

# Unique reducibility of multiple blocking sets

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**Abstract.** A weighted  $t$ -fold  $(n - k)$ -blocking set  $B$  of  $\text{PG}(n, q)$  always contains a minimal weighted  $t$ -fold  $(n - k)$ -blocking set. We prove that, if  $|B| < (t + 1)q^{n-k} + \theta_{n-k-1}$ , then the minimal weighted  $t$ -fold  $(n - k)$ -blocking set contained in  $B$  is unique.

**Mathematics Subject Classification (2010).** 51E21, 51E20.

**Keywords.** Projective space, multiple blocking sets.

## 1 Introduction

A  $t$ -fold  $(n - k)$ -blocking set of  $\text{PG}(n, q)$  is a set of points which meets every  $k$ -dimensional subspace in at least  $t$  points. To exclude the trivial cases we will always suppose that  $0 < k < n$ . If the points of the set are not all different, so the set is a *multiset* of points, then it is called a *weighted  $t$ -fold  $(n - k)$ -blocking set*. A *weight function* of  $\text{PG}(n, q)$  is a mapping from the point set of  $\text{PG}(n, q)$  to the set of nonnegative integers. For a point  $P$  the integer  $w(P)$  is the *weight* of  $P$ . There is a natural correspondence between multisets and weight functions of  $\text{PG}(n, q)$ : let the weight of a point be the multiplicity of that point in the set. For a weight function  $w$ , the weight of a set  $M$  of points is by definition the sum of the weights of all its points, denoted by  $w(M)$ , and  $w(\text{PG}(n, q)) =: |w|$  can be called the *total weight* of  $w$ . The multiset associated to a weight function  $w$  is a  $t$ -fold  $(n - k)$ -blocking set if and only if the weight of every  $k$ -dimensional subspace is at least  $t$ . If this is the case, then we will call the weight function  $w$  a  *$t$ -fold  $(n - k)$ -blocking set* for short.

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\*The author was supported by OTKA Grant K81310.

If  $w$  is a  $t$ -fold  $(n - k)$ -blocking set, then a point  $P$  is called a *non-essential* point of  $w$ , if the weight of every  $k$ -subspace containing  $P$  is at least  $t + 1$  and  $w(P) \geq 1$ . Then the weight function  $w'$  defined by

$$w'(Q) = \begin{cases} w(Q) & \text{if } Q \neq P, \\ w(P) - 1 & \text{if } Q = P \end{cases}$$

is also a  $t$ -fold  $(n - k)$ -blocking set.

If  $w$  and  $w'$  are weight functions, and  $w'(P) \leq w(P)$  for all points  $P \in \text{PG}(n, q)$ , then we will say that  $w'$  is *contained* in  $w$ , and denote this by  $w' \leq w$ .

The  $t$ -fold  $(n - k)$ -blocking set  $w$  is said to be *minimal* if  $w' \equiv w$  for any  $t$ -fold  $(n - k)$ -blocking set  $w'$  contained in  $w$ .

A  $t$ -fold  $(n - k)$ -blocking set is not minimal if and only if it has non-essential points. If we start reducing the weight of the non-essential points one by one, always checking carefully that the resulting set/weight function is still a  $t$ -fold  $(n - k)$ -blocking set, then after some steps we will arrive at a minimal  $t$ -fold  $(n - k)$ -blocking set. It is a natural question to ask if there are conditions which guarantee the uniqueness of this minimal  $t$ -fold  $(n - k)$ -blocking set. Here, two weight functions  $w'$  and  $w''$  are considered to be different if there is a point  $P$ , such that  $w'(P) \neq w''(P)$ .

In [12] such a condition is given for non-weighted 1-fold 1-blocking sets of  $\text{PG}(2, q)$ .

**Result 1.1.** (Szőnyi, [12]) *A non-weighted 1-fold 1-blocking set of  $\text{PG}(2, q)$ , with size smaller than  $2q + 1$  contains a unique minimal 1-fold 1-blocking set.*

This result was recently generalized to non-weighted 1-fold  $(n - k)$ -blocking sets of  $\text{PG}(n, q)$  in [9].

**Result 1.2.** (Lavrauw, Storme and Van de Voorde, [9]) *A non-weighted 1-fold  $(n - k)$ -blocking set of  $\text{PG}(n, q)$ , with size smaller than  $2q^{n-k}$  contains a unique minimal 1-fold  $(n - k)$ -blocking set.*

Using the standard notation  $\theta_m = \frac{q^{m+1}-1}{q-1}$  for the number of points of an  $m$ -dimensional subspace of  $\text{PG}(n, q)$ , our result is the following.

**Theorem 1.3.** *A weighted  $t$ -fold  $(n - k)$ -blocking set of  $\text{PG}(n, q)$ , with total weight smaller than*

$$(t + 1)q^{n-k} + \theta_{n-k-1}$$

*contains a unique minimal weighted  $t$ -fold  $(n - k)$ -blocking set.*

Note that Theorem 1.3 is stronger than Result 1.2. Examples in the last section show that the bound is sharp if  $t = 1$ , or if  $k = n - 1$ .

## 2 $t$ -fold $(n - k)$ -blocking sets containing two minimal $t$ -fold $(n - k)$ -blocking sets

Let  $w$  be a  $t$ -fold  $(n - k)$ -blocking set. We will now define a new weight function  $s_w$  on the points of  $\text{PG}(n, q)$ . For a point  $P$  let  $s_w(P)$  be the largest integer for which the weight function  $w'$  defined by

$$w'(Q) = \begin{cases} w(Q) & \text{if } Q \neq P, \\ w(P) - s_w(P) & \text{if } Q = P \end{cases}$$

is also a  $t$ -fold  $(n - k)$ -blocking set. Then  $w(P) \geq s_w(P) \geq 0$ , so if  $w(P) = 0$ , then  $s_w(P) = 0$ . It is also clear that  $w$  is minimal if and only if  $s_w \equiv 0$ .

**Lemma 2.1.** *For a  $t$ -fold  $(n - k)$ -blocking set  $w$  and  $P \in \text{PG}(n, q)$  the following are true:*

- (a)  $s_w(P) = \min\{w(P), \min_{P \in \Pi_k} (w(\Pi_k) - t)\}$ , where  $\Pi_k$  runs along the  $k$ -dimensional subspaces containing  $P$ ;
- (b)  $s_w(P) = \max_{w' \leq w} \{w(P) - w'(P)\}$ , where  $w'$  runs along the  $t$ -fold  $(n - k)$ -blocking sets contained in  $w$ .

**Lemma 2.2.** *If  $w$  is a  $t$ -fold  $(n - k)$ -blocking set which contains two different minimal  $t$ -fold  $(n - k)$ -blocking sets, then there is a weight function  $v \leq w$  and a line  $l^*$  with the following properties:*

- (a)  $v(\Pi_k) \geq t$  for any  $k$ -subspace  $\Pi_k$  not containing  $l^*$ ;

- (b)  $v(\Pi_k) \geq t - 1$  for any  $k$ -subspace  $\Pi_k$  containing  $l^*$ ;
- (c) there is a  $k$ -subspace  $\Pi_k^*$  containing  $l^*$ , for which  $v(\Pi_k^*) = t - 1$ ;
- (d)  $|w| \geq |v| + 2$ .

*Proof.* Let  $w'$  and  $w''$  be two different minimal  $t$ -fold  $(n - k)$ -blocking sets contained in  $w$ . Then there is a point  $P^* \in \text{PG}(n, q)$ , such that  $w'(P^*) > w''(P^*)$ . Define  $\tilde{w}$  as follows:

$$\tilde{w}(Q) = \begin{cases} w(Q) & \text{if } Q \neq P^*, \\ w'(P^*) & \text{if } Q = P^*. \end{cases}$$

Then  $\tilde{w}$  is a  $t$ -fold  $(n - k)$ -blocking set,  $w', w'' \leq \tilde{w}$ , and Lemma 2.1(b) yields that  $s_{\tilde{w}}(P^*) \geq \tilde{w}(P^*) - w''(P^*) = w'(P^*) - w''(P^*) > 0$ . (\*)

As  $\tilde{w}$  contains the minimal  $t$ -fold  $(n - k)$ -blocking set  $w'$ , we can start reducing the weight of the points with  $\tilde{w}(P) > w'(P)$ , one at a time, until we arrive at  $w'$ . Formally, let  $\tilde{w} = w_1 \geq w_2 \geq \dots \geq w_m = w'$  be a sequence of  $t$ -fold  $(n - k)$ -blocking sets, such that for  $i \in \{1, 2, \dots, m - 1\}$  the  $t$ -fold  $(n - k)$ -blocking sets  $w_i$  and  $w_{i+1}$  only differ in one point  $P_i$ , and  $w_{i+1}(P_i) = w_i(P_i) - 1$ . Clearly  $P_i \neq P^*$ , and the points  $P_i$  are not necessarily all different. It is also clear that  $\tilde{w} \neq w'$ , because  $\tilde{w} = w'$  would mean that  $w''$  is contained in  $w'$ , which is a contradiction, so  $m \geq 2$  follows.

By Lemma 2.1(a),  $s_{w_{i+1}} \leq s_{w_i}$ , in fact, for any point  $Q$ , either  $s_{w_{i+1}}(Q) = s_{w_i}(Q)$ , or  $s_{w_{i+1}}(Q) = s_{w_i}(Q) - 1$ . For the point  $P^*$  we have  $s_{\tilde{w}}(P^*) > 0$  by (\*), and  $s_{w'}(P^*) = 0$  by the minimality of  $w'$ . So there will be an  $i \in \{1, 2, \dots, m - 1\}$  such that  $s_{w_i}(P^*) = 1$  and  $s_{w_{i+1}}(P^*) = 0$ . The weight functions  $w_i$  and  $w_{i+1}$  only differ in the point  $P_i$ . Then by Lemma 2.1(a) there is a  $k$ -space  $\Pi_k^*$  which contains  $P_i$  and  $P^*$ , and has weight  $w_i(\Pi_k^*) = t + 1$ . Also by Lemma 2.1(a) this yields  $s_{w_i}(P_i) \leq 1$ , and as  $w_{i+1}(P_i) = w_i(P_i) - 1$ , so  $P_i$  is a non-essential point of  $w_i$ , then  $s_{w_i}(P_i) = 1$  follows. Thus for any  $k$ -dimensional subspace  $\Pi_k$ , which contains  $P^*$  and/or  $P_i$  we have  $w_i(\Pi_k) \geq t + 1$ .

Let  $l^*$  be the line connecting  $P_i$  and  $P^*$ , and define  $v$  to be the following weight function:

$$v(Q) = \begin{cases} w_i(Q) & \text{if } Q \notin \{P^*, P_i\}, \\ w_i(Q) - 1 & \text{if } Q \in \{P^*, P_i\}. \end{cases}$$

Clearly  $|w| \geq |w_i| = |v| + 2$ , and  $v$  is a weight function contained in  $w$ . The weight of a  $k$ -subspace  $\Pi_k$  is  $w_{i-1}(\Pi_k) - |\Pi_k \cap \{P^*, P_i\}|$ . Thus,  $v$ ,  $l^*$  and  $\Pi_k^*$  satisfy the properties given in the lemma.  $\square$

### 3 $t$ -fold nuclei

If  $t = 1$ ,  $n = 2$ ,  $k = 1$ , then Lemma 2.2 yields that if  $w$  is a 1-fold 1-blocking set of  $\text{PG}(2, q)$  containing two different minimal 1-fold 1-blocking sets, then  $w$  contains a weight function  $v$ , which defines a blocking set of the affine plane  $\text{AG}(2, q) := \text{PG}(2, q) \setminus l^*$ . Thus  $|w(\text{PG}(2, q))| \geq s(q) + 2$ , where  $s(q)$  denotes the size of the smallest 1-blocking set of  $\text{AG}(2, q)$ . There are several independent proofs for  $s(q) = 2q - 1$ , from which Result 1.1 follows (see Jamison [8], Brouwer and Schrijver [5], Blokhuis [2], Szőnyi [12]).

In [2],  $s(q) = 2q - 1$  is proved as a corollary of a theorem on *nuclei* of point sets. Now we generalize the notion of *nucleus* to multisets/weight functions.

**Definition 3.1.** (1) Let  $S$  be a multiset of  $\text{PG}(n, q)$ . A point  $P \notin S$  will be called a  $t$ -fold *nucleus* of  $S$  if every line through  $P$  meets  $S$  in at least  $t$  points, counted with multiplicities.

(2) Let  $w$  be a weight function of  $\text{PG}(n, q)$ . A point  $P \in \text{PG}(n, q)$  with  $w(P) = 0$  will be called a  $t$ -fold *nucleus* of  $w$  if every line through  $P$  has weight at least  $t$ .

For  $S$  to have nuclei, clearly  $|S| \geq t\theta_{n-1}$  is needed. Let  $|S| = t\theta_{n-1} + r$ ,  $r \geq 0$ .

Note that for  $|S| = t\theta_{n-1} - r$ ,  $r \geq 0$ , a ‘symmetric’ version of the definition can be: a point  $P \notin S$  is a  $t$ -fold nucleus of  $S$ , if every line through  $P$  meets  $S$  in at most  $t$  points, counted with multiplicities.

The notion of *nucleus* was first introduced by Mazzocca for affine sets for  $n = 2$ ,  $t = 1$  and  $r = 0$ . Blokhuis extended the notion to  $r \geq 0$  in [2] and to  $t \geq 1$  in [3], and Sziklai generalized the definition for sets of the projective space  $\text{PG}(n, q)$  in [11]. (The ‘symmetric’ version was introduced in [7] and [11].)

Denote by  $N^t(S)$  the set of  $t$ -fold nuclei of  $S$ , and let  $p$  be the characteristic of the field  $\text{GF}(q)$ .

**Result 3.2.** (Sziklai, [11]) *Let  $S$  be a set of points in  $\text{PG}(n, q)$  with  $|S| = t\theta_{n-1} + r$ ,  $r \geq 0$ . Let  $H_\infty$  be a given hyperplane,  $|S \cap H_\infty| = m_\infty$ . Then*

$$|N^t(S) \setminus H_\infty| \leq (r + 1)(q - 1),$$

*provided that  $\binom{t\theta_{n-1} + r - m_\infty}{r+1} \not\equiv 0 \pmod{p}$ .*

Result 3.2 was proved in the case when  $m_\infty = 0$ ,  $n = 2$  by Blokhuis and Wilbrink ( $r = 0$ ,  $t = 1$ , see [4]) and by Blokhuis (for  $r \geq 0$ ,  $t = 1$ , see [2], and for  $r \geq 0$ ,  $t \geq 1$  see [3]). The ‘symmetric’ version was also settled by Sziklai in [11].

As Result 3.2 is not applicable when  $\binom{t\theta_{n-1} + r - m_\infty}{r+1} \equiv 0 \pmod{p}$ , to obtain an upper bound in this case, Ball presented the following theorem.

**Result 3.3.** (Ball, [1]) *Let  $S$  be a set of points in  $\text{PG}(n, q)$  with  $|S| = t\theta_{n-1} + r$ ,  $r \geq 0$ , and let  $H_\infty$  be a given hyperplane,  $|S \cap H_\infty| = m_\infty$ . Then*

$$|N^t(S) \setminus H_\infty| \leq (r + 1 + j)(q - 1),$$

*provided that the binomial coefficient*

$$\binom{t\theta_{n-1} + r - m_\infty}{r + 1 + j} \not\equiv 0 \pmod{p}$$

*for some  $j \geq 0$ .*

The proof of Result 3.2 and 3.3 can be easily copied for multisets/weight functions and we obtain the following lemma.

**Lemma 3.4.** *Let  $w$  be a weight function on  $\text{PG}(n, q)$  and  $H_\infty$  a given hyperplane with  $w(H_\infty) = m_\infty$ . Suppose that  $w(\text{PG}(n, q)) = t\theta_{n-1} + r$ , with  $r \geq 0$ . Then if*

$$\binom{t\theta_{n-1} + r - m_\infty}{r + 1 + j} \not\equiv 0 \pmod{p}$$

*for some  $j \geq 0$ , then the number of  $t$ -fold nuclei of  $w$  in  $\text{PG}(n, q) \setminus H_\infty$  is at most  $(r + 1 + j)(q - 1)$ .*

*Proof.* If the binomial coefficient is nonzero, then  $w(\text{PG}(n, q) \setminus H_\infty) > 0$ , so the number of  $t$ -fold nuclei in  $\text{PG}(n, q) \setminus H_\infty$  is at most  $q^n - 1$ . Thus the statement is trivially true for  $r + 1 \geq \theta_{n-1}$ , so from now on we will suppose  $r < \theta_{n-1} - 1$ .

Identify the points of  $\text{AG}(n, q) := \text{PG}(n, q) \setminus H_\infty$  with the elements of  $\text{GF}(q^n)$ , and the points of  $H_\infty$  with the  $\theta_{n-1}$ -st roots of unity of  $\text{GF}(q^n)$  in the usual way. The points of  $\text{PG}(n, q)$  will be denoted by capital letters, and the corresponding elements of  $\text{GF}(q^n)$  by the same lowercase letters. Then for points  $A \neq B \in \text{AG}(n, q)$ , the line  $AB$  contains the ideal point  $C \in H_\infty$  if and only if  $(a - b)^{q-1} = c$  holds.

Let  $\mathcal{S} = \{a_1, a_2, \dots, a_{t\theta_{n-1}+r-m_\infty}\} \cup \{c_1, \dots, c_{m_\infty}\}$  be the multiset of elements of  $\text{GF}(q^n)$  corresponding to the points of nonzero weight of  $\text{PG}(n, q) \setminus H_\infty$  and  $H_\infty$  respectively, such that  $a \in \mathcal{S}$  has multiplicity  $w(A)$  in  $\mathcal{S}$  for the corresponding point  $A \in \text{PG}(n, q)$ .

Let  $X$  and  $Y$  be variables, and define

$$\mathcal{B}(X) = \{(X - a_i)^{q-1} \mid i = 1, \dots, t\theta_{n-1} + r - m_\infty\} \cup \{c_1, \dots, c_{m_\infty}\},$$

and

$$F(Y, X) = \prod_{b \in \mathcal{B}(X)} (Y - b).$$

Then

$$F(Y, X) = \sum_{j=0}^{t\theta_{n-1}+r} (-1)^j \sigma_j(\mathcal{B}(X)) Y^{t\theta_{n-1}+r-j},$$

where  $\sigma_j(\mathcal{B}(X))$  denotes the  $j$ th elementary symmetric polynomial of the set  $\mathcal{B}(X)$ .

Suppose that  $x \in \text{GF}(q^n)$  is an element corresponding to a  $t$ -fold nucleus of  $w$ . Then  $\mathcal{B}(x)$  contains every  $\theta_{n-1}$ -st root of unity with multiplicity at least  $t$ , so

$$F(Y, x) = (Y^{\theta_{n-1}} - 1)^t (Y^r + \text{terms of lower degree}).$$

As  $r < \theta_{n-1} - 1$ , the coefficients of the terms

$$Y^{t\theta_{n-1}-1}, Y^{t\theta_{n-1}-2}, \dots, Y^{(t-1)\theta_{n-1}+r+1}$$

are 0 in  $F(Y, x)$ . Thus  $\sigma_{r+1+j}(\mathcal{B}(x)) = 0$  for  $0 \leq j \leq \theta_{n-1} - r - 2$ .

The degree of  $\sigma_{r+1+j}(\mathcal{B}(X))$  as a polynomial of  $X$  is at most  $(r+1+j)(q-1)$ , with equality precisely if the binomial coefficient

$$\binom{t\theta_{n-1} + r - m_\infty}{r+1+j}$$

does not vanish. In this case  $\sigma_{r+1+j}(\mathcal{B}(X))$  is not the zero polynomial, and every nucleus is a root of it, hence the number of nuclei is at most its degree:  $(r+1+j)(q-1)$ .  $\square$

We will now use Lemma 3.4 for  $n = 2$ ,  $j = 0$  and  $m_\infty = t - 1$ .

**Lemma 3.5.** *Suppose that  $v$  is a weight function of  $\text{PG}(2, q)$  such that there is a line  $l_\infty$ , with  $v(l_\infty) = t - 1$ , while all other lines have weight at least  $t$ . Then  $|v| \geq (t+1)q - 1$ .*

*Proof.* Assume first that  $t \leq q - 2$ . Suppose on the contrary that  $v$  is such a weight function, yet the total weight of  $v$  is less than  $(t+1)q - 1$ . We may suppose  $|v| = (t+1)q - 2$  (or else increase the weight of some of the points of  $\text{PG}(2, q) \setminus l_\infty$ ). All lines other than  $l_\infty$  have weight at least  $t$ , which means that all the points of  $\text{PG}(2, q) \setminus l_\infty$  with weight 0 are  $t$ -fold nuclei of  $v$ . As  $v(\text{PG}(2, q) \setminus l_\infty) = (t+1)q - 2 - (t-1) = tq + q - t - 1$ ,  $\text{PG}(2, q) \setminus l_\infty$  has at most  $tq + q - t - 1$  points with positive  $v$  weight (and exactly this many if every point of  $\text{PG}(2, q) \setminus l_\infty$  has weight  $\leq 1$ ). So  $v$  has at least  $q^2 - (tq + q - t - 1) = q^2 - tq - q + t + 1$   $t$ -fold nuclei.

We will use Lemma 3.4 to prove that this is not possible. As

$$|v| = (t+1)q - 2 = t(q+1) + q - t - 2$$

and

$$\binom{t(q+1) + q - t - 2 - (t-1)}{q - t - 2 + 1} = \binom{tq + q - t - 1}{q - t - 1} \not\equiv 0 \pmod{p}$$

by Lucas' theorem, so Lemma 3.4 yields that the number of  $t$ -fold nuclei of  $v$  is at most  $(q-t-1)(q-1) = q^2 - tq - 2q + t + 1$ , a contradiction. The same arguments prove that, if  $|v| = (t+1)q - 1$ , then  $v(P) \leq 1$  for all points  $P \in \text{PG}(2, q) \setminus l_\infty$ .

For  $t \geq q - 1$ , the assertion can be proved by summing the weights of all lines through a carefully selected point  $P$ . If  $P \in \text{PG}(2, q) \setminus l_\infty$  and  $v(P) = 0$ ,



then  $|v| \geq t(q+1) = tq + t \geq tq + q - 1$ . If  $P \in l_\infty$  and  $v(P) = 0$ , then  $|v| \geq tq + t - 1$  and so if  $t \geq q$ , then we are done. If  $t = q - 1$  and all points of  $\text{PG}(2, q) \setminus l_\infty$  have positive weight, then  $v(\text{PG}(2, q) \setminus l_\infty) \geq q^2$ , so  $|v| \geq q^2 + t - 1 > (t+1)q - 1$ . With this we have proved that if we can select a point  $P \in \text{PG}(2, q)$  with  $v(P) = 0$ , then the assertion is true.

Assume now that  $v(P) > 0$  for every point, let  $m = \min_P v(P)$  and define a new weight function  $\tilde{v}$ , by  $\tilde{v}(P) := v(P) - m$ . Then  $\tilde{v}(l_\infty) = t - m(q+1) - 1$  and  $\tilde{v}(l) \geq t - m(q+1)$  for any line  $l \neq l_\infty$ . If  $t - m(q+1) \leq q - 2$  then we can use the first part of the proof to prove  $|\tilde{v}| \geq (t - m(q+1) + 1)q - 1$ . If  $t - m(q+1) \geq q - 1$  then we can use the second part, as there will be a point with zero  $\tilde{v}$  weight. Then

$$|v| = |\tilde{v}| + m(q^2 + q + 1) \geq (t - m(q+1) + 1)q - 1 + m(q^2 + q + 1) = (t+1)q - 1 + m.$$

Hence the result is established.  $\square$

## 4 Proof of the main theorem

**Theorem 1.3.** *A weighted  $t$ -fold  $(n - k)$ -blocking set of  $\text{PG}(n, q)$ , with total weight smaller than*

$$(t + 1)q^{n-k} + \theta_{n-k-1}$$

*contains a unique minimal weighted  $t$ -fold  $(n - k)$ -blocking set.*

*Proof.* Assume that  $w$  is a weighted  $t$ -fold  $(n - k)$ -blocking set of  $\text{PG}(n, q)$  which contains two different minimal  $t$ -fold  $(n - k)$ -blocking sets. We will prove  $|w| \geq (t + 1)q^{n-k} + \theta_{n-k-1}$ . By Lemma 2.2 there is a weight function  $v \leq w$ , a line  $l^*$  and a  $k$ -subspace  $\Pi_k^*$  containing  $l^*$ , such that

- (a)  $v(\Pi_k) \geq t$ , for every  $k$ -subspace  $\Pi_k$  not containing  $l^*$ ;
- (b)  $v(\Pi_k) \geq t - 1$  for every  $k$ -subspace  $\Pi_k$  containing  $l^*$ ;
- (c)  $v(\Pi_k^*) = t - 1$ ;
- (d)  $|w| \geq |v| + 2$ .

### Case 1

Assume first that  $k = 1$ . Then  $\Pi_k^* = l^*$  is a line, and  $v(l^*) = t - 1$ , while the  $v$  weight of any other line is at least  $t$ . If  $n = 2$ , then  $|v| \geq (t + 1)q - 1$  by Lemma 3.5, which proves the theorem in this case. Now assume  $n \geq 3$  and let  $\Pi$  be a plane containing the line  $l^*$ . Then the weight function  $v$  restricted to the plane  $\Pi$  fulfills the requirements of Lemma 3.5, so  $v(\Pi) \geq (t + 1)q - 1$ . This is true for all the planes containing the line  $l^*$ , so clearly  $|v| \geq \theta_{n-2} \cdot ((t + 1)q - 1 - (t - 1)) + t - 1 = (t + 1)q^{n-1} + \theta_{n-2} - 2$ .

### Case 2

For  $n \geq 3$  and  $k \geq 2$  we will use induction on  $n$  to prove that

$$|v| \geq (t + 1)q^{n-k} + \theta_{n-k-1} - 2.$$

**Case 2a** Let  $V \in \Pi_k^* \setminus l^*$  be a point with  $v(V) = 0$ . Consider the quotient space  $\text{PG}(n, q)/V \cong \text{PG}(n - 1, q)$ , and the weight function  $\tilde{v}$  induced by  $v$  on  $\text{PG}(n - 1, q)$ . Clearly  $\tilde{v}(\text{PG}(n - 1, q)) = v(\text{PG}(n, q))$ . The plane  $\langle V, l^* \rangle$  corresponds to a line, and a  $k$ -space containing  $V$  corresponds to a  $(k - 1)$ -space. It is not hard to check that  $\tilde{v}$  fulfills requirements (a)-(c) with  $\langle V, l^* \rangle/V$  as  $l^*$  and  $\Pi_k^*/V$  as  $\Pi_{k-1}^*$ , and so by induction

$$\tilde{v}(\text{PG}(n - 1, q)) \geq (t + 1)q^{n-k} + \theta_{n-k-1} - 2.$$

**Case 2b** Suppose now that for all  $P \in \Pi_k^* \setminus l^*$ :  $v(P) > 0$ , but there is a point  $v(V) = 0$ . Then  $t - 1 \geq \theta_k - (q + 1)$ . Increase the weight of one point ( $\neq V$ ) of  $l^*$  by one to obtain the new weight function  $v'$ , which is now a  $t$ -fold  $(n - k)$ -blocking set of  $\text{PG}(n, q)$ . We will prove that  $|v'| \geq tq^{n-k} + \theta_{n-k} - 1$ . This is generally not true for  $t$ -fold  $(n - k)$ -blocking sets of  $\text{PG}(n, q)$ , only if  $t$  is large enough.

Assume, on the contrary, that  $|v'| \leq tq^{n-k} + \theta_{n-k} - 2$ . Then we can find a line  $\Sigma_1$  containing  $V$ , such that

$$v'(\Sigma_1) \leq \frac{t - (q^{k-1} + q^{k-2} + \dots + q)}{q^{k-1}},$$

because if all lines through  $V$  had  $v'$  weight more than

$$\frac{t - (q^{k-1} + q^{k-2} + \dots + q)}{q^{k-1}},$$

then all these weights would be at least  $\geq \frac{t - (q^{k-1} + q^{k-2} + \dots + q)}{q^{k-1}} + \frac{1}{q^{k-1}}$ , and then the total weight of  $v'$  would be

$$\begin{aligned} |v'| &\geq \left( \frac{t - q^{k-1} - q^{k-2} - \dots - q}{q^{k-1}} + \frac{1}{q^{k-1}} \right) \cdot \theta_{n-1} \\ &= tq^{n-k} + \left( \frac{t}{q^{k-1}} - \frac{q^k + q^{k-1} + \dots + q^2}{q^{k-1}} \right) \theta_{n-2} - \frac{q^{k-1} + q^{k-2} + \dots + q}{q^{k-1}} + \frac{\theta_{n-1}}{q^{k-1}} \\ &> tq^{n-k} + \frac{q^{n-1} + q^{n-2} + \dots + q^k}{q^{k-1}} = tq^{n-k} + \theta_{n-k} - 1. \end{aligned}$$

We will now prove that if  $1 \leq j \leq k-2$  and  $\Sigma_j$  is a  $j$ -space with

$$v'(\Sigma_j) \leq \frac{t - (q^{k-j} + q^{k-j-1} + \dots + q)}{q^{k-j}},$$

then we can find a  $(j+1)$ -space  $\Sigma_{j+1} \supset \Sigma_j$ , with

$$v'(\Sigma_{j+1}) \leq \frac{t - (q^{k-j-1} + \dots + q)}{q^{k-j-1}}.$$

If this were not true, then we would have

$$\begin{aligned} |v'| &> \left( \frac{t - (q^{k-j-1} + \dots + q)}{q^{k-j-1}} - v'(\Sigma_j) \right) \cdot \theta_{n-j-1} + v'(\Sigma_j) \\ &\geq \left( \frac{t - (q^{k-j-1} + \dots + q)}{q^{k-j-1}} - \frac{t - (q^{k-j} + q^{k-j-1} + \dots + q)}{q^{k-j}} \right) \cdot \theta_{n-j-1} \\ &\quad + \frac{t - (q^{k-j} + q^{k-j-1} + \dots + q)}{q^{k-j}} = tq^{n-k} + \theta_{n-k} + 1. \end{aligned}$$

Thus we can find a  $(k-1)$ -space  $\Sigma_{k-1}$ , with  $v'(\Sigma_{k-1}) \leq \frac{t-q}{q}$ . But all  $k$ -spaces containing  $\Sigma_{k-1}$  have  $v'$  weight at least  $t$ , so

$$|v'| \geq \left( t - \frac{t}{q} + 1 \right) \cdot \theta_{n-k} + \frac{t-q}{q} = tq^{n-k} + \theta_{n-k} - 1,$$

a contradiction.

**Case 2c** There is one more case remaining to be proved: if  $v(P) > 0$  for all points  $P \in \text{PG}(n, q)$ . Then let  $m := \min_P v(P)$  and let  $\tilde{v} := v - m$ . Then

$\tilde{v}$  fulfills requirements (a)-(c) with  $\tilde{t} := t - m \cdot \theta_k$ . Cases 2a and 2b prove  $|\tilde{v}| \geq \tilde{t}q^{n-k} + \theta_{n-k} - 2$  and then

$$\begin{aligned} |v| = |\tilde{v}| + m \cdot \theta_n &\geq (t - m \cdot \theta_k)q^{n-k} + \theta_{n-k} - 2 + m \cdot \theta_n \\ &= tq^{n-k} + \theta_{n-k} - 2 + m\theta_{n-k-1}. \end{aligned}$$

□

## 5 Examples

In this section we investigate the sharpness of Theorem 1.3. We are looking for weighted  $t$ -fold  $(n - k)$ -blocking sets of size  $(t + 1)q^{n-k} + \theta_{n-k-1}$ , which contain two different minimal  $t$ -fold  $(n - k)$ -blocking sets.

### 5.1 The case $t = 1$

**Example 1** Let  $\Pi^1$  and  $\Pi^2$  be two  $(n - k)$ -dimensional subspaces of  $\text{PG}(n, q)$  meeting in an  $(n - k - 1)$ -dimensional subspace. Then  $B := \Pi^1 \cup \Pi^2$  contains two different minimal 1-fold  $(n - k)$ -blocking sets ( $\Pi^1$  and  $\Pi^2$ ), and  $|B| = 2q^{n-k} + \theta_{n-k-1}$ . ■

**Corollary 5.1.** *Theorem 1.3 is sharp, if  $t = 1$ .*

The following proposition is a corollary of Theorem 1.3, but in fact equivalent to it if  $t = 1$  and  $k = 1$ . Corollary 5.3 can also be found in [13].

**Proposition 5.2.** *Let  $B$  be a minimal 1-fold  $(n - 1)$ -blocking set of  $\text{PG}(n, q)$ , and  $P \in B$ . Then there are at least  $\geq 2q^{n-1} + \theta_{n-2} - |B|$  tangents through  $P$ .*

*Proof.* Suppose that there are  $k$  tangents through  $P$ . Take points  $P_1, P_2, \dots, P_k$ , one from each of the tangents,  $P_i \neq P$ . Clearly  $(B \setminus \{P\}) \cup \{P_1, \dots, P_k\}$  is a 1-fold  $(n - 1)$ -blocking set. It contains a minimal 1-fold  $(n - 1)$ -blocking set  $B'$ , and  $B \neq B'$ . Thus  $B \cup \{P_1, \dots, P_k\}$  contains two different minimal 1-fold  $(n - 1)$ -blocking sets, so  $|B| + k \geq 2q^{n-1} + \theta_{n-2}$ . □

**Corollary 5.3.** *Let  $B$  be any 1-fold  $(n - 1)$ -blocking set of  $\text{PG}(n, q)$ , and  $P \in B$  an essential point of  $B$ . Then there are at least  $\geq 2q^{n-1} + \theta_{n-2} - |B|$  tangents through  $P$ .*

**Construction 1** Let  $B$  be a 1-fold  $(n - 1)$ -blocking set which has a point  $P \in B$ , through which there are exactly  $2q^{n-1} + \theta_{n-2} - |B|$  tangents to  $B$ . Then adding a point to every tangent will result in a 1-fold  $(n - 1)$ -blocking set of size  $2q^{n-1} + \theta_{n-2}$ , which contains two different minimal 1-fold  $(n - 1)$ -blocking sets.  $\square$

**Construction 2** Embed construction 1 in an  $(n - k + 1)$ -dimensional subspace of  $\text{PG}(n, q)$  to obtain 1-fold  $(n - k)$ -blocking sets of size  $2q^{n-k} + \theta_{n-k-1}$ , which contain two different minimal 1-fold  $(n - k)$ -blocking sets.  $\square$

Note that blocking sets used in the above construction exist: the so called Rédei type blocking sets always contain points which are on exactly  $2q^{n-1} + \theta_{n-2} - |B|$  tangents (see [10]).

## 5.2 The case $t \geq 2$

We will use the following notation: for the multisets  $B_1$  and  $B_2$ , with associated weight functions  $w_1$  and  $w_2$  respectively,  $B_1 \cup B_2$  will denote the multiset defined by the weight function  $\max\{w_1, w_2\}$ , while  $B_1 + B_2$  will denote the multiset defined by the weight function  $w_1 + w_2$ .

Note that the proof of Lemma 3.5 yields that for  $n = 2$ ,  $k = 1$  it is not possible to have  $v(\text{PG}(2, q)) = (t + 1)q - 1$ , if  $t \geq q + 1$ , and so the proof of Theorem 1.3 yields that the bound cannot be sharp if  $t \geq q + 1$ . Also from the proofs of Lemma 3.5 and Theorem 1.3 it follows that if  $t \leq q - 2$  and  $B$  is a weighted  $t$ -fold  $(n - k)$ -blocking set which contains two different minimal  $t$ -fold  $(n - k)$ -blocking sets and  $|B| = (t + 1)q^{n-k} + \theta_{n-k-1}$ , then only points on one line (the line  $l^*$ ) can be multiple points.

**Example 2** Let  $\Pi$  be a plane of  $\text{PG}(n, k)$ , let  $l_1, l_2, \dots, l_t$  be different lines in  $\Pi$  through a common point  $P$ , and  $l_{t+1}$  a further line of  $\Pi$ , with  $P \notin l_{t+1}$ . Then the multiset  $B := (l_1 + l_2 + \dots + l_t) \cup l_{t+1}$  is a  $t$ -fold 1-blocking set in  $\text{PG}(n, q)$ ,  $|B| = t(q + 1) + (q + 1 - t) = (t + 1)q + 1$ , and  $l_1 + l_2 + \dots + l_t$  and  $l_1 \cup (l_2 + \dots + l_t) \cup l_{t+1}$  are two minimal  $t$ -fold 1-blocking sets contained in  $B$ ; the latter one differs from  $B$  only in the point  $P$ .  $\square$

**Corollary 5.4.** *Theorem 1.3 is sharp if  $k = n - 1$ ,  $2 \leq t \leq q$ .*

The following proposition is again a corollary of Theorem 1.3, which is in fact equivalent to it if  $k = 1$ . For  $n = 2$  and with an upper bound on the size of  $B$ , it can also be found in [6].

**Proposition 5.5.** *Let  $B$  be a minimal  $t$ -fold  $(n - 1)$ -blocking set of  $\text{PG}(n, q)$ , and  $P \in B$ . Then there are at least  $\geq (t + 1)q^{n-1} + \theta_{n-2} - |B|$   $t$ -secants through  $P$ .*

*Proof.* Suppose that there are  $k$   $t$ -secants through  $P$ . Take points  $P_1, P_2, \dots, P_k$ , one from each of the  $t$ -secants,  $P_i \neq P$ . Clearly the  $t$ -fold  $(n - 1)$ -blocking set  $B \setminus \{P\} + \{P_1, \dots, P_k\}$  contains a minimal  $t$ -fold  $(n - 1)$ -blocking set  $B'$ , and  $B \neq B'$ . Thus  $B + \{P_1, \dots, P_k\}$  contains two different minimal  $t$ -fold  $(n - 1)$ -blocking sets, so  $|B| + k \geq (t + 1)q^{n-1} + \theta_{n-2}$ .  $\square$

**Construction 3** Let  $B$  be a minimal  $t$ -fold  $(n - 1)$ -blocking set which has a point  $P \in B$ , through which there are exactly  $(t + 1)q^{n-1} + \theta_{n-2} - |B|$   $t$ -secants to  $B$ . Then adding a point to every  $t$ -secant will result in a  $t$ -fold  $(n - 1)$ -blocking set of size  $(t + 1)q^{n-1} + \theta_{n-2}$  and containing two different minimal  $t$ -fold  $(n - 1)$ -blocking sets.  $\square$

**Construction 4** Embed Construction 3 in an  $(n - k + 1)$ -dimensional subspace of  $\text{PG}(n, q)$  to obtain  $t$ -fold  $(n - k)$ -blocking sets of size  $(t + 1)q^{n-k} + \theta_{n-k-1}$ , which contain two different minimal  $t$ -fold  $(n - k)$ -blocking sets.  $\square$

For  $n = 2$ ,  $k = 1$  and  $2 \leq t \leq q$  one can find  $t$ -fold 1-blocking sets in  $\text{PG}(2, q)$  which have points that are on exactly  $(t + 1)q + 1 - |B|$   $t$ -secants to  $B$ : take the sum of  $t$  Rédei type blocking sets which have a common Rédei line, and share exactly one point, that is not on the Rédei line. Example 2 is a special case of this: the sum of  $t$  lines sharing a common point. Then, with Construction 4, we get examples for  $n \geq 3$ ,  $k = n - 1$  and  $1 \leq t \leq q$ . Unfortunately, for  $t \geq 2$ ,  $n \geq 3$  and  $k = 1$ , in the minimal  $t$ -fold  $(n - 1)$ -blocking sets examined by the author all points have at least  $t\theta_{n-1} - (q + 1 - t)q^{n-2} - |B|$   $t$ -secants to  $B$ . Thus it may be conjectured that the correct bound in Theorem 1.3 should be

$$t\theta_{n-k} + (q + 1 - t)q^{n-k-1}.$$

## Acknowledgment

I am very grateful to Leo Storme for his valuable suggestions and help.

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