

Traffic Steering in Software Defined Networks: Planning and Online Routing

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Outline

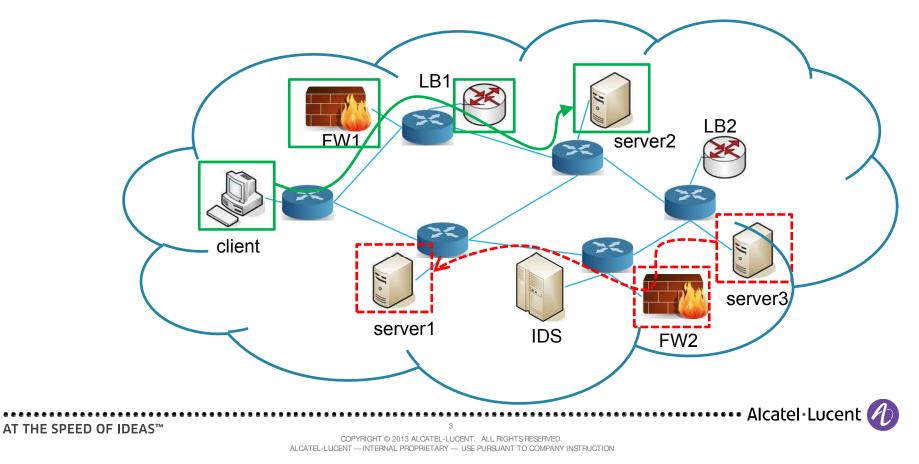
- 1. Introduction to Policy Aware Routing and SDN Framework
- 2. Problem Statement
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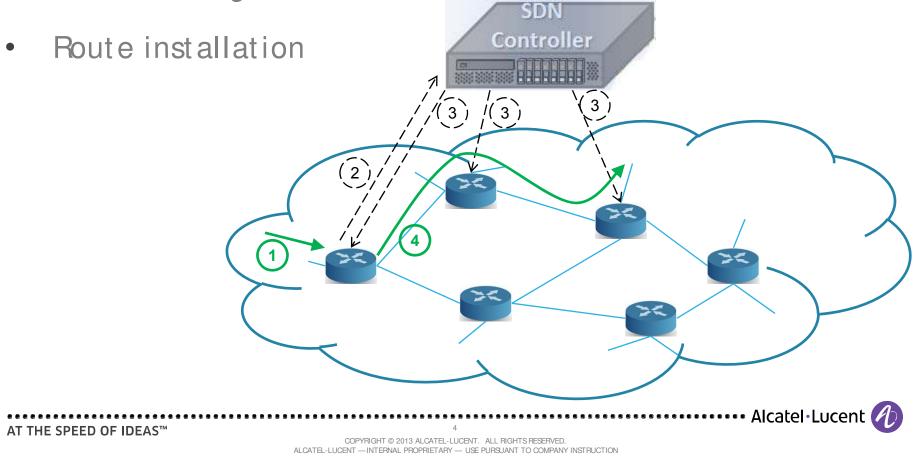
Policy Aware Routing – Middlebox Traversal

- Firewall, load balancer, intrusion and detection system
- Data center and wide area networks



Software Defined Network – Flow Scheduling

- Policy lookup
- Flow steering



Problem Statement

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- 1. Offline network planning problem: Given the location of the middleboxes and projected traffic, determine if there is enough capacity in the network to access the middlebox functions.
 FPTAS that achieves
 64 (1) approximate solution
- 2. Online traffic steering problem: When a flow has to be instantiated in the network, the SDN controller finds a physical path for the current request that permits as many future flows to be instantiated in the network as possible.



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General Methodology Description

Linear programming (exponential)

> $\max \lambda$ s.t. $\sum_{p \in \mathcal{P}_d} x_{dp} \ge \lambda h_d, \ \forall d$ $\sum_{d} \sum_{p:e \in p} x_{dp} \le c_e, \ \forall e$ $x_{dp} \ge 0, \ \forall d, \forall p$

 $\min \alpha(\mathcal{L}) = \sum_{e} c_e l_e$ s.t. $\sum_{e \in p} l_e \geq z_d, \ \forall d, \forall p \in \mathcal{P}_d$ $\sum_{d} h_d z_d \ge 1$ $l_e, z_d \ge 0, \ \forall e, \forall d$

augment shortest path lengths z_d

Primal-dual FPTAS (iterative)

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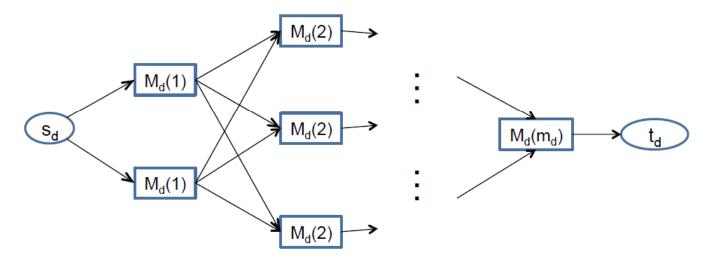
- Find the shortest path p for demand d under the current length system I_e ;
- Allocate flow $\Delta f_d := \min \{h'_d, \min_{e \in p_d(\mathcal{L})} \{c_e\}\}$.
- Augment the length of arcs e on path p by $l_e := l_e (1 + \epsilon \Delta f_d / c_e)$.

intuition: highly congested arcs get larger lengths to avoid being further exploited •••••••••••• Alcatel Lucent 🥢

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SPAR: Sizing Policy Aware Routed Networks

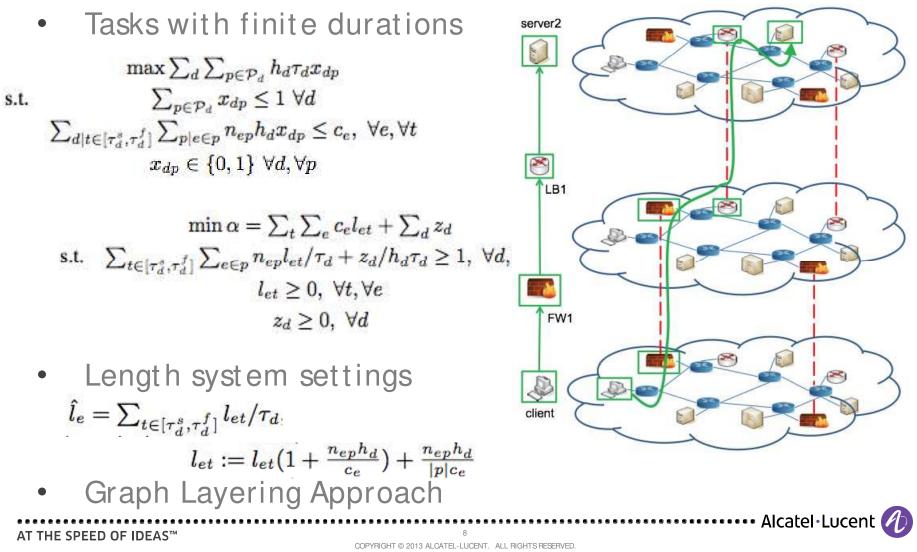
• Policy aware shortest path (segment+APSP+assembly)



- Lazy dual update
 - Perform one pass of APSP for all demands and segments;
 - Delay update until all policy aware shortest paths are found and allocated.



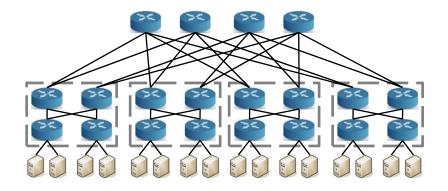
COATS: Competitive Online Algorithm for Traffic Steering



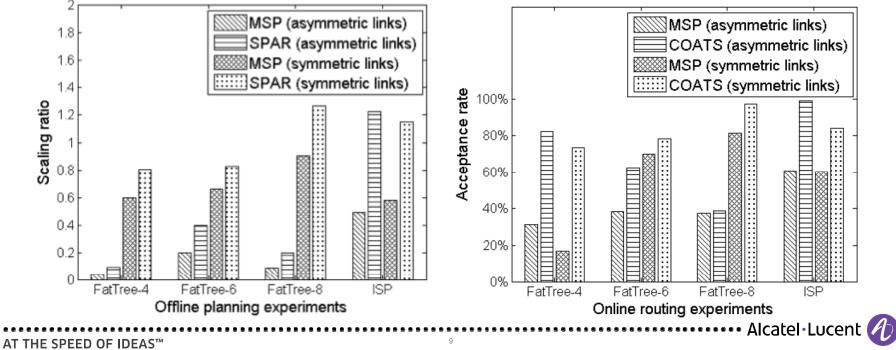
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Simulation Results

• Fat Tree + ISP networks



• SPAR & COATS



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Conclusions

- We have developed a FPTAS, SPAR, to solve the offline networking problem for policy aware routing. Lazy dual update is of independent interest and can be used for other constrained routing problems.
- We have also developed a log-competitive algorithm, COATS, to solve the online traffic steering problem. Tasks with finite durations are taken into account, as in real systems.
- Simulation results in both data center and ISP networks show that our proposed algorithms work very well.
- Future work: further exploit the challenges in policy aware routing and how they affect network design and management.

Thank you!

