that the best 8 -ary constellation is either the 8 -PSK, or 8 -cross, or $(1,7)$ depending on the desired spectral efficiency. As for the case of 16 -ary constellations, the most attractive signal set is rectangular 16-QAM for all spectral efficiencies of practical interest. These results have been confirmed by computer simulations of BICM schemes combining turbo codes and several signal sets. The study carried out in this correspondence does not formally prove the optimality of these particular signal sets for BICM design. To the best of our knowledge, we have, however, considered in this work the 11 signal sets that present the best SER performance and/or the lowest $N_{\text {min }}$, among all possible 8 - and 16 -ary constellations. Therefore, it is our belief that there is no constellation that significantly outperforms 8 -PSK, 8 -cross, $(1,7)$, or 16-QAM when combined with an error-correcting code in a BICM system.

An important conclusion of this work is that signal sets having a simple structure, such as 8-PSK and 16-QAM, are very attractive for BICM. This result strongly suggests that some other higher order constellations such as cross 32-QAM, 64-QAM, cross 128-QAM, 256-QAM, etc., are also very much of interest for the design of power-efficient BICM schemes. On the other hand, signal sets displaying optimal error performance in the absence of coding are, generally, not of interest for BICM design.

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# Two Generalized Complex Orthogonal Space-Time Block Codes of Rates 7/11 and 3/5 for 5 and 6 Transmit Antennas 

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#### Abstract

Space-time block codes from orthogonal designs have two advantages, namely, fast maximum-likelihood (ML) decoding and full diversity. Rate 1 real (pulse amplitude modulation-PAM) space-time codes (real orthogonal designs) for multiple transmit antennas have been constructed from the real Hurwitz-Radon families, which also provides the rate $1 / 2$ complex (quadrature amplitude modulation-QAM) space-time codes (complex orthogonal designs) for any number of transmit antennas.


[^0]Rate 3 / 4 complex orthogonal designs (space-time codes) for three and four transmit antennas have existed in the literature but no high rate ( $>\mathbf{1 / 2}$ ) complex orthogonal designs for other numbers of transmit antennas exist. In this correspondence, we present rate $7 / 11$ and rate $3 / 5$ generalized complex orthogonal designs for five and six transmit antennas, respectively.
Index Terms-Diversity, (generalized) complex orthogonal designs, space-time block codes.

## I. Introduction

Space-time coding for multiple transmit antenna systems in broad-band wireless communications has attracted considerable attention lately, see for example [1], [2]-[11]. To design "good" space-time block codes is a challenging problem. Since a space-time block code is a collection of some matrices, even for a small block size and a reasonable rate, the set of a space-time block code can be significantly large and therefore, its maximum-likelihood (ML) decoding may have a high complexity. On the other hand, the performance of a space-time block code depends on the diversity of the code. Therefore, a "good" space-time block code should possess two properties: i) the decoding at the receiver is reasonably fast; and ii) the diversity of the code is not small. Based on these two properties, space-time block codes based on orthogonal designs have been first proposed in Alamouti [3] for two transmit antennas and then generalized in Tarokh, Jafarkhani, and Calderbank [4] for $n \geq 2$ transmit antennas by connecting the space-time codes to orthogonal designs and the Hurwitz-Radon theory, see, for example, [12], [13]. A similar scheme was suggested in Ganesan and Stoica [6] from a maximum signal-to-noise ratio (SNR) approach.

Consider a space-time block code from complex orthogonal design $G$ for $n$ transmit antennas and block length $p$, i.e., $G$ is a $p \times n$ generator. Consider a signal constellation $\mathcal{S}$, for example the binary phase shift keying (BPSK) $\{1,-1\}$ or the quaternary phase shift keying (QPSK) $\{1,-1, j,-j\}$. Let $x_{1}, x_{2}, \ldots, x_{k}$ be information symbols in $\mathcal{S}$. The entries of the complex orthogonal design of $G$ are formed from complex linear combinations of $x_{1}, x_{1}^{*}, x_{2}, x_{2}^{*}, \ldots, x_{k}, x_{k}^{*}$ such that the columns of $G$ are orthogonal to each other, where $x^{*}$ is the complex conjugate of $x$. More detailed definition is reviewed later. See [4] for more about its encoding scheme and fast decoding algorithm. The rate of the code is $R=k / p$, which means that each codeword with block length $p$ carries $k$ information symbols. For a fixed $n$ and rate $R$, it is desired to have the block length $p$ as small as possible for decreasing the time delay in decoding. It was shown in [2], [4] that $R \leq 1$, i.e., $p \geq k$. Clearly, for a given $k$, the smallest possible block length $p$ is $k$. When the signal constellation $\mathcal{S}$ has all real symbols, such as pulse amplitude modulation (PAM), rate $R=1$, i.e., $p=k$, real orthogonal designs for any fixed transmit antenna number $n$ have been given in, for example, [4] from the real Hurwitz-Radon families [12]. These real orthogonal designs also provide a method to construct complex orthogonal designs of rate $R=k /(2 k)=1 / 2$ for any fixed transmit antenna number $n$ [4] where the block length $p=2 k$ and $k$ depends on $n$ (see details later). Examples of rate $3 / 4$ complex orthogonal designs for $n=3$ and $n=4$ transmit antennas have appeared in [4], [6]-[8]. It was proved in [4] that $4 \times 4$ complex orthogonal designs of rate 1 do not exist and a simpler proof was given in [6] by using the amicable design theory [12], [13]. This implies that the only square complex orthogonal design of rate 1 is the $2 \times 2$ complex orthogonal design proposed in [3]. To the best of our knowledge, there are no known high-rate ( $R>1 / 2$ ) complex orthogonal designs for transmit antenna number $n \geq 5$.

The main contribution of this correspondence is to present rate $7 / 11$ and rate $3 / 5$ generalized complex orthogonal designs for $n=5$ and $n=6$ transmit antennas, respectively. We also present some constructions of rate $1 / 2$ complex orthogonal designs with smaller block length $p$ than those presented in [4].

## II. (Generalized) Complex Orthogonal Designs

This correspondence follows the terminologies of [4]. We first review the concept of the (generalized) complex orthogonal design and some known designs of rate greater than or equal to $1 / 2$ [3], [4], [6]-[8]. We then present two generalized complex orthogonal designs of rate $7 / 11$ and rate $3 / 5$ for $n=5$ and $n=6$ transmit antennas, respectively. We also present some simpler complex orthogonal designs of rate $1 / 2$ for $n=5,6,7,8$ transmit antennas.

Definition 1: A generalized complex orthogonal design (GCOD) in variables $x_{1}, x_{2}, \ldots, x_{k}$ is a $p \times n$ matrix $G$ such that:

- the entries of $G$ are complex linear combinations of $x_{1}, x_{2}, \ldots, x_{k}$ and their complex conjugates $x_{1}^{*}, x_{2}^{*}, \ldots, x_{k}^{*}$;
- $G^{\mathcal{H}} G=D$, where $G^{\mathcal{H}}$ is the complex conjugate and transpose of $G$, and $D$ is an $n \times n$ diagonal matrix with the $(i, i)$ th diagonal element of the form

$$
l_{i, 1}\left|x_{1}\right|^{2}+l_{i, 2}\left|x_{2}\right|^{2}+\cdots+l_{i, k}\left|x_{k}\right|^{2}
$$

where all the coefficients $l_{i, 1}, l_{i, 2}, \ldots, l_{i, k}$ are strictly positive numbers.
The rate of $G$ is defined as $R=k / p$. If

$$
G^{\mathcal{H}} G=\left(\left|x_{1}\right|^{2}+\left|x_{2}\right|^{2}+\cdots+\left|x_{k}\right|^{2}\right) I_{n \times n}
$$

then $G$ is called a complex orthogonal design (COD).
Tarokh, Jafarkhani, and Calderbank [4] first mentioned that the rate of space-time block codes from generalized complex orthogonal designs cannot be greater than 1 , i.e., $R=k / p \leq 1$. Later, it was proved in [9] that this rate must be less than 1 for more than two transmit antennas. For a fixed number of transmit antennas $n$ and rate $R$, it is desired to have the block length $p$ as small as possible.

The first space-time block code from complex orthogonal design was proposed in Alamouti [3] for two transmit antennas. It is the following $2 \times 2$ COD in variables $x_{1}$ and $x_{2}$

$$
G_{2}=\left[\begin{array}{rr}
x_{1} & x_{2}  \tag{1}\\
-x_{2}^{*} & x_{1}^{*}
\end{array}\right]
$$

Clearly, the rate of $G_{2}$ achieves the maximum rate 1 . For space-time block codes from (generalized) complex orthogonal designs, rate 1 is achievable only for two transmit antennas.

For $n=3$ and $n=4$ transmit antennas, there are complex orthogonal designs of rate $R=3 / 4$ [4], [6]-[8], for example,

$$
G_{3}=\left[\begin{array}{rcr}
x_{1} & x_{2} & x_{3}  \tag{2}\\
-x_{2}^{*} & x_{1}^{*} & 0 \\
x_{3}^{*} & 0 & -x_{1}^{*} \\
0 & x_{3}^{*} & -x_{2}^{*}
\end{array}\right]
$$

for three transmit antennas, and

$$
G_{4}=\left[\begin{array}{rrrr}
x_{1} & x_{2} & x_{3} & 0  \tag{3}\\
-x_{2}^{*} & x_{1}^{*} & 0 & x_{3} \\
x_{3}^{*} & 0 & -x_{1}^{*} & x_{2} \\
0 & x_{3}^{*} & -x_{2}^{*} & -x_{1}
\end{array}\right]
$$

for four transmit antennas. In fact, $G_{3}$ is obtained by taking the first three columns of $G_{4}$. For $n \geq 5$, Tarokh, Jafarkhani, and Calderbank [4] gave a general construction for complex orthogonal design of rate $R=1 / 2$. The block length $p$ of this design is

$$
p=2 \min \left(2^{4 c+d}\right)
$$

where the minimization is taken over the set

$$
\left\{c, d: 0 \leq d<4, c \geq 0, \text { and } 8 c+2^{d} \geq n\right\} .
$$

We observe that, when $5 \leq n \leq 8, p$ is 16 .

By now, the existing designs of rate greater than $1 / 2$ are only $G_{2}, G_{3}$, and $G_{4}$ with rates $1,3 / 4$, and $3 / 4$, respectively. It has been proved in [4], [5] that if there exists a $p \times n$ GCOD in variables $x_{1}, x_{2}, \ldots, x_{k}$ such that $l_{i, 1}=l_{i, 2}=\cdots=l_{i, k}$ for each $i$, then there exists a COD in the same variables and of the same size. Notice that all of the existing designs are CODs. However, the constraint of CODs is not necessary to construct space-time block codes. GCODs can also provide the advantages of the fast ML decoding and the full diversity. In fact, the diagonal form of $D$ guarantees the fast ML decoding, since the orthogonal columns of $G$ can separate the transmitted symbols $x_{1}, x_{2}, \ldots, x_{k}$ from each other at the decoder. Also, the strictly positive coefficients $l_{i, 1}, l_{i, 2}, \ldots, l_{i, k}$ imply the full rank of $G$. This guarantees the full diversity advantage of coding. For more details about the coding scheme and the fast ML decoding, we refer the reader to [4].

## A. Two GCODs of Rates $7 / 11$ and $3 / 5$ for $n=5$ and $n=6$

## Transmit Antennas

We present here two generalized complex orthogonal designs of rate greater than $1 / 2$. The first one is an $11 \times 5$ matrix given by

$$
G_{5}=\left[\begin{array}{rrrrr}
x_{1} & x_{2} & x_{3} & 0 & x_{4}  \tag{4}\\
-x_{2}^{*} & x_{1}^{*} & 0 & x_{3} & x_{5} \\
x_{3}^{*} & 0 & -x_{1}^{*} & x_{2} & x_{6} \\
0 & x_{3}^{*} & -x_{2}^{*} & -x_{1} & x_{7} \\
x_{4}^{*} & 0 & 0 & -x_{7}^{*} & -x_{1}^{*} \\
0 & x_{4}^{*} & 0 & x_{6}^{*} & -x_{2}^{*} \\
0 & 0 & x_{4}^{*} & x_{5}^{*} & -x_{3}^{*} \\
0 & -x_{5}^{*} & x_{6}^{*} & 0 & x_{1} \\
x_{5}^{*} & 0 & x_{7}^{*} & 0 & x_{2} \\
-x_{6}^{*} & -x_{7}^{*} & 0 & 0 & x_{3} \\
x_{7} & -x_{6} & -x_{5} & x_{4} & 0
\end{array}\right]
$$

for $n=5$ transmit antennas. $G_{5}$ is constructed from $G_{4}$ as follows. At first, we keep $G_{4}$ as a $4 \times 4$ submatrix of $G_{5}$ and add symbols $x_{4}, x_{5}, x_{6}, x_{7}$ into the fifth column of $G_{5}$. Then, we arrange the entries of $G_{5}$ from the fifth row to the end such that all of the five columns are orthogonal to each other and the number of the total rows should be as small as possible. From the resulting matrix in (4), we can check that $G_{5}^{\mathcal{H}} G_{5}=D$, where $D$ is a $5 \times 5$ diagonal matrix with the $(i, i)$ th diagonal element $D(i, i)$ of the form

$$
D(1,1)=D(2,2)=D(3,3)=D(4,4)=\sum_{m=1}^{7}\left|x_{m}\right|^{2}
$$

and

$$
D(5,5)=2 \sum_{m=1}^{3}\left|x_{m}\right|^{2}+\sum_{m=4}^{7}\left|x_{m}\right|^{2}
$$

Notice that the symbols $x_{1}, x_{2}, x_{3}$ and their complex conjugates $x_{1}^{*}, x_{2}^{*}, x_{3}^{*}$ appear in the fifth column of $G_{5}$. Clearly, the rate of $G_{5}$ is $R=7 / 11=0.6364$, and the block length $(p=11)$ of $G_{5}$ is smaller than that ( $p=16$ ) of rate $1 / 2$ complex orthogonal design given in [4] for $n=5$ transmit antennas. A by-product of constructing $G_{5}$ is a $7 \times 4$ complex orthogonal design as follows:

$$
\left[\begin{array}{rrrr}
x_{1}^{*} & 0 & 0 & -x_{4}^{*}  \tag{5}\\
0 & x_{1}^{*} & 0 & x_{3}^{*} \\
0 & 0 & x_{1}^{*} & x_{2}^{*} \\
0 & -x_{2}^{*} & x_{3}^{*} & 0 \\
x_{2}^{*} & 0 & x_{4}^{*} & 0 \\
-x_{3}^{*} & -x_{4}^{*} & 0 & 0 \\
x_{4} & -x_{3} & -x_{2} & x_{1}
\end{array}\right]
$$

where rate $R=4 / 7=0.5714$, and the block length $p(=7)$ in this example is not a power of 2 as appeared in the previous examples of COD.

The second one is a $30 \times 6$ matrix given by

$$
G_{6}=\left[\begin{array}{cccccc}
x_{1} & x_{2} & x_{3} & 0 & x_{4} & x_{8}  \tag{6}\\
-x_{2}^{*} & x_{1}^{*} & 0 & x_{3} & x_{5} & x_{9} \\
x_{3}^{*} & 0 & -x_{1}^{*} & x_{2} & x_{6} & x_{10} \\
0 & x_{3}^{*} & -x_{2}^{*} & -x_{1} & x_{7} & x_{11} \\
x_{4}^{*} & 0 & 0 & -x_{7}^{*} & -x_{1}^{*} & x_{12} \\
0 & x_{4}^{*} & 0 & x_{6}^{*} & -x_{2}^{*} & x_{13} \\
0 & 0 & x_{4}^{*} & x_{5}^{*} & -x_{3}^{*} & x_{14} \\
0 & x_{5}^{*} & -x_{6}^{*} & 0 & -x_{1} & x_{15} \\
x_{5}^{*} & 0 & x_{7}^{*} & 0 & x_{2} & x_{16} \\
x_{6}^{*} & x_{7}^{*} & 0 & 0 & -x_{3} & x_{17} \\
x_{7} & -x_{6} & -x_{5} & x_{4} & 0 & x_{18} \\
x_{8}^{*} & 0 & 0 & -x_{11}^{*} & -x_{15}^{*} & -x_{1}^{*} \\
0 & x_{8}^{*} & 0 & x_{10}^{*} & x_{16}^{*} & -x_{2}^{*} \\
0 & 0 & x_{8}^{*} & x_{9}^{*} & -x_{17}^{*} & -x_{3}^{*} \\
0 & 0 & 0 & x_{18}^{*} & x_{8}^{*} & -x_{4}^{*} \\
0 & 0 & -x_{18}^{*} & 0 & x_{9}^{*} & -x_{5}^{*} \\
0 & -x_{18}^{*} & 0 & 0 & x_{10}^{*} & -x_{6}^{*} \\
x_{18}^{*} & 0 & 0 & 0 & x_{11}^{*} & -x_{7}^{*} \\
0 & -x_{9}^{*} & x_{10}^{*} & 0 & x_{12}^{*} & x_{1} \\
x_{9}^{*} & 0 & x_{11}^{*} & 0 & x_{13}^{*} & x_{2} \\
-x_{10}^{*} & -x_{11}^{*} & 0 & 0 & x_{14}^{*} & x_{3} \\
-x_{12}^{*} & -x_{13}^{*} & -x_{14}^{*} & 0 & 0 & x_{4} \\
-x_{16}^{*} & -x_{15}^{*} & 0 & -x_{14}^{*} & 0 & x_{5} \\
-x_{17}^{*} & 0 & x_{15}^{*} & -x_{13}^{*} & 0 & x_{6} \\
0 & -x_{17}^{*} & -x_{16}^{*} & x_{12}^{*} & 0 & x_{7} \\
0 & x_{14} & -x_{13} & -x_{15} & x_{11} & 0 \\
x_{14} & 0 & -x_{12} & -x_{16} & x_{10} & 0 \\
-x_{13} & x_{12} & 0 & x_{17} & x_{9} & 0 \\
x_{15} & -x_{16} & x_{17} & 0 & x_{8} & 0 \\
-x_{11} & x_{10} & x_{9} & -x_{8} & x_{18} & 0
\end{array}\right]
$$

for $n=6$ transmit antennas. Actually, $G_{6}$ is constructed form $G_{5}$ as follows. At first, we keep $G_{5}$ as an $11 \times 5$ submatrix of $G_{6}$ and add symbols $x_{8}, x_{9}, \ldots, x_{18}$ into the sixth column of $G_{6}$. Then, we arrange the entries of $G_{6}$ from the twelfth row to the end such that all of the six columns of $G_{6}$ are orthogonal to each other and the number of the total rows should be as small as possible. The resulting matrix in (6) is of size $30 \times 6$. By a tedious check, we have $G_{6}^{\mathcal{H}} G_{6}=D$, where $D$ is a $6 \times 6$ diagonal matrix with the $(i, i)$ th diagonal element $D(i, i)$ of the form

$$
D(1,1)=D(2,2)=D(3,3)=D(4,4)=\sum_{m=1}^{18}\left|x_{m}\right|^{2}
$$

and

$$
\begin{aligned}
& D(5,5)=\sum_{m=1}^{18}\left|x_{m}\right|^{2}+\sum_{m=1}^{3}\left|x_{m}\right|^{2}+\sum_{m=8}^{11}\left|x_{m}\right|^{2} \\
& D(6,6)=2 \sum_{m=1}^{7}\left|x_{m}\right|^{2}+\sum_{m=8}^{18}\left|x_{m}\right|^{2} .
\end{aligned}
$$

Notice that the symbols $x_{1}, x_{2}, \ldots, x_{7}$ and their conjugates $x_{1}^{*}, x_{2}^{*}, \ldots, x_{7}^{*}$ appear in the sixth column of $G_{6}$. Clearly, the rate of $G_{6}$ is $R=18 / 30=0.6$.

The same procedure may be used to construct generalized complex orthogonal designs for other numbers of transmit antennas. However, it is hard to obtain other designs with rate greater than $1 / 2$. For example, $G_{6}$ may be used to construct $G_{7}$ for seven transmit antennas
as follow. We keep $G_{6}$ as a $30 \times 6$ submatrix of $G_{7}$ and add symbols $x_{19}, x_{20}, \ldots, x_{48}$ into the seventh column of $G_{7}$. However, it is hard to arrange the entries of $G_{7}$ from the thirty-first row to the end such that all of the seven columns are orthogonal to each other and the number of the total rows should be as small as possible. Notice that in this case, the symbols $x_{1}, x_{2}, \ldots, x_{18}$ and their complex conjugates $x_{1}^{*}, x_{2}^{*}, \ldots, x_{18}^{*}$ should appear in the seventh column of $G_{7}$. Therefore, the block length of $G_{7}$ will be at least $30+18+18=66$.

## B. Simpler CODs of Rate $1 / 2$ for $n=5,6,7,8$ Transmit Antennas

For a $p \times n$ complex orthogonal design for $n$ transmit antennas, the block length $p$ is related to the time delay in decoding. Clearly, for a fixed transmit antenna number $n$ and rate $R$, it is desired to have the block length $p$ as small as possible.

In [4], the constructions of rate $1 / 2$ complex orthogonal designs for $n \geq 5$ take advantage of the real orthogonal designs of full rate 1 which are available for any $n \geq 2$. For $n(n=5,6,7,8)$ transmit antennas, the complex orthogonal design in [4] is given by taking the first $n$ columns of the following $16 \times 8$ matrix:

$$
\left[\begin{array}{rrrrrrrr}
x_{1} & x_{2} & x_{3} & x_{4} & x_{5} & x_{6} & x_{7} & x_{8}  \tag{7}\\
-x_{2} & x_{1} & x_{4} & -x_{3} & x_{6} & -x_{5} & -x_{8} & x_{7} \\
-x_{3} & -x_{4} & x_{1} & x_{2} & x_{7} & x_{8} & -x_{5} & -x_{6} \\
-x_{4} & x_{3} & -x_{2} & x_{1} & x_{8} & -x_{7} & x_{6} & -x_{5} \\
-x_{5} & -x_{6} & -x_{7} & -x_{8} & x_{1} & x_{2} & x_{3} & x_{4} \\
-x_{6} & x_{5} & -x_{8} & x_{7} & -x_{2} & x_{1} & -x_{4} & x_{3} \\
-x_{7} & x_{8} & x_{5} & -x_{6} & -x_{3} & x_{4} & x_{1} & -x_{2} \\
-x_{8} & -x_{7} & x_{6} & x_{5} & -x_{4} & -x_{3} & x_{2} & x_{1} \\
x_{1}^{*} & x_{2}^{*} & x_{3}^{*} & x_{4}^{*} & x_{5}^{*} & x_{6}^{*} & x_{7}^{*} & x_{8}^{*} \\
-x_{2}^{*} & x_{1}^{*} & x_{4}^{*} & -x_{3}^{*} & x_{6}^{*} & -x_{5}^{*} & -x_{8}^{*} & x_{7}^{*} \\
-x_{3}^{*} & -x_{4}^{*} & x_{1}^{*} & x_{2}^{*} & x_{7}^{*} & x_{8}^{*} & -x_{5}^{*} & -x_{6}^{*} \\
-x_{4}^{*} & x_{3}^{*} & -x_{2}^{*} & x_{1}^{*} & x_{8}^{*} & -x_{7}^{*} & x_{6}^{*} & -x_{5}^{*} \\
-x_{4}^{*} & -x_{6}^{*} & -x_{7}^{*} & -x_{8}^{*} & x_{1}^{*} & x_{2}^{*} & x_{3}^{*} & x_{4}^{4} \\
-x_{*}^{*} & x_{4}^{*} & -x_{8}^{*} & x_{7}^{*} & -x_{2}^{*} & x_{1}^{*} & -x_{4}^{*} & x_{3}^{*} \\
-x_{7}^{*} & x_{8}^{*} & x_{5}^{*} & -x_{6}^{6} & -x_{3}^{*} & x_{4}^{*} & x_{1}^{*} & -x_{2}^{*} \\
-x_{8}^{*} & -x_{7}^{*} & x_{6}^{*} & x_{5}^{*} & -x_{4}^{*} & -x_{3}^{2} & x_{2}^{*} & x_{1}^{4}
\end{array}\right]
$$

where the block length $p=16$.
For $n(n=5,6,7,8)$ transmit antennas, we have a simpler design of rate $1 / 2$ by taking the first $n$ columns of the following $8 \times 8$ matrix:

$$
\left[\begin{array}{rcrrrrrr}
x_{1} & x_{2} & x_{3} & 0 & x_{4} & 0 & 0 & 0  \tag{8}\\
-x_{2}^{*} & x_{1}^{*} & 0 & x_{3} & 0 & x_{4} & 0 & 0 \\
x_{3}^{*} & 0 & -x_{1}^{*} & x_{2} & 0 & 0 & x_{4} & 0 \\
0 & x_{3}^{*} & -x_{2}^{*} & -x_{1} & 0 & 0 & 0 & x_{4} \\
x_{4}^{*} & 0 & 0 & 0 & -x_{1}^{*} & x_{2} & -x_{3} & 0 \\
0 & x_{4}^{*} & 0 & 0 & -x_{2}^{*} & -x_{1} & 0 & -x_{3} \\
0 & 0 & x_{4}^{*} & 0 & -x_{3}^{*} & 0 & x_{1} & x_{2} \\
0 & 0 & 0 & x_{4}^{*} & 0 & -x_{3}^{*} & -x_{2}^{*} & x_{1}^{*}
\end{array}\right]
$$

systematic construction of complex orthogonal designs of rate $1 / 2$ for any number of transmit antennas based on the Hurwitz-Radon theory. The previously known designs of rate larger than $1 / 2$ and less than 1 were given only for three and four transmit antennas with rate $3 / 4$. In this correspondence, we presented two generalized complex orthogonal designs of rate $7 / 11$ and rate $3 / 5$ for five and six transmit antennas, respectively. Although the construction method may be used for other numbers of transmit antennas, it is hard to obtain other designs with rate larger than $1 / 2$.

As pointed out in [4], what we have known about orthogonal designs is only a tip of the iceberg. We hope that the two designs we presented will stimulate future work. Recently, it was proved in [10] that for more than two transmit antennas, the rate of CODs cannot be greater than $3 / 4$; and the rate of GCODs cannot be greater than $4 / 5$. A tutorial on space-time block codes from (generalized) complex orthogonal designs can be found in [11].

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## III. Conclusion and Some Comments

Space-time block codes from (generalized) complex orthogonal designs have two advantages: fast ML decoding and full diversity. It has been proved in [9] that the rate of these space-time block codes must be less than 1 except for the case of wo transmit antennas. There is a


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