\operatorname{CALL} \operatorname{CCSLC}(A(1, I 3), X(1,1, I 3), R, M, L, M)\)
10 CONTINUE
C
C COMPUTE THE SOLUTION OF THE GIVEN SYSTEM BY
C THE INVERSE DISCRETE FOURIER TRANSFORMATION
C
CALL SALWC(X,R,R(K+1),ML,K,ML,-1)
C
D0 \(40 \mathrm{I} 3=1\), K
DO \(30 \mathrm{I} 2=1\), L
DO \(20 \mathrm{II}=1, \mathrm{M}\)
\(\mathrm{X}(\mathrm{I} 1, \mathrm{I} 2, \mathrm{I} 3)=\mathrm{X}(\mathrm{I} 1, \mathrm{I} 2, \mathrm{I} 3) / \mathrm{RK}\)
CONTINUE
20 CONTIN
30 CONTINUE
40 CONTINUE
RETURN
END
```

SUBROUTINE CCGSLC(A,X,R,M,L,K,LDA)
INTEGER M,l,K,LDA
COMPLEX A(LDA,K),X(M,L,K),R(1)
CCGSLC SOLVES THE COMPLEX LINEAR SYSTEM
$A * X=B$
WITH THE CCG - MATRIX A .
ON ENTRY
A COMPLEX (M**2*L, K)
THE FIRST ROW OF OUTER BLOCKS OF THE CCG - MATRIX .
EACH OUTER BLOCK IS REPRESENTED BY ITS FIRST ROW
OF INNER BLOCKS. EACH INNER BLOCK IS REPRESENTED
by COLUMNS. ON RETURN A has been destroyed .
X COMPIEX ( $M * L * K$ )
THE RIGHT HAND SIDE VECTOR B .
R $\operatorname{COMPLEX}(\operatorname{MAX}(M, 2 * L, 2 * K))$
A WORK VECTOR .
M INTEGER
THE ORDER OF THE INNER BLOCKS OF THE MATRIX A .
L INTEGER
THE NUMBER OF INNER BLOCKS IN A ROW OR COLUMN OF AN OUTER BLOCK OF THE MATRIX A .
K INTEGER
THE NUMBER OF OUTER BLOCKS IN A ROW OR COLUNN OF THE MATRIX A.
LDA INTEGER
THE LEADING DIMENSION OF THE ARRAY A .
ON RETURN
X THE SOLUTION VECTOR .
TOEPLITZ PACKAGE. THIS VERSION DATED 07/23/82.
SUBROUTINES AND FUNCTIONS
TOEPLITZ PACKAGE ... CGSLLC,SALWC
rORTRAN ... FLOAT
INTERNAL VARIABLES
INTEGER I1,I2,I3,ML,MM
REAL RK
$\mathrm{RK}=\mathrm{FLOAT}(\mathrm{K})$
$M M=M * * 2$

```
\(\mathrm{ML}=\mathrm{M} * \mathrm{~L}\)

C
C
C C
        DO 10 I3 \(=1\), \(K\)
        CALL \(\operatorname{CGSLC}(A(1,13), X(1,1, I 3), R, M, L, M M)\)
    10 CONTINUE
C
C
    REDUCE THE CCG - MATRIX TO A BLOCK-DIAGONAL MATRIX
    bY THE INVERSE DISCRETE FOLRIER TRANSFORMATION .
    CALL SALWC (A,R,R(K+1),MM*L,K,LDA,-1)
    COMPUTE THE DISCRETE FOURIER TRANSFORMATION OF
    THE RIGHT HAND SIDE VECTOR .
    CALL SALWC (X,R,R(K+1),ML,K,ML,1)
SOLVE THE BLOCK-DIAGONAL SYSTEM, BLOCKS OF WIIICH
    ARE CG - MATRICES .
    COMPUTE THE SOLUTION OF THE GIVEN SYSTEM BY
    THE INVERSE DISCRETE FOURIER TRANSFORMATION .
    CALL SALWC (X,R,R(K+1),ML,K,ML,-1)
    DO \(40 \mathrm{I} 3=1\), K
    DO 30 I2 = 1, L
        DO 20 I1 \(=1, M\)
            \(\mathrm{X}(\mathrm{I} 1, \mathrm{I} 2, \mathrm{I} 3)=\mathrm{X}(\mathrm{I} 1, \mathrm{I} 2, \mathrm{I} 3) / \mathrm{RK}\)
        CONTINUE
    CONTINUE
CONTINUE
    RETURN
    END

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[^0]:    *The one member not governed by the naming convention is the service. subroutine SALWC (SALWZ in double precision), called by most of the two-level and all of the three-level system solving subroutines.

