VARIETIES WITH DEFINABLE FACTOR CONGRUENCES

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ABSTRACT. We study direct product representations of algebras in varieties. We collect several conditions expressing that these representations are definable in a first-order-logic sense, among them the concept of Definable Factor Congruences (DFC). The main results are that DFC is a Mal'cev property and that it is equivalent to all other conditions formulated; in particular we prove that $\mathcal V$ has DFC if and only if $\mathcal V$ has $\vec 0$ & $\vec 1$ and Boolean Factor Congruences. We also obtain an explicit first-order definition Φ of the kernel of the canonical projections via the terms associated to the Mal'cev condition for DFC, in such a manner that it is preserved by taking direct products and direct factors. The main tool is the use of central elements, which are a generalization of both central idempotent elements in rings with identity and neutral complemented elements in a bounded lattice.

1. Introduction

An algebra is a nonempty set together with an arbitrary but fixed collection of finitary operations. A variety is an equationally-definable class of algebras over the same language. A congruence of an algebra A is the kernel $\{(a,b) \in A : f(a) = f(b)\}$ of a homomorphism f with domain A; it is a factor congruence of A if f is a projection onto a direct factor of A. Thus, a direct product representation is determined by the pair of complementary factor congruences given by the canonical projections.

In this universal-algebraic setting, one key concept for the deeper study of direct product representations is that of *central element*. This tool can be developed fruitfully in varieties with $\vec{0}$ & $\vec{1}$, which we now define.

A variety with $\vec{0}$ & $\vec{1}$ is a variety \mathcal{V} in which there exist unary terms $0_1(w), \ldots, 0_l(w), 1_1(w), \ldots, 1_l(w)$ such that

$$\mathcal{V} \models \vec{0}(w) = \vec{1}(w) \rightarrow x = y,$$

where w, x and y are distinct variables, $\vec{0} = (0_1, \dots, 0_l)$ and $\vec{1} = (1_1, \dots, 1_l)$.

The terms $\vec{0}$ and $\vec{1}$ are analogous, in a rather general manner, to identity (top) and null (bottom) elements in rings (lattices), and its existence in a variety is equivalent to the fact that no nontrivial algebra in the variety has a trivial subalgebra. Throughout this paper we will assume that \mathcal{V} is a variety with $\vec{0}$ & $\vec{1}$ such that the terms $\vec{0}$ and $\vec{1}$ are closed. Of course, this can be achieved when the language has a constant symbol, and we will make this assumption in order to simplify and clarify our treatment. The proofs remain valid in the general case.

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If $\vec{a} \in A^l$ and $\vec{b} \in B^l$, we will write $[\vec{a}, \vec{b}]$ in place of $((a_1, b_1), \dots, (a_l, b_l)) \in (A \times B)^l$. If $A \in \mathcal{V}$, we say that $\vec{e} \in A^l$ is a *central element* of A if there exists an isomorphism $A \to A_1 \times A_2$ such that

$$\vec{e} \mapsto [\vec{0}, \vec{1}].$$

Two central elements \vec{e}, \vec{f} will be called *complementary* if there exists an isomorphism $A \to A_1 \times A_2$ such that $\vec{e} \mapsto [\vec{0}, \vec{1}]$ and $\vec{f} \mapsto [\vec{1}, \vec{0}]$.

Central elements are a generalization of both central idempotent elements in rings with identity and neutral complemented elements in a bounded lattice. It is well known that in the classical cases, central elements are powerful tools since they translate the concepts of factor congruence and of direct product representation into first-order logic. In the general framework of an arbitrary variety $\mathcal V$ with $\vec 0$ & $\vec 1$, we may write this definability property in the following fashion:

there exists a first-order formula $\Phi(x, y, \vec{z}, \vec{w})$ in the language of \mathcal{V} such that for all $A, B \in \mathcal{V}$, and $a, c \in A, b, d \in B$,

$$A \times B \models \Phi((a,b),(c,d),[\vec{0},\vec{1}],[\vec{1},\vec{0}])$$
 if and only if $a=c$.

It was shown in [12] that in varieties with $\vec{0}$ & $\vec{1}$ that have the Fraser-Horn-Hu property, central elements have this property, and the formula Φ may be chosen of the form $\exists \bigwedge p = q$. Moreover, when the variety is congruence modular, Φ can be chosen to be a conjunction of equations. One last family of varieties with $\vec{0}$ & $\vec{1}$ (which generalize the above-mentioned examples) is that of varieties in which factor congruences are *compact* [13].

Which is the most general context in which central elements concentrate the information concerning the direct product representations? One answer is given by the following condition: each pair of complementary central elements determines uniquely a pair of complementary factor congruences associated to them. This property may be written as:

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for A \in \mathcal{V}, the map (\theta, \theta^*) \mapsto \text{unique } (\vec{e}, \vec{f}) \in A^l \times A^l \text{ satisfying } \vec{0} \theta \vec{e} \theta^* \vec{1} \text{ and } \vec{1} \theta \vec{f} \theta^* \vec{0}
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is a bijection between the set of pairs of complementary factor congruences of A and the set of pairs of complementary central elements of A.

Another important concept involved in this paper is that of an algebra with Boolean Factor Congruences (BFC), that is, an algebra in which the set of factor congruences is a distributive sublattice of its congruence lattice. Though it is not apparent, this concept is intimately connected to the direct product construction. In the classical work of Chang, Jónsson and Tarski [3] it is proved that BFC is equivalent to the strict refinement property (a strengthening of the property that states that every two direct product representations have a common refinement). It is also noteworthy that in several works on sheaf representations [2, 4, 5, 9] BFC has played an important role. For example, in Bigelow and Burris [1] it is shown that in a variety with BFC the Boolean product representations with directly indecomposable factors are unique and coincide with the Pierce sheaf [10]. We refer the reader to the work of Willard [14] as a key reference on BFC.

In the present work, we will prove that all the conditions mentioned in this introduction are indeed equivalent for the class of varieties with $\vec{0}$ & $\vec{1}$.

Theorem 1.1. Let V be a variety with $\vec{0} \& \vec{1}$. The following are equivalent:

- (1) \mathcal{V} has the Weak Determining Property: for $A \in \mathcal{V}$, the map $(\theta, \theta^*) \mapsto unique\ (\vec{e}, \vec{f}) \in A^l \times A^l$ satisfying $\vec{0} \theta \vec{e} \theta^* \vec{1}$ and $\vec{1} \theta \vec{f} \theta^* \vec{0}$ is a bijection between the set of pairs of complementary factor congruences of A and the set of pairs of complementary central elements of A.
- (2) V has the Determining Property: For A ∈ V, the map (θ, θ*) → unique e ∈ A^l satisfying 0 θ e θ* 1 is a bijection between the set of pairs of complementary factor congruences of A and the set of central elements of A.
- (3) V has Definable Factor Congruences (DFC): There exists a first-order formula $\Phi(x, y, \vec{z})$ in the language of V such that for all $A, B \in V$, and $a, c \in A$, $b, d \in B$,

$$A \times B \models \Phi((a,b),(c,d),[\vec{0},\vec{1}])$$
 if and only if $a = c$.

(4) There exists a first-order formula $\Phi(x, y, \vec{z}, \vec{w})$ in the language of V such that for all $A, B \in V$, and $a, c \in A, b, d \in B$,

$$A \times B \models \Phi((a,b),(c,d),[\vec{0},\vec{1}],[\vec{1},\vec{0}])$$
 if and only if $a = c$.

(5) V has BFC.

Moreover, when the above equivalent conditions hold, the formula Φ in (3) can be chosen to be preserved by direct products and direct factors and for every $A \in \mathcal{V}$, the map

$$\vec{e} \mapsto \Phi^A(\cdot, \cdot, \vec{e})$$

is a bijection between the set of central elements and the Boolean algebra of factor congruences of A.

We now briefly sketch the contents of each section. In Section 2 we give a Mal'cev condition for a Left Determining Property to be defined there; this condition is entirely analogous to a Mal'cev condition for BFC. The terms obtained in this section are the building blocks for our definability constructions. Section 3 provides an explicit formula Φ satisfying (3) of Theorem 1.1. This formula is constructed in such a way that it is preserved by direct products and direct factors; this last assertion is proved in the Appendix. In Section 4, we characterize in first-order logic the (pairs of complementary) central elements in a variety with DFC and show that the coordinates (in a direct product representation) of a central element are central elements. Several results obtained in the previous sections are compiled in Section 5 to finish the proof of the Main Theorem. Two examples are treated in Section 6. In §6.1 we present a variety with $\vec{0}$ & $\vec{1}$ that does not have DFC; this shows that our definitions do not trivialize. In §6.2 we give an optimal formula for the case of semilattices.

Throughout this paper the following notation will be used. For $A \in \mathcal{V}$ and $\vec{a}, \vec{b} \in A^n$, $\operatorname{Cg}^A(\vec{a}, \vec{b})$ will denote the congruence generated by the set $\{(a_k, b_k) : 1 \leq k \leq n\}$. The symbols ∇ and Δ will stand for the universal and trivial congruences, respectively. We will use $\theta \times \theta^* = \Delta$ in place of " θ and θ^* are complementary factor

congruences". The term algebra (in the language of \mathcal{V}) and the \mathcal{V} -free algebra on X will be denoted by T(X) and F(X), respectively. The i-th component of an element a in a direct product $\Pi_i A_i$ will be called $a^i = \pi_i(a)$; hence, if $a \in A_0 \times A_1$, $a = (a^0, a^1)$. If elements a, b of an algebra A are related by a congruence $\theta \in \text{Con } A$, we will write interchangeably $(a, b) \in \theta$, $a \theta b$ or $a \stackrel{\theta}{=} b$. This notation generalizes to tuples, viz., $\vec{a} \theta \vec{b}$ means $(a_i, b_i) \in \theta$ for all i.

2. The Left Determining Property: A Mal'cev condition

We will use the following Left Determining Property:

For every
$$A \in \mathcal{V}$$
, $\vec{e} \in A^l$ and $\varphi, \varphi^*, \theta, \theta^* \in \text{Con } A$, if $\varphi \times \varphi^* = \Delta$, $\theta \times \theta^* = \Delta$, $\vec{0} \theta \vec{e} \theta^* \vec{1}$ and $\vec{0} \varphi \vec{e} \varphi^* \vec{1}$, then $\theta = \varphi$.

It is not difficult to see the implications: Determining Property \Rightarrow Left Determining Property \Rightarrow Weak Determining Property.

The following theorem gives a Mal'cev condition for the Left Determining Property. Let s_i, t_i be (2i+l)-ary terms (in the language of \mathcal{V}) for each $i=1,\ldots,n$ and let A be an algebra in the language of \mathcal{V} (not necessarily in \mathcal{V}). For $(c,d,\vec{c},a_1,b_1,\ldots,a_n,b_n)\in A^{2+l+2n}$, we define $\sigma(c,d,\vec{e},a_1,b_1,\ldots,a_n,b_n)$ to be the tuple $(x,y,\vec{z},x_1,y_1,\ldots,x_n,y_n)$ given by the following recursion:

$$x := c$$
 $x_j := s_j(x, y, \vec{z}, x_1, y_1, \dots, x_{j-1}, y_{j-1}),$
 $y := c$ $y_j := b_j,$
 $\vec{z} := \vec{0}.$

We define σ^* , ρ , ρ^* analogously.

•
$$\sigma^*(c,d,\vec{e},a_1,b_1,\ldots,a_n,b_n):=(x,y,\vec{z},x_1,y_1,\ldots,x_n,y_n)$$
 where:
$$x:=c \qquad x_j:=t_j(x,y,\vec{z},x_1,y_1,\ldots,x_{j-1},y_{j-1}),$$
 $y:=d \qquad y_j:=b_j,$ $\vec{z}:=\vec{1}.$

• $\rho(c, d, \vec{e}, a_1, b_1, \dots, a_n, b_n) := (x, y, \vec{z}, x_1, y_1, \dots, x_n, y_n)$ where:

$$x := c$$
 $x_j := a_j,$ $y := d$ $y_j := s_j(x, y, \vec{z}, x_1, y_1, \dots, x_{j-1}, y_{j-1}),$ $\vec{z} := \vec{0}.$

•
$$\rho^*(c,d,\vec{e},a_1,b_1,\ldots,a_n,b_n) := (x,y,\vec{z},x_1,y_1,\ldots,x_n,y_n)$$
 where:
$$x := c \qquad x_j := a_j,$$

$$y := d \qquad y_j := t_j(x,y,\vec{z},x_1,y_1,\ldots,x_{j-1},y_{j-1}),$$

$$\vec{z} := \vec{1}.$$

We first state without proof a lemma concerning these functions.

Lemma 2.1. For every $c, d, \vec{e}, a_1, b_1, \ldots, a_n, b_n \in A$, we have the following identities:

$$(2.1) \quad \operatorname{Cg}(c,d) \vee \operatorname{Cg}(\vec{e},\vec{0}) \vee \bigvee_{i} \operatorname{Cg}(a_{i},s_{i}(c,d,\vec{e},a_{1},b_{1},\ldots,a_{i-1},b_{i-1})) \\ = \operatorname{Cg}((c,d,\vec{e},a_{1},b_{1},\ldots,a_{n},b_{n}),\sigma(c,d,\vec{e},a_{1},b_{1},\ldots,a_{n},b_{n})), \\ \operatorname{Cg}(\vec{e},\vec{1}) \vee \bigvee_{i} \operatorname{Cg}(a_{i},t_{i}(c,d,\vec{e},\ldots,a_{i-1},b_{i-1})) \\ = \operatorname{Cg}((c,d,\vec{e},\ldots,a_{n},b_{n}),\sigma^{*}(c,d,\vec{e},\ldots,a_{n},b_{n})), \\ \operatorname{Cg}(\vec{e},\vec{0}) \vee \bigvee_{i} \operatorname{Cg}(b_{i},s_{i}(c,d,\vec{e},\ldots,a_{i-1},b_{i-1})) \\ = \operatorname{Cg}((c,d,\vec{e},\ldots,a_{n},b_{n}),\rho(c,d,\vec{e},\ldots,a_{n},b_{n})), \\ \operatorname{Cg}(\vec{e},\vec{1}) \vee \bigvee_{i} \operatorname{Cg}(b_{i},t_{i}(c,d,\vec{e},\ldots,a_{i-1},b_{i-1})) \\ = \operatorname{Cg}((c,d,\vec{e},\ldots,a_{n},b_{n}),\rho^{*}(c,d,\vec{e},\ldots,a_{n},b_{n})).$$

In the proofs that follow, we will repeatedly find elements in an algebra that solve congruential "equations" of the form

$$a \stackrel{\theta}{\equiv} x \stackrel{\theta^*}{\equiv} b$$

when $\theta \times \theta^* = \Delta$. Using the functions σ , σ^* , ρ and ρ^* just defined, we can assert conclusions from the way elements such as x are constructed. This is the content of the following immediate consequences of Lemma 2.1:

Corollary 2.2. Given $c, d, \vec{e} \in A$ and $\theta, \theta^* \in \text{Con } A$ such that $\vec{0} \theta \vec{e} \theta^* \vec{1}$ and $c \theta d$, for every a_i and b_i with i = 1, ..., n such that

$$s_1(c, d, \vec{e}) \stackrel{\theta}{\equiv} a_1 \stackrel{\theta^*}{\equiv} t_1(c, d, \vec{e}),$$

$$s_2(c, d, \vec{e}, a_1, b_1) \stackrel{\theta}{\equiv} a_2 \stackrel{\theta^*}{\equiv} t_2(c, d, \vec{e}, a_1, b_1),$$
...

$$s_{i+1}(c, d, \vec{e}, a_1, b_1, \dots, a_i, b_i) \stackrel{\theta}{=} a_{i+1} \stackrel{\theta^*}{=} t_{i+1}(c, d, \vec{e}, a_1, b_1, \dots, a_i, b_i)$$

we have

(2.2)
$$t(c, d, \vec{e}, a_1, b_1, \dots, a_n, b_n) \stackrel{\theta}{\equiv} t(\sigma(c, d, \vec{e}, a_1, b_1, \dots, a_n, b_n)),$$
$$t(c, d, \vec{e}, a_1, b_1, \dots, a_n, b_n) \stackrel{\theta^*}{\equiv} t(\sigma^*(c, d, \vec{e}, a_1, b_1, \dots, a_n, b_n)),$$

for every (2n+l+2)-ary term t in the language of V.

The next result is entirely analogous.

Corollary 2.3. Suppose $c, d, \vec{e} \in A$ and $\varphi, \varphi^* \in \text{Con } A$ such that $\vec{0} \varphi \vec{e} \varphi^* \vec{1}$. If a_i and b_i satisfy

$$s_1(c,d,\vec{e}) \stackrel{\varphi}{=} b_1 \stackrel{\varphi^*}{=} t_1(c,d,\vec{e}),$$

$$\dots$$

$$s_{i+1}(c,d,\vec{e},a_1,b_1,\dots,a_i,b_i) \stackrel{\varphi}{=} b_{i+1} \stackrel{\varphi^*}{=} t_{i+1}(c,d,\vec{e},a_1,b_1,\dots,a_i,b_i),$$

we obtain

(2.3)
$$t(c, d, \vec{e}, a_1, b_1, \dots, a_n, b_n) \stackrel{\varphi}{=} t(\rho(c, d, \vec{e}, a_1, b_1, \dots, a_n, b_n)),$$
$$t(c, d, \vec{e}, a_1, b_1, \dots, a_n, b_n) \stackrel{\varphi^*}{=} t(\rho^*(c, d, \vec{e}, a_1, b_1, \dots, a_n, b_n))$$

for every (2n+l+2)-ary term t in the language of V.

We will also need the following (Grätzer's) version of Mal'cev's key observation on principal congruences.

Lemma 2.4. Let A be any algebra and let $a, b \in A$, $\vec{a}, \vec{b} \in A^n$. Then $(a, b) \in \operatorname{Cg}^A(\vec{a}, \vec{b})$ if and only if there exist (n+m)-ary terms $p_1(\vec{x}, \vec{u}), \ldots, p_k(\vec{x}, \vec{u})$, with k odd, and $\vec{u} \in A^m$ such that:

$$\begin{split} a &= p_1(\vec{a}, \vec{u}), \\ p_i(\vec{b}, \vec{u}) &= p_{i+1}(\vec{b}, \vec{u}), i \ odd, \\ p_i(\vec{a}, \vec{u}) &= p_{i+1}(\vec{a}, \vec{u}), i \ even, \\ p_k(\vec{b}, \vec{u}) &= b. \end{split}$$

We will use $|\alpha|$ to denote the length of a word α and ε will denote the empty word.

Theorem 2.5. Let V be a variety with $\vec{0} \& \vec{1}$. V has Left Determining Property if and only if there exist integers N=2k and n, (2i+l)-ary terms s_i and t_i for each $i=1,\ldots,n$, and for every word α in the alphabet $\{1,\ldots,N\}$ of length no greater than N there are terms L_{α} , R_{α} such that

$$|\alpha| = N$$

(2.4)
$$L_{\alpha}(\rho(\vec{X})) \approx R_{\alpha}(\rho(\vec{X})),$$
$$L_{\alpha}(\rho^{*}(\vec{X})) \approx R_{\alpha}(\rho^{*}(\vec{X})),$$

$$|\alpha| = 0$$

(2.5)
$$x \approx L_{\varepsilon}(x, y, \vec{z}, x_1, y_1, \dots, x_n, y_n), \\ R_{\varepsilon}(x, y, \vec{z}, x_1, y_1, \dots, x_n, y_n) \approx y,$$

(2.6)
$$L_{\varepsilon}(\rho(\vec{X})) \approx L_1(\rho(\vec{X})),$$

(2.7)
$$R_j(\rho(\vec{X})) \approx L_{j+1}(\rho(\vec{X})) \quad \text{if } 1 \leq j \leq N-1,$$

(2.8)
$$R_N(\rho(\vec{X})) \approx R_{\varepsilon}(\rho(\vec{X})).$$

$$0 < |\alpha| < N$$

If $|\alpha|$ is even, then

(2.9)
$$L_{\alpha}(\rho(\vec{X})) \approx L_{\alpha 1}(\rho(\vec{X})),$$

$$(2.10) R_{\alpha i}(\rho(\vec{X})) \approx L_{\alpha(j+1)}(\rho(\vec{X})) if 1 \le j \le k-1,$$

(2.11)
$$R_{\alpha k}(\rho(\vec{X})) \approx R_{\alpha}(\rho(\vec{X})),$$
$$L_{\alpha}(\rho^{*}(\vec{X})) \approx L_{\alpha(k+1)}(\rho^{*}(\vec{X})),$$

(2.12)
$$R_{\alpha j}(\rho^*(\vec{X})) \approx L_{\alpha(j+1)}(\rho^*(\vec{X})) \quad \text{if } k+1 \le j \le N-1,$$
$$R_{\alpha N}(\rho^*(\vec{X})) \approx R_{\alpha}(\rho^*(\vec{X})).$$

If $|\alpha|$ is odd, then

$$L_{\alpha}(\sigma(\vec{X})) \approx L_{\alpha 1}(\sigma(\vec{X})),$$

(2.13)
$$R_{\alpha j}(\sigma(\vec{X})) \approx L_{\alpha(j+1)}(\sigma(\vec{X})) \quad \text{if } 1 \leq j \leq k-1,$$

$$R_{\alpha k}(\sigma(\vec{X})) \approx R_{\alpha}(\sigma(\vec{X})),$$

$$L_{\alpha}(\sigma^*(\vec{X})) \approx L_{\alpha(k+1)}(\sigma^*(\vec{X})),$$

(2.14)
$$R_{\alpha j}(\sigma^*(\vec{X})) \approx L_{\alpha(j+1)}(\sigma^*(\vec{X})) \quad \text{if } k+1 \le j \le N-1,$$
$$R_{\alpha N}(\sigma^*(\vec{X})) \approx R_{\alpha}(\sigma^*(\vec{X})),$$

where $\vec{X} = (x, y, \vec{z}, x_1, y_1, \dots, x_n, y_n)$ and σ , σ^* , ρ and ρ^* are defined relative to s_i , t_i , on $T(\vec{X})$.

Proof. (\Leftarrow) Assume the existence of the terms, and suppose $\varphi \times \varphi^* = \Delta$, $\theta \times \theta^* = \Delta$, $\vec{0} \theta e \theta^* \vec{1}$, $\vec{0} \varphi e \varphi^* \vec{1}$, and $c \theta d$. We want to see $c \varphi d$. There exist unique a_i, b_i satisfying the following relations:

$$s_1(c, d, \vec{e}) \stackrel{\theta}{=} a_1 \stackrel{\theta^*}{=} t_1(c, d, \vec{e}),$$

$$s_1(c, d, \vec{e}) \stackrel{\varphi}{=} b_1 \stackrel{\varphi^*}{=} t_1(c, d, \vec{e}),$$

. . .

$$s_{j+1}(c, d, \vec{e}, a_1, b_1, \dots, a_j, b_j) \stackrel{\theta}{\equiv} a_{j+1} \stackrel{\theta^*}{\equiv} t_{j+1}(c, d, \vec{e}, a_1, b_1, \dots, a_j, b_j),$$

$$s_{j+1}(c, d, \vec{e}, a_1, b_1, \dots, a_j, b_j) \stackrel{\varphi}{\equiv} b_{j+1} \stackrel{\varphi^*}{\equiv} t_{j+1}(c, d, \vec{e}, a_1, b_1, \dots, a_j, b_j).$$

Note that their definition combines the schemes in Corollaries 2.2 and 2.3. So, by equations (2.2) and (2.3) we have, taking $t := L_{\alpha}, R_{\alpha}$:

(2.15)

$$L_{\alpha}(\sigma(c,d,\vec{e},a_1,b_1,\dots)) \stackrel{\theta}{=} L_{\alpha}(c,d,\vec{e},a_1,b_1,\dots,a_n,b_n) \stackrel{\theta^*}{=} L_{\alpha}(\sigma^*(c,d,\vec{e},a_1,b_1,\dots)),$$

$$L_{\alpha}(\rho(c,d,\vec{e},a_1,b_1,\dots)) \stackrel{\varphi}{=} L_{\alpha}(c,d,\vec{e},a_1,b_1,\dots,a_n,b_n) \stackrel{\varphi^*}{=} L_{\alpha}(\rho^*(c,d,\vec{e},a_1,b_1,\dots)),$$

$$(2.16)$$

$$R_{\alpha}(\sigma(c,d,\vec{e},a_1,b_1,\dots)) \stackrel{\theta}{\equiv} R_{\alpha}(c,d,\vec{e},a_1,b_1,\dots,a_n,b_n) \stackrel{\theta^*}{\equiv} R_{\alpha}(\sigma^*(c,d,\vec{e},a_1,b_1,\dots)),$$

$$R_{\alpha}(\rho(c,d,\vec{e},a_1,b_1,\dots)) \stackrel{\varphi}{\equiv} R_{\alpha}(c,d,\vec{e},a_1,b_1,\dots,a_n,b_n) \stackrel{\varphi^*}{\equiv} R_{\alpha}(\rho^*(c,d,\vec{e},a_1,b_1,\dots))$$

for every α . We'll prove inductively that

(2.17)
$$L_{\alpha}(c, d, \vec{e}, a_1, b_1, \dots, a_n, b_n) = R_{\alpha}(c, d, \vec{e}, a_1, b_1, \dots, a_n, b_n)$$

for all $\alpha \neq \varepsilon$. Take α such that $|\alpha| = N$. Then

$$L_{\alpha}(c, d, \vec{e}, a_1, b_1, \dots, a_n, b_n) \stackrel{\varphi}{=} L_{\alpha}(\rho(c, d, \vec{e}, a_1, b_1, \dots, a_n, b_n))$$
 by (2.15)

$$= R_{\alpha}(\rho(c, d, \vec{e}, a_1, b_1, \dots, a_n, b_n))$$
 using (2.4)

$$\stackrel{\varphi}{=} R_{\alpha}(c, d, \vec{e}, a_1, b_1, \dots, a_n, b_n)$$
 by (2.16)

and

$$L_{\alpha}(c, d, \vec{e}, a_1, b_1, \dots, a_n, b_n) \stackrel{\varphi^*}{\equiv} L_{\alpha}(\rho^*(c, d, \vec{e}, a_1, b_1, \dots, a_n, b_n)) \quad \text{by (2.15)}$$

$$= R_{\alpha}(\rho^*(c, d, \vec{e}, a_1, b_1, \dots, a_n, b_n)) \quad \text{using (2.4)}$$

$$\stackrel{\varphi^*}{\equiv} R_{\alpha}(c, d, \vec{e}, a_1, b_1, \dots, a_n, b_n) \quad \text{by (2.16)}.$$

Hence $(L_{\alpha}(c,d,\vec{e},a_1,b_1,\ldots,a_n,b_n), R_{\alpha}(c,d,\vec{e},a_1,b_1,\ldots,a_n,b_n)) \in \varphi \cap \varphi^* = \Delta$ and then $L_{\alpha}(c,d,\vec{e},a_1,b_1,\ldots,a_n,b_n) = R_{\alpha}(c,d,\vec{e},a_1,b_1,\ldots,a_n,b_n)$.

Take $\alpha \neq \varepsilon$ of odd length and assume

$$L_{\alpha j}(c, d, \vec{e}, a_1, b_1, \dots, a_n, b_n) = R_{\alpha j}(c, d, \vec{e}, a_1, b_1, \dots, a_n, b_n)$$

for every j = 1, ..., N. We check that

$$L_{\alpha}(c,d,\vec{e},a_1,b_1,\ldots,a_n,b_n) \stackrel{\theta}{=} R_{\alpha}(c,d,\vec{e},a_1,b_1,\ldots,a_n,b_n),$$

$$L_{\alpha}(c,d,\vec{e},a_{1},b_{1},\ldots,a_{n},b_{n}) \stackrel{\theta}{\equiv} L_{\alpha}(\sigma(c,d,\vec{e},a_{1},b_{1},\ldots,a_{n},b_{n})) \quad \text{by (2.15)}$$

$$= L_{\alpha 1}(\sigma(c,d,\vec{e},a_{1},b_{1},\ldots,a_{n},b_{n})) \quad \text{by (2.13)}$$

$$\stackrel{\theta}{\equiv} L_{\alpha 1}(c,d,\vec{e},a_{1},b_{1},\ldots,a_{n},b_{n}) \quad \text{by (2.15)}$$

$$= R_{\alpha 1}(c,d,\vec{e},a_{1},b_{1},\ldots,a_{n},b_{n}) \quad \text{by ind. hypothesis}$$

$$\stackrel{\theta}{\equiv} R_{\alpha 1}(\sigma(c,d,\vec{e},a_{1},b_{1},\ldots,a_{n},b_{n})) \quad \text{by (2.16)}$$

$$\stackrel{\theta}{\equiv} \cdots \quad \text{using (2.13)}...$$

$$= R_{\alpha k}(\sigma(c,d,\vec{e},a_{1},b_{1},\ldots,a_{n},b_{n})) \quad \text{using (2.13)}$$

$$\stackrel{\theta}{\equiv} R_{\alpha}(c,d,\vec{e},a_{1},b_{1},\ldots,a_{n},b_{n}).$$

In the same way for θ^* , we have

$$L_{\alpha}(c, d, \vec{e}, a_{1}, b_{1}, \dots, a_{n}, b_{n}) \stackrel{\theta^{*}}{\equiv} L_{\alpha}(\sigma^{*}(c, d, \vec{e}, a_{1}, b_{1}, \dots, a_{n}, b_{n})) \quad \text{by (2.15)}$$

$$= L_{\alpha(k+1)}(\sigma^{*}(c, d, \vec{e}, a_{1}, b_{1}, \dots, a_{n}, b_{n})) \quad \text{by (2.14)}$$

$$\stackrel{\theta^{*}}{\equiv} L_{\alpha(k+1)}(c, d, \vec{e}, a_{1}, b_{1}, \dots, a_{n}, b_{n}) \quad \text{by (2.15)}$$

$$= R_{\alpha(k+1)}(c, d, \vec{e}, a_{1}, b_{1}, \dots, a_{n}, b_{n}) \quad \text{by ind. hyp.}$$

$$\stackrel{\theta^{*}}{\equiv} R_{\alpha(k+1)}(\sigma^{*}(c, d, \vec{e}, a_{1}, b_{1}, \dots, a_{n}, b_{n})) \quad \text{by (2.16)}$$

$$\stackrel{\theta^{*}}{\equiv} \dots \quad \text{using (2.14)}$$

$$= R_{\alpha N}(\sigma^{*}(c, d, \vec{e}, a_{1}, b_{1}, \dots, a_{n}, b_{n})) \quad \text{using (2.14)}$$

$$\stackrel{\theta^{*}}{\equiv} R_{\alpha}(c, d, \vec{e}, a_{1}, b_{1}, \dots, a_{n}, b_{n}) \quad \text{by (2.16)}.$$

Hence $(L_{\alpha}(c, d, \vec{e}, a_1, b_1, \dots, a_n, b_n), R_{\alpha}(c, d, \vec{e}, a_1, b_1, \dots, a_n, b_n)) \in \theta \cap \theta^* = \Delta$, and therefore they are equal.

If $\alpha \neq \varepsilon$ has even length, then

$$L_{\alpha}(c,d,\vec{e},a_{1},b_{1},\ldots,a_{n},b_{n}) \stackrel{\mathcal{E}}{=} L_{\alpha}(\rho(c,d,\vec{e},a_{1},b_{1},\ldots,a_{n},b_{n})) \quad \text{by (2.15)}$$

$$= L_{\alpha 1}(\rho(c,d,\vec{e},a_{1},b_{1},\ldots,a_{n},b_{n})) \quad \text{by (2.9)}$$

$$\stackrel{\mathcal{E}}{=} L_{\alpha 1}(c,d,\vec{e},a_{1},b_{1},\ldots,a_{n},b_{n}) \quad \text{by (2.15)}$$

$$= R_{\alpha 1}(c,d,\vec{e},a_{1},b_{1},\ldots,a_{n},b_{n}) \quad \text{by ind. hypothesis}$$

$$\stackrel{\mathcal{E}}{=} R_{\alpha 1}(\rho(c,d,\vec{e},a_{1},b_{1},\ldots,a_{n},b_{n})) \quad \text{by (2.16)}$$

$$\stackrel{\mathcal{E}}{=} \ldots \quad \text{using (2.10)}...$$

$$= R_{\alpha k}(\rho(c,d,\vec{e},a_{1},b_{1},\ldots,a_{n},b_{n})) \quad \text{using (2.11)}$$

$$\stackrel{\mathcal{E}}{=} R_{\alpha}(\rho(c,d,\vec{e},a_{1},b_{1},\ldots,a_{n},b_{n})) \quad \text{using (2.11)}$$

$$\stackrel{\mathcal{E}}{=} R_{\alpha}(\rho(c,d,\vec{e},a_{1},b_{1},\ldots,a_{n},b_{n})) \quad \text{by (2.16)}$$

$$\text{proves } (L_{\alpha}(c,d,\vec{e},a_{1},b_{1},\ldots,a_{n},b_{n}), R_{\alpha}(c,d,\vec{e},a_{1},b_{1},\ldots,a_{n},b_{n})) \quad \text{by (2.15)}$$

$$= L_{\alpha}(c,d,\vec{e},a_{1},b_{1},\ldots,a_{n},b_{n}) \quad \text{by (2.15)}$$

$$= L_{\alpha}(c,d,\vec{e},a_{1},b_{1},\ldots,a_{n},b_{n}) \quad \text{by (2.15)}$$

$$= L_{\alpha}(c,d,\vec{e},a_{1},b_{1},\ldots,a_{n},b_{n}) \quad \text{by ind. hyp.}$$

$$\stackrel{\mathcal{E}^{*}}{=} L_{\alpha}(c,d,\vec{e},a_{1},b_{1},\ldots,a_{n},b_{n}) \quad \text{by ind. hyp.}$$

$$\stackrel{\mathcal{E}^{*}}{=} R_{\alpha}(c,d,\vec{e},a_{1},b_{1},\ldots,a_{n},b_{n}) \quad \text{by (2.16)}$$

$$\stackrel{\mathcal{E}^{*}}{=} \ldots \quad \text{using (2.12)}$$

$$= R_{\alpha}(\rho^{*}(c,d,\vec{e},a_{1},b_{1},\ldots,a_{n},b_{n})) \quad \text{using (2.12)}$$

$$\stackrel{\mathcal{E}^{*}}{=} R_{\alpha}(c,d,\vec{e},a_{1},b_{1},\ldots,a_{n},b_{n}) \quad \text{by (2.16)}$$

completes this case. Finally, we have:

$$c = L_{\varepsilon}(\rho(c, d, \vec{e}, a_1, b_1, \dots, a_n, b_n)) \qquad \text{using } (2.5)$$

$$= L_1(\rho(c, d, \vec{e}, a_1, b_1, \dots, a_n, b_n)) \qquad \text{by } (2.6)$$

$$\stackrel{\varphi}{\equiv} L_1(c, d, \vec{e}, a_1, b_1, \dots, a_n, b_n) \qquad \text{by } (2.15)$$

$$= R_1(c, d, \vec{e}, a_1, b_1, \dots, a_n, b_n) \qquad \text{by } (2.17)$$

$$\stackrel{\varphi}{\equiv} R_1(\rho(c, d, \vec{e}, a_1, b_1, \dots, a_n, b_n)) \qquad \text{by } (2.16)$$

$$\stackrel{\varphi}{\equiv} \dots \qquad \qquad \text{using } (2.7) \dots$$

$$= R_N(\rho(c, d, \vec{e}, a_1, b_1, \dots, a_n, b_n))$$

$$= R_{\varepsilon}(\rho(c, d, \vec{e}, a_1, b_1, \dots, a_n, b_n)) \qquad \text{using } (2.8)$$

$$= d \qquad \text{using } (2.5).$$

This proves $(c,d) \in \varphi$.

 (\Rightarrow) For each set of variables Y, define

Take the following congruences on F:

$$Y^* := Y \cup \{x_{p,q} : p, q \in T(Y)\} \cup \{y_{p,q} : p, q \in T(Y)\},$$

$$Y^{0*} := Y,$$

$$Y^{(n+1)*} := (Y^{n*})^*,$$

$$Y^{\infty} := \bigcup_{n \ge 1} Y^{n*},$$

where $x_{p,q}$ and $y_{p,q}$ are new variables. Take $Z := \{x, y, z_1, \ldots, z_l\}$ and $F := F(Z^{\infty})$. Define the index of $p \in T(Z^{\infty})$ as $ind(p) = \min\{j : p \in T(Z^{j*})\}$; it is evident that if $ind(x_{p,q}) \leq ind(x_{r,s})$, neither p nor q can be terms depending on $x_{r,s}$. The same holds for $ind(x_{p,q}) \leq ind(y_{r,s})$ and symmetrically, and for $ind(y_{p,q}) \leq ind(y_{r,s})$.

$$\begin{split} \theta &:= \operatorname{Cg}(\vec{0}, \vec{z}) \vee \operatorname{Cg}(x, y) \vee \bigvee \{ \operatorname{Cg}(p, x_{p,q}) : p, q \in F \}, \qquad \delta_0 = \epsilon_0 := \Delta^F, \\ \theta^* &:= \operatorname{Cg}(\vec{1}, \vec{z}) \vee \bigvee \{ \operatorname{Cg}(x_{p,q}, q) : p, q \in F \}, \qquad \delta_{n+1} := (\theta \vee \epsilon_n) \cap (\theta^* \vee \epsilon_n), \\ \varphi &:= \operatorname{Cg}(\vec{0}, \vec{z}) \vee \bigvee \{ \operatorname{Cg}(p, y_{p,q}) : p, q \in F \}, \qquad \epsilon_{n+1} := (\varphi \vee \delta_n) \cap (\varphi^* \vee \delta_n), \\ \varphi^* &:= \operatorname{Cg}(\vec{1}, \vec{z}) \vee \bigvee \{ \operatorname{Cg}(y_{p,q}, q) : p, q \in F \}, \qquad \delta_\infty := \bigvee_{n \geq 0} \delta_n = \bigvee_{n \geq 0} \epsilon_n. \end{split}$$

By construction, $\varphi \circ \varphi^* = \theta \circ \theta^* = \nabla^F$, $\vec{0} \theta \vec{z} \theta^* \vec{1}$, $\vec{0} \varphi \vec{z} \varphi^* \vec{1}$, and $x \theta y$. Observe that if $(a,b) \in (\varphi \vee \delta_{\infty}) \cap (\varphi^* \vee \delta_{\infty})$, then there exists an $n \geq 0$ such that $(a,b) \in (\varphi \vee \delta_n) \cap (\varphi^* \vee \delta_n)$. But this congruence is exactly ϵ_{n+1} ; hence $(a,b) \in \epsilon_{n+1} \subseteq \delta_{\infty}$. We may conclude $(\varphi \vee \delta_{\infty}) \cap (\varphi^* \vee \delta_{\infty}) = \delta_{\infty}$. The same happens with θ and θ^* ; hence

$$(\varphi \vee \delta_{\infty})/\delta_{\infty} \times (\varphi^* \vee \delta_{\infty})/\delta_{\infty} = \Delta, \qquad (\theta \vee \delta_{\infty})/\delta_{\infty} \times (\theta^* \vee \delta_{\infty})/\delta_{\infty} = \Delta$$

in F/δ_{∞} . Then, by the Left Determining Property we have $(x/\delta_{\infty},y/\delta_{\infty}) \in (\varphi \vee \delta_{\infty})/\delta_{\infty}$ and hence $(x,y) \in \varphi \vee \delta_{\infty}$. We may find an even integer N=2k such that $(x,y) \in \varphi \circ^{2N} \delta_N^N$, where δ_N^N is the result of replacing each occurrence of "V" in the definition of δ_N by \circ^N , the n-fold relational product.

We will inductively define terms L_{α} and R_{α} , for α a word of length at most N in the alphabet $\{1, \ldots, N\}$ such that:

$$(2.18) x = L_{\varepsilon}, y = R_{\varepsilon},$$

$$(2.19) (L_{\varepsilon}, L_1) \in \varphi, (R_N, R_{\varepsilon}) \in \varphi,$$

$$(2.20) (L_i, R_i) \in \delta_N^N if 1 \le i \le N,$$

(2.21)
$$(R_i, L_{(i+1)}) \in \varphi$$
 if $1 \le i \le N - 1$.

For $\alpha \neq \varepsilon$ with $|\alpha| < N$ an odd integer,

$$(2.22) (L_{\alpha}, L_{\alpha 1}) \in \theta, (R_{\alpha k}, R_{\alpha}) \in \theta,$$

$$(2.23) (L_{\alpha}, L_{\alpha(k+1)}) \in \theta^*, (R_{\alpha N}, R_{\alpha}) \in \theta^*,$$

$$(2.24) (L_{\alpha i}, R_{\alpha i}) \in \epsilon_{N-|\alpha|}^{N} if 1 \le i \le N,$$

$$(2.25)$$
 $(R_{\alpha i}, L_{\alpha(i+1)}) \in \theta$ if $1 \le i \le k-1$,

(2.26)
$$(R_{\alpha i}, L_{\alpha(i+1)}) \in \theta^*$$
 if $k+1 \le i \le N-1$

and for $\alpha \neq \varepsilon$ with $|\alpha| < N$ an even integer:

$$(2.27) (L_{\alpha}, L_{\alpha 1}) \in \varphi, (R_{\alpha k}, R_{\alpha}) \in \varphi,$$

$$(2.28) (L_{\alpha}, L_{\alpha(k+1)}) \in \varphi^*, (R_{\alpha N}, R_{\alpha}) \in \varphi^*,$$

$$(2.29) (L_{\alpha i}, R_{\alpha i}) \in \delta_{N-|\alpha|}^{N} if 1 \le i \le N,$$

$$(2.30) (R_{\alpha i}, L_{\alpha(i+1)}) \in \varphi if 1 \le i \le k-1,$$

(2.31)
$$(R_{\alpha i}, L_{\alpha(i+1)}) \in \varphi^*$$
 if $k+1 \le i \le N-1$.

We take $L_{\varepsilon} := x$ and $R_{\varepsilon} := y$. Since we know $(x, y) \in \varphi \circ^{2N} \delta_N^N$, we define L_i , R_i for i = 1, ..., N to be terms satisfying

$$x \varphi L_1 \delta_N^N R_1 \varphi L_2 \delta_N^N \cdots \varphi L_N \delta_N^N R_N \varphi y.$$

Note that these terms satisfy (2.18)–(2.31) whenever they can be checked.

Suppose we have defined the terms corresponding to words with length less than or equal to j and that they satisfy equations among (2.18)–(2.31) that involve words of length j or shorter. Then we shall define terms corresponding to words with length equal to j+1 such that the totality of terms defined satisfy equations among (2.18)–(2.31) which involve words of length j+1 or shorter. We have two cases:

Case 1: j odd. Take α , with $|\alpha| = j$. We have L_{α} and R_{α} and by (2.29), they satisfy $(L_{\alpha}, R_{\alpha}) \in \delta_{N-j+1}^{N} = (\theta \circ^{N} \epsilon_{N-j}^{N}) \cap (\theta^* \circ^{N} \epsilon_{N-j}^{N})$. We define $L_{\alpha i}$ and $R_{\alpha i}$ for $i = 1, \ldots, N$ such that:

(2.32)
$$L_{\alpha} \theta L_{\alpha 1} \epsilon_{N-j}^{N} R_{\alpha 1} \theta L_{\alpha 2} \cdots R_{\alpha(k-1)} \theta L_{\alpha k} \epsilon_{N-j}^{N} R_{\alpha k} \theta R_{\alpha},$$

$$L_{\alpha} \theta^{*} L_{\alpha(k+1)} \epsilon_{N-j}^{N} R_{\alpha(k+1)} \theta^{*} L_{\alpha(k+2)} \cdots L_{\alpha N} \epsilon_{N-j}^{N} R_{\alpha N} \theta^{*} R_{\alpha}.$$

The equations among (2.18)–(2.31) which involve terms L_{μ} and R_{μ} with $|\mu| = j+1$ are (2.22)–(2.26). All of them can be inferred from (2.32).

Case 2: j even. Take α , with $|\alpha| = j$. We define $L_{\alpha i}$ and $R_{\alpha i}$ for i = 1, ..., N. By (2.24) and by the definition of ϵ_{N-j+1}^N we may define our terms satisfying:

$$L_{\alpha} \varphi L_{\alpha 1} \delta_{N-j}^{N} R_{\alpha 1} \varphi L_{\alpha 2} \cdots R_{\alpha(k-1)} \varphi L_{\alpha k} \delta_{N-j}^{N} R_{\alpha k} \varphi R_{\alpha},$$

$$L_{\alpha} \varphi L_{\alpha(k+1)} \delta_{N-j}^{N} R_{\alpha(k+1)} \varphi^{*} L_{\alpha(k+2)} \cdots R_{\alpha N} \varphi^{*} R_{\alpha}.$$

From this we immediately conclude (2.27)–(2.31).

Let $V \subseteq Z^{\infty}$ be a finite set of variables such that if we replace θ , θ^* , φ and φ^* , respectively, by the following compact congruences:

$$\theta_0 := \operatorname{Cg}(\vec{0}, \vec{z}) \vee \operatorname{Cg}(x, y) \vee \bigvee \{ \operatorname{Cg}(p, x_{p,q}) : x_{p,q} \in V \},$$

$$\theta^*_0 := \operatorname{Cg}(\vec{1}, \vec{z}) \vee \bigvee \{ \operatorname{Cg}(x_{p,q}, q) : x_{p,q} \in V \},$$

$$\varphi_0 := \operatorname{Cg}(\vec{0}, \vec{z}) \vee \bigvee \{ \operatorname{Cg}(p, y_{p,q}) : y_{p,q} \in V \},$$

$$\varphi^*_0 := \operatorname{Cg}(\vec{1}, \vec{z}) \vee \bigvee \{ \operatorname{Cg}(y_{p,q}, q) : y_{p,q} \in V \},$$

then we still obtain congruential relations (2.18)–(2.31) (excepting (2.20), (2.29) and (2.24)). It is clear that if we enlarge the set V to a new set X, the properties enumerated will still hold.

Let V_0 be the union of V and the (finite) set of variables occurring in terms L_{α} , R_{α} with α a word. Define:

$$V_{n+1} := V_n \cup \bigcup \{ Var(p), Var(q) : x_{p,q} \in V_n \text{ or } y_{p,q} \in V_n \}.$$

Hence, for some M we have $V_M = V_{M+1}$; set

$$X := (V_M \cup \{x_{p,q} : y_{p,q} \in V_M\} \cup \{y_{p,q} : x_{p,q} \in V_M\}) \setminus \{x, y, z_1, \dots, z_l\}.$$

Order X totally so that $ind: X \to \omega$ is nondecreasing and $x_{p,q}$ is the immediate predecessor of $y_{p,q}$, and add x, y, z_1, \ldots, z_l at the beginning; we have

$$\vec{X} = (x, y, \vec{z}, x_{s_1, t_1}, y_{s_1, t_1}, \dots, x_{s_n, t_n}, y_{s_n, t_n}) = (x, y, \vec{z}, x_1, y_1, \dots, x_n, y_n).$$

We may consider then $L_{\alpha} = L_{\alpha}(\vec{X})$ and the same for R_{α} , and by the remarks after the definition of Z^{∞} , we may assume $s_i = s_i(x, y, \vec{z}, x_1, y_1, \dots, x_{i-1}, y_{i-1})$ and the same for t_i . Finally, define $\sigma, \rho, \sigma^*, \rho^*$ on the term algebra $T(\vec{X})$ with respect to s_i, t_i .

We claim that these L's, R's, s's and t's satisfy the Mal'cev condition. Let's check it for identity (2.9). Take α with $0 < |\alpha| < N$ an even integer. By Lemma 2.1 we have

$$\varphi = \operatorname{Cg}(\vec{X}, \rho(\vec{X})).$$

Since we have $L_{\alpha} \varphi L_{\alpha 1}$ by equation (2.27), Lemma 2.4 gives us terms p_i such that for some tuple \vec{u} , F satisfies:

$$L_{\alpha} = p_1(\vec{X}, \vec{u}),$$

$$p_1(\rho(\vec{X}), \vec{u}) = p_2(\rho(\vec{X}), \vec{u}),$$

$$p_2(\vec{X}, \vec{u}) = p_3(\vec{X}, \vec{u}),$$

$$\dots$$

$$p_m(\rho(\vec{X}), \vec{u}) = L_{\alpha 1}.$$

Since $\vec{u} = \vec{u}(\vec{X}, \vec{Y})$ can be construed as members of $T(Z^{\infty})$, we obtain the following laws for \mathcal{V} :

$$L_{\alpha}(\vec{X}) \approx p_1(\vec{X}, \vec{u}(\vec{X}, \vec{Y})),$$

$$p_1(\rho(\vec{X}), \vec{u}(\vec{X}, \vec{Y})) \approx p_2(\rho(\vec{X}), \vec{u}(\vec{X}, \vec{Y})),$$

$$p_2(\vec{X}, \vec{u}(\vec{X}, \vec{Y})) \approx p_3(\vec{X}, \vec{u}(\vec{X}, \vec{Y})),$$

$$\dots$$

$$p_m(\rho(\vec{X}), \vec{u}(\vec{X}, \vec{Y})) \approx L_{\alpha 1}(\vec{X}).$$

Replacing \vec{X} by $\rho(\vec{X})$ everywhere and noting that $\rho(\rho(\vec{X})) = \rho(\vec{X})$, we have

$$L_{\alpha}(\rho(\vec{X})) \approx p_{1}(\rho(\vec{X}), \vec{u}(\rho(\vec{X}), \vec{Y})),$$

$$p_{1}(\rho(\vec{X}), \vec{u}(\rho(\vec{X}), \vec{Y})) \approx p_{2}(\rho(\vec{X}), \vec{u}(\rho(\vec{X}), \vec{Y})),$$

$$p_{2}(\rho(\vec{X}), \vec{u}(\rho(\vec{X}), \vec{Y})) \approx p_{3}(\rho(\vec{X}), \vec{u}(\rho(\vec{X}), \vec{Y})),$$

$$\cdots$$

$$p_{m}(\rho(\vec{X}), \vec{u}(\rho(\vec{X}), \vec{Y})) \approx L_{\alpha 1}(\rho(\vec{X})),$$

and by transitivity,

$$\mathcal{V} \models L_{\alpha}(\rho(\vec{X})) \approx L_{\alpha 1}(\rho(\vec{X})),$$

which is what we were looking for. The other identities are obtained similarly. \Box

The proof of the previous theorem follows the line of a proof for a Mal'cev condition for BFC. One such condition that closely parallels ours was personally communicated to us by R. Willard.

3. A CANONICAL FORM OF DFC

We will assume in this section that \mathcal{V} has the Determining Property. Since the Determining Property implies the Left Determining Property, we may define the following formulas in the language of \mathcal{V} :

$$\Psi_m := \bigwedge_{|\alpha| = m} \left(\left(\bigwedge_{\gamma \neq \varepsilon} L_{\alpha\gamma}(\vec{X}) = R_{\alpha\gamma}(\vec{X}) \right) \ \to \ L_{\alpha}(\vec{X}) = R_{\alpha}(\vec{X}) \right),$$

where every subindex varies over words of length less than or equal to N; so an expression of the form " $\bigwedge_{\gamma \neq \varepsilon} L_{\alpha \gamma} = R_{\alpha \gamma}$ " should be read as " $\bigwedge \{L_{\alpha \gamma} = R_{\alpha \gamma} : \gamma \neq \varepsilon \text{ and } |\alpha \gamma| \leq N\}$ ". Hence $\Psi_N = (\bigwedge_{|\beta|=N} L_{\beta}(\vec{X}) = R_{\beta}(\vec{X}))$. (The antecedent "vanishes".)

Lemma 3.1. Let $A \in \mathcal{V}$ and let $\varphi, \varphi^* \in \text{Con } A$ such that $\varphi \times \varphi^* = \Delta$ and $\vec{0} \varphi \vec{e} \varphi^* \vec{1}$. Then for all $c, d \in A$, A satisfies $\Phi_1(c, d, \vec{e})$, where

(3.1)
$$\Phi_1(x, y, \vec{z}) := \exists y_1 \forall x_1 \dots \exists y_n \forall x_n \bigwedge_{m=1}^k \Psi_{2m}$$

with n, k as in Theorem 2.5.

Proof. Take b_1 such that

$$s_1(c, d, \vec{e}) \stackrel{\varphi}{\equiv} b_1 \stackrel{\varphi^*}{\equiv} t_1(c, d, \vec{e}).$$

Assuming b_i is already chosen and a_i is given, define b_{i+1} such that

$$s_{i+1}(c,d,\vec{e},a_1,b_1,\ldots,a_i,b_i) \stackrel{\varphi}{\equiv} b_{i+1} \stackrel{\varphi^*}{\equiv} t_{i+1}(c,d,\vec{e},a_1,b_1,\ldots,a_i,b_i).$$

The construction of the b_i 's then corresponds to the equations in Corollary 2.3. Hence, (2.3) implies that A satisfies

$$L_{\alpha}(c,d,\vec{e},a_{1},b_{1},\ldots,a_{n},b_{n}) \stackrel{\varphi}{=} L_{\alpha}(\rho(c,d,\vec{e},a_{1},b_{1},\ldots,a_{n},b_{n})),$$

$$R_{\alpha}(c,d,\vec{e},a_{1},b_{1},\ldots,a_{n},b_{n}) \stackrel{\varphi}{=} R_{\alpha}(\rho(c,d,\vec{e},a_{1},b_{1},\ldots,a_{n},b_{n})),$$

$$L_{\alpha}(c,d,\vec{e},a_{1},b_{1},\ldots,a_{n},b_{n}) \stackrel{\varphi^{*}}{=} L_{\alpha}(\rho^{*}(c,d,\vec{e},a_{1},b_{1},\ldots,a_{n},b_{n})),$$

$$R_{\alpha}(c,d,\vec{e},a_{1},b_{1},\ldots,a_{n},b_{n}) \stackrel{\varphi^{*}}{=} R_{\alpha}(\rho^{*}(c,d,\vec{e},a_{1},b_{1},\ldots,a_{n},b_{n})),$$

for all α . These together with equations (2.4) imply that for each β with $|\beta| = N$,

(3.2)
$$L_{\beta}(c, d, \vec{e}, a_1, b_1, \dots, a_n, b_n) \stackrel{\varphi}{\equiv} R_{\beta}(c, d, \vec{e}, a_1, b_1, \dots, a_n, b_n),$$
$$L_{\beta}(c, d, \vec{e}, a_1, b_1, \dots, a_n, b_n) \stackrel{\varphi^*}{\equiv} R_{\beta}(c, d, \vec{e}, a_1, b_1, \dots, a_n, b_n).$$

Since $\varphi \cap \varphi^* = \Delta$, this yields $A \models \Psi_N(c, d, \vec{e}, a_1, b_1, \dots, a_n, b_n)$.

Take nonempty α with $|\alpha| < N$ even and suppose

$$A \models \bigwedge_{\gamma \neq \varepsilon} L_{\alpha\gamma}(c, d, \vec{e}, a_1, b_1, \dots, a_n, b_n) = R_{\alpha\gamma}(c, d, \vec{e}, a_1, b_1, \dots, a_n, b_n).$$

We can see that $L_{\alpha}(c, d, \vec{e}, a_1, b_1, \dots, a_n, b_n) \stackrel{\varphi}{\equiv} R_{\alpha}(c, d, \vec{e}, a_1, b_1, \dots, a_n, b_n)$, by using equations (2.9), (2.10) and (2.11) as follows:

$$L_{\alpha}(c,d,\vec{e},a_{1},b_{1},\ldots,a_{n},b_{n}) \stackrel{\varphi}{\equiv} L_{\alpha}(\rho(c,d,\vec{e},a_{1},b_{1},\ldots,a_{n},b_{n})) \quad \text{by (2.3)}$$

$$= L_{\alpha 1}(\rho(c,d,\vec{e},a_{1},b_{1},\ldots,a_{n},b_{n})) \quad \text{by (2.9)}$$

$$\stackrel{\varphi}{\equiv} L_{\alpha 1}(c,d,\vec{e},a_{1},b_{1},\ldots,a_{n},b_{n}) \quad \text{by (2.3)}$$

$$= R_{\alpha 1}(c,d,\vec{e},a_{1},b_{1},\ldots,a_{n},b_{n}) \quad \text{by hypothesis}$$

$$\stackrel{\varphi}{\equiv} R_{\alpha 1}(\rho(c,d,\vec{e},a_{1},b_{1},\ldots,a_{n},b_{n})) \quad \text{by (2.3)}$$

$$= \cdots \qquad \qquad \text{using (2.10), (3.2)}$$

$$= \cdots \qquad \qquad \text{and iterating...}$$

$$= R_{\alpha k}(\rho(c,d,\vec{e},a_{1},b_{1},\ldots,a_{n},b_{n})) \quad \text{using (2.11)}$$

$$\stackrel{\varphi}{\equiv} R_{\alpha}(\rho(c,d,\vec{e},a_{1},b_{1},\ldots,a_{n},b_{n})) \quad \text{by (2.3)}.$$

It can be proved in an entirely analogous fashion (by using equations (2.12)) that $L_{\alpha}(c, d, \vec{e}, a_1, b_1, \dots, a_n, b_n) \stackrel{\varphi^*}{=} R_{\alpha}(c, d, \vec{e}, a_1, b_1, \dots, a_n, b_n)$, which yields

$$A \models L_{\alpha}(c, d, \vec{e}, a_1, b_1, \dots, a_n, b_n) = R_{\alpha}(c, d, \vec{e}, a_1, b_1, \dots, a_n, b_n)$$

and we have proved the lemma.

Theorem 3.2. Let V be a variety with the Determining Property, let $A \in V$ and let $\theta, \theta^* \in \text{Con } A$ such that $\theta \times \theta^* = \Delta$, and $\vec{0} \theta \vec{e} \theta^* \vec{1}$. Then $c \stackrel{\theta}{\equiv} d$ if and only if $A \models \Phi_1(c, d, \vec{e}) \land \Phi_2(c, d, \vec{e})$, where Φ_1 and Φ_2 are defined as follows:

$$\Phi_1(x, y, \vec{z}) := \exists y_1 \forall x_1 \dots \exists y_n \forall x_n \bigwedge_{m=1}^k \Psi_{2m},
\Phi_2(x, y, \vec{z}) := \exists x_1 \forall y_1 \dots \exists x_n \forall y_n \bigwedge_{m=1}^k \Psi_{2m-1}.$$

Proof. Since the definition of Φ_1 is the same as the one given by formula (3.1) in Lemma 3.1, we only have to worry about $\Phi_2(c, d, \vec{e})$.

 (\Rightarrow) Assume $c \stackrel{\theta}{\equiv} d$. Much in the same way as in the proof of Lemma 3.1, define a_1 such that

$$s_1(c,d,\vec{e}) \stackrel{\theta}{\equiv} a_1 \stackrel{\theta^*}{\equiv} t_1(c,d,\vec{e});$$

and assuming a_i is already chosen and b_i is given, let

$$s_{i+1}(c, d, \vec{e}, a_1, b_1, \dots, a_i, b_i) \stackrel{\theta}{\equiv} a_{i+1} \stackrel{\theta^*}{\equiv} t_{i+1}(c, d, \vec{e}, a_1, b_1, \dots, a_i, b_i).$$

This choice conforms to the pattern of Corollary 2.2, so we obtain

(3.3)
$$L_{\alpha}(c, d, \vec{e}, a_1, b_1, \dots, a_n, b_n) \stackrel{\theta}{=} L_{\alpha}(\sigma(c, d, \vec{e}, a_1, b_1, \dots, a_n, b_n)),$$
$$R_{\alpha}(c, d, \vec{e}, a_1, b_1, \dots, a_n, b_n) \stackrel{\theta}{=} R_{\alpha}(\sigma(c, d, \vec{e}, a_1, b_1, \dots, a_n, b_n)),$$

(3.4)
$$L_{\alpha}(c, d, \vec{e}, a_1, b_1, \dots, a_n, b_n) \stackrel{\theta^*}{\equiv} L_{\alpha}(\sigma^*(c, d, \vec{e}, a_1, b_1, \dots, a_n, b_n)),$$

$$R_{\alpha}(c, d, \vec{e}, a_1, b_1, \dots, a_n, b_n) \stackrel{\theta^*}{\equiv} R_{\alpha}(\sigma^*(c, d, \vec{e}, a_1, b_1, \dots, a_n, b_n))$$

by equations (2.2).

If we suppose now that

$$A \models \bigwedge_{\gamma \neq \varepsilon} L_{\alpha\gamma}(c, d, \vec{e}, a_1, b_1, \dots, a_n, b_n) = R_{\alpha\gamma}(c, d, \vec{e}, a_1, b_1, \dots, a_n, b_n)$$

for some α with $|\alpha| < N$ odd, we'll be able to prove $L_{\alpha}(c, d, \vec{e}, a_1, b_1, \dots, a_n, b_n) = R_{\alpha}(c, d, \vec{e}, a_1, b_1, \dots, a_n, b_n)$ by showing (in the same way as in Lemma 3.1) that:

- $L_{\alpha}(c,d,\vec{e},a_1,b_1,\ldots,a_n,b_n) \stackrel{\theta}{\equiv} R_{\alpha}(c,d,\vec{e},a_1,b_1,\ldots,a_n,b_n)$ (this can be accomplished using (2.13) and (3.3)), and
- $L_{\alpha}(c, d, \vec{e}, a_1, b_1, \dots, a_n, b_n) \stackrel{\theta^*}{\equiv} R_{\alpha}(c, d, \vec{e}, a_1, b_1, \dots, a_n, b_n)$ (by (2.14) and (3.4)).
- (\Leftarrow) Assume $A \models \Phi_2(c,d,\vec{e})$. Take b_1 such that

$$s_1(c, d, \vec{e}) \stackrel{\theta}{\equiv} b_1 \stackrel{\theta^*}{\equiv} t_1(c, d, \vec{e}).$$

Let a_1 be given by the outermost existential quantifier of Φ_2 .

Assuming b_i is already chosen and a_i is the corresponding witness for Φ_2 , let

$$s_{i+1}(c, d, \vec{e}, a_1, b_1, \dots, a_i, b_i) \stackrel{\theta}{\equiv} b_{i+1} \stackrel{\theta^*}{\equiv} t_{i+1}(c, d, \vec{e}, a_1, b_1, \dots, a_i, b_i).$$

This selection conforms to the scheme of Corollary 2.3 (with $\varphi := \theta$ and $\varphi^* := \theta^*$) and satisfies the matrix of Φ_1 , as was seen in the proof of Lemma 3.1. Hence we have, respectively,

(3.5)
$$L_{\alpha}(c, d, \vec{e}, a_1, b_1, \dots, a_n, b_n) \stackrel{\theta}{\equiv} L_{\alpha}(\rho(c, d, \vec{e}, a_1, b_1, \dots, a_n, b_n)),$$
$$R_{\alpha}(c, d, \vec{e}, a_1, b_1, \dots, a_n, b_n) \stackrel{\theta^*}{\equiv} R_{\alpha}(\rho(c, d, \vec{e}, a_1, b_1, \dots, a_n, b_n))$$

and

$$A \models \left(\bigwedge_{m=1}^{N} \Psi_{m}\right)(c, d, \vec{e}, a_{1}, b_{1}, \dots, a_{n}, b_{n}).$$

From an easy inspection of the form of Ψ_m , it can be deduced that

(3.6)
$$A \models \bigwedge_{j=1}^{N} L_{j}(c, d, \vec{e}, a_{1}, b_{1}, \dots, a_{n}, b_{n}) = R_{j}(c, d, \vec{e}, a_{1}, b_{1}, \dots, a_{n}, b_{n}).$$

Therefore,

$$c = L_{\varepsilon}(c, d, \vec{e}, a_1, b_1, \dots, a_n, b_n)$$

$$\stackrel{\theta}{\equiv} L_{\varepsilon}(\rho(c, d, \vec{e}, a_1, b_1, \dots, a_n, b_n))$$

$$= L_1(\rho(c, d, \vec{e}, a_1, b_1, \dots, a_n, b_n))$$

$$\stackrel{\theta}{\equiv} L_1(c, d, \vec{e}, a_1, b_1, \dots, a_n, b_n)$$

$$= R_1(c, d, \vec{e}, a_1, b_1, \dots, a_n, b_n)$$

$$\stackrel{\theta}{\equiv} R_1(\rho(c, d, \vec{e}, a_1, b_1, \dots, a_n, b_n))$$

$$\stackrel{\theta}{\equiv} R_1(\rho(c, d, \vec{e}, a_1, b_1, \dots, a_n, b_n))$$

$$\stackrel{\theta}{\equiv} L_2(\rho(c, d, \vec{e}, a_1, b_1, \dots, a_n, b_n))$$

$$\stackrel{\theta}{\equiv} \dots$$

$$= R_N(\rho(c, d, \vec{e}, a_1, b_1, \dots, a_n, b_n))$$

$$\stackrel{\theta}{\equiv} R_{\varepsilon}(\rho(c, d, \vec{e}, a_1, b_1, \dots, a_n, b_n))$$

$$= R_{\varepsilon}(\rho(c, d, \vec{e}, a_1, b_1, \dots, a_n, b_n))$$

$$\stackrel{\theta}{\equiv} R_{\varepsilon}(c, d, \vec{e}, a_1, b_1, \dots, a_n, b_n))$$

$$\stackrel{\theta}{\equiv} R_{\varepsilon}(c, d, \vec{e}, a_1, b_1, \dots, a_n, b_n)$$

$$\stackrel{\theta}{\equiv} R_{\varepsilon}(c, d, \vec{e}, a_1, b_1, \dots, a_n, b_n)$$

$$\stackrel{\theta}{\equiv} R_{\varepsilon}(c, d, \vec{e}, a_1, b_1, \dots, a_n, b_n)$$

$$\stackrel{\theta}{\equiv} R_{\varepsilon}(c, d, \vec{e}, a_1, b_1, \dots, a_n, b_n)$$

$$\stackrel{\theta}{\equiv} R_{\varepsilon}(c, d, \vec{e}, a_1, b_1, \dots, a_n, b_n)$$

$$\stackrel{\theta}{\equiv} R_{\varepsilon}(c, d, \vec{e}, a_1, b_1, \dots, a_n, b_n)$$

$$\stackrel{\theta}{\equiv} R_{\varepsilon}(c, d, \vec{e}, a_1, b_1, \dots, a_n, b_n)$$

$$\stackrel{\theta}{\equiv} R_{\varepsilon}(c, d, \vec{e}, a_1, b_1, \dots, a_n, b_n)$$

$$\stackrel{\theta}{\equiv} R_{\varepsilon}(c, d, \vec{e}, a_1, b_1, \dots, a_n, b_n)$$

$$\stackrel{\theta}{\equiv} R_{\varepsilon}(c, d, \vec{e}, a_1, b_1, \dots, a_n, b_n)$$

$$\stackrel{\theta}{\equiv} R_{\varepsilon}(c, d, \vec{e}, a_1, b_1, \dots, a_n, b_n)$$

$$\stackrel{\theta}{\equiv} R_{\varepsilon}(c, d, \vec{e}, a_1, b_1, \dots, a_n, b_n)$$

$$\stackrel{\theta}{\equiv} R_{\varepsilon}(c, d, \vec{e}, a_1, b_1, \dots, a_n, b_n)$$

$$\stackrel{\theta}{\equiv} R_{\varepsilon}(c, d, \vec{e}, a_1, b_1, \dots, a_n, b_n)$$

$$\stackrel{\theta}{\equiv} R_{\varepsilon}(c, d, \vec{e}, a_1, b_1, \dots, a_n, b_n)$$

Hence $c \stackrel{\theta}{\equiv} d$.

4. Central elements in a variety with DFC

In the Appendix, we will prove a preservation result (Theorem A.4) that implies that formula $\Phi_1 \wedge \Phi_2$ of Theorem 3.2 is preserved by taking direct factors and direct products (taking τ_{α} to be " $L_{\alpha}(\vec{X}) = R_{\alpha}(\vec{X})$ "). Call Φ the conjunction $\Phi_1 \wedge \Phi_2$.

Lemma 4.1. Assume V has the Determining Property. Then there is a set Σ of first-order formulas such that for every $A \in V$, \vec{e} , $\vec{f} \in A^l$ we have that \vec{e} and \vec{f} are complementary central elements if and only if $A \models \zeta(\vec{e}, \vec{f})$ for every $\zeta \in \Sigma$. Moreover, each formula in Σ is preserved by taking direct factors.

Proof. The following formulas in the language of \mathcal{V} will assert the properties needed to force $\Phi(\cdot, \cdot, \vec{e})$ and $\Phi(\cdot, \cdot, \vec{f})$ to define the pair of complementary factor congruences associated with \vec{e} and \vec{f} .

- $CAN(\vec{e}, \vec{f}) = \bigwedge_{i=1}^{l} \Phi(0_i, e_i, \vec{e}) \wedge \bigwedge_{i=1}^{l} \Phi(1_i, f_i, \vec{e})$ This formula says that \vec{e} is related to $\vec{0}$ and \vec{f} to $\vec{1}$ via $\Phi(\cdot, \cdot, \vec{e})$.
- $PROD(\vec{e}, \vec{f}) = \forall x, y \exists z \ \left(\Phi(x, z, \vec{e}) \land \Phi(z, y, \vec{f}) \right)$ The relational product of $\Phi(\cdot, \cdot, \vec{e})$ and $\Phi(\cdot, \cdot, \vec{f})$ is the universal congruence.
- $INT(\vec{e}, \vec{f}) = \forall x, y \ \left(\Phi(x, y, \vec{e}) \land \Phi(x, y, \vec{f}) \rightarrow x = y \right)$ Their intersection is Δ .
- $REF(\vec{e}, \vec{f}) = \forall x \ \Phi(x, x, \vec{e})$ $\Phi(\cdot, \cdot, \vec{e})$ is reflexive.
- $SYM(\vec{e}, \vec{f}) = \forall x, y, z \left(\Phi(x, y, \vec{e}) \land \Phi(y, z, \vec{e}) \land \Phi(z, x, \vec{f}) \rightarrow z = x \right)$
- $TRANS(\vec{e}, \vec{f}) = \forall x, y, z, u \left(\Phi(x, y, \vec{e}) \land \Phi(y, z, \vec{e}) \land \Phi(x, u, \vec{e}) \land \Phi(u, z, \vec{f}) \rightarrow u = z \right)$

The reader may verify that these two formulas (in conjunction with the previous ones) say that $\Phi(\cdot,\cdot,\vec{e})$ is symmetric and transitive.

• For each m-ary function symbol F, define:

$$PRES_{F}(\vec{e}, \vec{f}) = \forall u_{1}, v_{1}, \dots, u_{m}, v_{m}$$

$$\left(\bigwedge_{j} \Phi(u_{j}, v_{j}, \vec{e})\right) \wedge \Phi(F(u_{1}, \dots, u_{m}), z, \vec{e}) \wedge \Phi(z, F(v_{1}, \dots, v_{m}), \vec{f}) \rightarrow$$

$$\rightarrow z = F(v_{1}, \dots, v_{m}).$$

These formulas ensure $\Phi(\cdot,\cdot,\vec{e})$ is preserved by the basic operations of \mathcal{V} . Finally, define CAN', REF', SYM', TRANS' and $PRES'_F$ to be the result of interchanging \vec{e} with \vec{f} in CAN, REF, SYM, TRANS and $PRES_F$, respectively, and let Σ be the union of the following two sets:

$$\{CAN, PROD, INT, REF, SYM, TRANS, CAN', REF', SYM', TRANS'\},$$

 $\{PRES_F, PRES'_F : F \text{ a function symbol}\}.$

Now it is immediate to check that \vec{e} and \vec{f} are complementary central elements if they satisfy all formulas in Σ . To see the converse, note that if \vec{e} and \vec{f} are complementary central elements, there is an isomorphism $A \to A_0 \times A_1$ such that \vec{e} , \vec{f} correspond to $[\vec{0}, \vec{1}], [\vec{1}, \vec{0}]$, respectively, and Theorem 3.2 guarantees that Σ will hold.

To see that Σ is preserved by direct factors, we first note that each one of $CAN, CAN' \ PROD, PROD', REF$ and REF' is obtained by quantification of a formula preserved by taking direct factors by Theorem A.4. In the second place, the remainder of axioms in Σ are of the form $\forall \vec{x} \ (\tau(\vec{e}, \vec{f}, \vec{x}) \to x_i = x_j)$ with $REF(\vec{e}, \vec{f}) \land REF'(\vec{e}, \vec{f}) \to \exists \vec{x} \ \tau(\vec{e}, \vec{f}, \vec{x})$ universally valid, and since $\forall \vec{x} \ (\tau(\vec{e}, \vec{f}, \vec{x}) \to x_i = x_j) \land \exists \vec{x} \ \tau(\vec{e}, \vec{f}, \vec{x})$ is preserved by taking direct factors (whenever $\tau(\vec{e}, \vec{f}, \vec{x})$ is preserved by taking direct factors and direct products), we have the result.

Corollary 4.2. Assume V has the Determining Property. Then if $[\vec{e}_0, \vec{e}_1]$ is a central element of $A_0 \times A_1$, then \vec{e}_i is a central element of A_i , i = 0, 1.

Proof. Immediate by the previous lemma.

We use \mathcal{V}_{DI} to denote the class of directly indecomposable members of \mathcal{V} .

Corollary 4.3. If V is a variety over a finite language and if it has the Determining Property, then V_{DI} is axiomatizable by a set of first-order sentences.

Proof. The set $\Sigma = \Sigma(\vec{e}, \vec{f})$ in Lemma 4.1 is finite if the language is finite. Hence

$$\vec{0} \neq \vec{1} \land \forall \vec{e}, \vec{f} \land \Sigma(\vec{e}, \vec{f}) \rightarrow ((\vec{e} = \vec{0} \land \vec{f} = \vec{1}) \lor (\vec{e} = \vec{1} \land \vec{f} = \vec{0}))$$

together with axioms for \mathcal{V} defines \mathcal{V}_{DI} .

Lemma 4.4. The Determining Property implies BFC.

Proof. By Bigelow and Burris [1], we only need to check that if $A = A_0 \times A_1$, and θ is a factor congruence on A, then

$$\{((a,b),(c,b)):b\in A_1 \text{ and } \exists a',c'(a,a')\,\theta\,(c,c')\}\subseteq\theta.$$

Let $\vec{e} = [\vec{e}_0, \vec{e}_1]$ be the central element associated to θ , so " $x \theta y$ " is defined by $\Phi(x, y, \vec{e})$. We have

$$(a, a') \theta(c, c') \text{ iff } A_0 \times A_1 \models \Phi((a, a'), (c, c'), [\vec{e}_0, \vec{e}_1]).$$

By Theorem A.4, this implies

$$A_0 \models \Phi(a, c, \vec{e_0}).$$

Now Corollary 4.2 ensures \vec{e}_1 is central in A_1 , and hence $A_1 \models \Phi(b, b, \vec{e}_1)$. Since Φ is preserved by direct products, we obtain

$$A_0 \times A_1 \models \Phi((a,b),(c,b),[\vec{e}_0,\vec{e}_1]),$$

and then $(a, b) \theta (c, b)$.

5. The main theorem

Proof of Theorem 1.1. (5) \Rightarrow (2) Suppose we have a pair of complementary factor congruences φ and φ^* such that $\vec{0} \varphi \vec{e} \varphi^* \vec{1}$. Suppose now that also $\theta \times \theta^* = \Delta$ and $\vec{0} \theta \vec{e} \theta^* \vec{1}$. Then $\vec{0} \stackrel{\varphi \circ \theta^*}{\equiv} \vec{1}$ and hence $\varphi \vee \theta^* = \nabla$. So we have

$$(\varphi \vee \theta^*) \cap \theta = \theta.$$

By BFC, we obtain $\varphi \cap \theta = \theta$ and then $\varphi \subseteq \theta$. By symmetry, we obtain $\varphi = \theta$ and $\varphi^* = \theta^*$.

- $(2) \Rightarrow (3)$ Theorem 3.2.
- $(3) \Rightarrow (4)$ Obvious.
- $(4) \Rightarrow (1)$ Immediate.
- $(2)\Rightarrow(5)$ Lemma 4.4.
- (1) \Rightarrow (5) Define $\tilde{0}_i$ and $\tilde{1}_i$ with $i=1,\ldots,2l$ in the following way:

$$(\tilde{0}_1, \dots, \tilde{0}_{2l}) := (0_1, \dots 0_l, 1_1, \dots 1_l),$$

 $(\tilde{1}_1, \dots, \tilde{1}_{2l}) := (1_1, \dots 1_l, 0_1, \dots, 0_l).$

It can be easily checked (using the Weak Determining Property) that with these $\tilde{0}$'s and $\tilde{1}$'s we have the Determining Property. Since $(2)\Rightarrow(5)$, we have our result.

6. Examples

6.1. $\vec{0}$ & $\vec{1}$ does not imply BFC. The variety \mathcal{V}_L with language $\{+, *, 0, 1\}$ given by the following set of equations Σ :

$$x + 0 = x,$$

$$x + 1 = x * 1,$$

$$x * 0 = 0$$

has $\vec{0}$ & $\vec{1}$. Next we will define various algebras in \mathcal{V}_L . In the first place take $L_{\omega} := \langle \omega, +, *, 0, 1 \rangle$, where

$$0+1:=0,$$
 $0*1:=0,$ $1+1:=1,$ $x+0:=x,$ $x*0:=0$

for all $x \in \omega$ and

$$z + y := 2, \qquad z * y := 2$$

for all $z, y \in \omega$ not previously considered. For each $n \geq 2$, L_n will denote the subalgebra of L_{ω} with universe $n = \{0, 1, \dots, n-1\}$. Now define D_n to be the subalgebra of $L_2 \times L_{\omega}$ with universe $(2 \times n) \cup \{(1, n)\}$.

Define the following subsets of $2 \times \omega$:

$$P_0 := \{(0, j) \mid 3 \le j\},\$$

$$P_1 := \{(1, j) \mid 3 \le j\}.$$

Then $2 \times \omega = (2 \times 3) \cup P_0 \cup P_1$. Note that for all $z \in (2 \times \omega) \setminus \{(0,0),(1,0)\}$ and for all $x, y \in P_1$ we have:

(6.1)
$$\begin{aligned} x + z &= y + z, \\ z + x &= z + y, \\ x * z &= y * z, \\ z * x &= z * y. \end{aligned}$$

Lemma 6.1. Every injective partial function $f: D_n \to (L_2 \times L_n)$ which fixes $(2 \times 3) \cup P_0$ is a partial isomorphism between D_n and $L_2 \times L_n$.

Proof. It's straightforward to see (using equations (6.1)) that if $B \subseteq P_1$ and σ is any permutation of P_1 , then $(2 \times 3) \cup P_0 \cup B$ and $(2 \times 3) \cup P_0 \cup \sigma(B)$ are subalgebras of $L_2 \times L_{\omega}$ and

$$\bar{\sigma}(x) := \begin{cases} x & x \in (2 \times 3) \cup P_0, \\ \sigma(x) & x \in B \end{cases}$$

is an isomorphism between them.

Since f is a restriction of such an isomorphism $\bar{\sigma}$, it is a partial isomorphism. \square

Recall that $(\mathcal{V}_L)_{DI}$ is the class of directly indecomposable members of \mathcal{V}_L .

Lemma 6.2. Let V be a variety. If V_{DI} is axiomatizable by a set of first-order sentences, then it is finitely axiomatizable relative to V.

Proof. First note that an ultraproduct of directly decomposable algebras is again decomposable. Let Σ be a set of first-order sentences axiomatizing \mathcal{V}_{DI} . By way of contradiction, suppose that \mathcal{V}_{DI} is not finitely axiomatizable relative to \mathcal{V} . Hence,

for each finite $\Sigma_0 \subseteq \Sigma$ there exists $A_{\Sigma_0} \in \mathcal{V} \setminus \mathcal{V}_{DI}$ satisfying Σ_0 . Now it is easy to construct an ultraproduct U of these decomposable algebras in such a way that U satisfies Σ , an absurdity.

Theorem 6.3. $(\mathcal{V}_L)_{DI}$ is not axiomatizable by a set of first-order sentences.

Proof. We will first prove that player " \exists " has a winning strategy for the back-and-forth (or "Ehrenfeucht") game of length n-3 played on D_n and $L_2 \times L_n$. The strategy is as follows:

- If \forall chooses an element in $(2 \times 3) \cup P_0$ (in either algebra), \exists will choose the same element in the other algebra.
- If \forall chooses an element in P_1 , \exists will choose an element in the P_1 -part of the other algebra, which has not been chosen up to this point.

There are n-3 elements in $P_1 \cap (L_2 \times L_n)$, so these instructions work up to n-3 moves. Let's call g the partial function defined by this game. By Lemma 6.1, g is a partial isomorphism and we have proved our first claim.

Now suppose φ is a sentence such that $(\mathcal{V}_L)_{DI} \models \varphi$. By the above strategy we have that for every sufficiently large n, $D_n \models \varphi$ if and only if $L_2 \times L_n \models \varphi$. By taking n such that 2n + 1 = cardinal of D_n is a prime number, we obtain $D_n \in (\mathcal{V}_L)_{DI}$. We conclude that there are decomposable algebras satisfying φ , and hence $(\mathcal{V}_L)_{DI}$ cannot be defined by a single first-order sentence. Using Lemma 6.2 we have our result.

Corollary 6.4. V_L does not have DFC.

Proof. Since DFC is equivalent to the Determining Property, we can use Corollary 4.3 and Theorem 6.3.

An indication that $(\mathcal{V}_L)_{DI}$ might not be definable was discovered by using the "Universal Algebra Calculator" program, designed by Ralph Freese and Emil Kiss [6].

6.2. **Semilattice expansions.** Throughout this section, we will suppose that \mathcal{V} is a variety with $\vec{0}$ & $\vec{1}$ for which there exists a binary term \vee such that for every $A \in \mathcal{V}$, \vee^A is a semilattice operation on A. We will keep the assumption that the language of \mathcal{V} has at least one constant.

First, we observe that by Lemma 2.4 together with the observation that $(x,y) \in \nabla^F = \operatorname{Cg}^F(\vec{0},\vec{1})$ (where $F \in \mathcal{V}$ is the free algebra freely generated by $\{x,y\}$), we obtain (2+l)-ary terms $u_i(x,y,\vec{z})$, $i=1,\ldots,k$, such that the following identities hold in \mathcal{V} :

$$(6.2) x \approx u_1(x, y, \vec{0}),$$

$$u_i(x, y, \vec{1}) \approx u_{i+1}(x, y, \vec{1}) \text{ with } i \text{ odd},$$

$$u_i(x, y, \vec{0}) \approx u_{i+1}(x, y, \vec{0}) \text{ with } i \text{ even},$$

$$u_k(x, y, \vec{1}) \approx y.$$

Proposition 6.5. The formula

$$\Phi(x,y,\vec{z}) := \forall u \left(\bigwedge_{i=1}^k \left(u_i(x,y,\vec{0}) \vee u = u_i(x,y,\vec{z}) \vee u \right) \longrightarrow (x \vee u = y \vee u) \right)$$

satisfies (3) of Theorem 1.1 for V.

Proof. Let $A, B \in \mathcal{V}$, and $a \in A, b, d \in B$. First we prove that

$$A \times B \models \Phi((a, b), (a, d), [\vec{0}, \vec{1}]).$$

Suppose that for some (u, v) we have

$$A \times B \models \bigwedge_{i=1}^{k} u_i((a,b),(a,d),[\vec{0},\vec{0}]) \lor (u,v) = u_i((a,b),(a,d),[\vec{0},\vec{1}]) \lor (u,v).$$

Then

$$B \models \bigwedge_{i=1}^{k} u_i(b, d, \vec{0}) \lor v = u_i(b, d, \vec{1}) \lor v.$$

But the above equations in combination with (6.2) produce

$$b \lor v = d \lor v$$

and hence

$$(a,b) \lor (u,v) = (a,d) \lor (u,v).$$

Now suppose

$$A \times B \models \Phi((a,b),(c,d),[\vec{0},\vec{1}]).$$

The reader can check that considering $u = (a, \bigvee_{i=1}^k u_i(b, d, \vec{0}) \vee u_i(b, d, \vec{1}))$, it can be proved that $a \vee c = a$, and similarly with $u = (c, \bigvee_{i=1}^k u_i(b, d, \vec{0}) \vee u_i(b, d, \vec{1}))$ and $a \vee c = c$; hence a = c.

The following example will show that the complexity of formula Φ in the above proposition cannot be improved for the general case.

Proposition 6.6. Let \mathcal{V}^{\vee} be the variety with language $\{+, *, 0, 1, \vee\}$ defined by the axioms of \mathcal{V}_L plus identities saying that \vee is a semilattice operation for which $0 \vee 1 = 0$. Then there exists neither a positive nor an existential formula satisfying (4) of Theorem 1.1 for \mathcal{V}^{\vee} .

Proof. Define join-semilattice operations on L_2 , L_4 and L_5 such that they are totally ordered with the order given by 0 > 1 > 2 > 3 > 4. Suppose that Φ satisfies (4) of Theorem 1.1, and consider $L_5 \times L_2$.

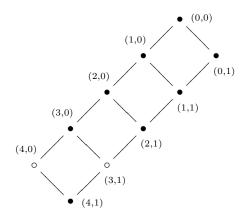


FIGURE 1. The semilattice structure of $L_5 \times L_2$.

The shaded dots in Figure 1 form a subalgebra of $L_5 \times L_2$; call it L. The reader may check that $F: L_4 \times L_2 \to L$, where

$$F(x) := \begin{cases} x & x \neq (3,1), \\ (4,1) & x = (3,1) \end{cases}$$

is an isomorphism. Since Φ witnesses (4) of Theorem 1.1 for \mathcal{V}^{\vee} , we have

$$L_4 \times L_2 \models \Phi((3,0),(3,1),(0,1),(1,0)).$$

Applying F everywhere,

$$L \models \Phi((3,0), (4,1), (0,1), (1,0)).$$

If we had Φ an existential formula, we would obtain

$$L_5 \times L_2 \models \Phi((3,0),(4,1),(0,1),(1,0)),$$

since L is a subalgebra of $L_5 \times L_2$. We would then conclude 3 = 4, an absurdity.

The fact that there is no positive formula Φ is an immediate consequence of the next two claims. We say that \mathcal{V} has *compact factor congruences* if every factor congruence of every algebra in \mathcal{V} is compact.

Claim 1. \mathcal{V}^{\vee} does not have compact factor congruences.

Proof. If \mathcal{V}^{\vee} had compact factor congruences, there would exist $\mathbf{0}_1(w), \ldots, \mathbf{0}_N(w)$, $\mathbf{1}_1(w), \ldots, \mathbf{1}_N(w)$ such that for every algebra $A = A_1 \times A_2 \in \mathcal{V}$, $(\lambda_1, \lambda_2) \in A$,

$$\ker \pi_1 = \operatorname{Cg}^A([\vec{\mathbf{0}}(\lambda_1), \vec{\mathbf{0}}(\lambda_2)], [\vec{\mathbf{0}}(\lambda_1), \vec{\mathbf{1}}(\lambda_2)]),$$

$$\ker \pi_2 = \operatorname{Cg}^A([\vec{\mathbf{1}}(\lambda_1), \vec{\mathbf{1}}(\lambda_2)], [\vec{\mathbf{0}}(\lambda_1), \vec{\mathbf{1}}(\lambda_2)]),$$

by Lemma 4 in [13]. Since the language contains constants, we can replace these new **0**'s and **1**'s by closed terms, and hence

$$\ker \pi_1 = \operatorname{Cg}^A([\vec{\mathbf{0}}, \vec{\mathbf{0}}], [\vec{\mathbf{0}}, \vec{\mathbf{1}}]),$$
$$\ker \pi_2 = \operatorname{Cg}^A([\vec{\mathbf{1}}, \vec{\mathbf{1}}], [\vec{\mathbf{0}}, \vec{\mathbf{1}}]).$$

Now, checking the axioms of \mathcal{V}^{\vee} , we conclude that for every closed term t in the language of \mathcal{V}^{\vee} , t is equivalent to 0 or t is equivalent to 1 over \mathcal{V}^{\vee} ; hence we should have

$$\ker \pi_1 = \operatorname{Cg}^A((0,0),(0,1)) \vee \operatorname{Cg}^A((1,0),(1,1)).$$

But the reader may check that if we take $A = L_5 \times L_2$, the equivalence relation depicted in Figure 2 is a congruence that contains the right hand side of the last equality and is clearly different from ker π_1 .

Claim 2. Let \mathcal{V} be a variety with $\vec{0}$ & $\vec{1}$. Suppose that Φ is a positive formula that satisfies (4) of Theorem 1.1. Then \mathcal{V} has compact factor congruences.

Proof. Let $A \in \mathcal{V}$. We will prove that if $\varphi \times \varphi^* = \Delta$, $\vec{0} \varphi \vec{e} \varphi^* \vec{1}$ and $\vec{1} \varphi \vec{f} \varphi^* \vec{0}$, then $\varphi = \operatorname{Cg}(\vec{0}, \vec{e}) \vee \operatorname{Cg}(\vec{1}, \vec{f})$, and hence is compact. Call $\theta := Cg(\vec{0}, \vec{e}) \vee \operatorname{Cg}(\vec{1}, \vec{f})$. Trivially, $\theta \subseteq \varphi$. Assume $x \varphi y$; by (4) of Theorem 1.1, we obtain $A \models \Phi(x, y, \vec{e}, \vec{f})$. Since Φ is positive, it is preserved by homomorphic images and then

$$A/\theta \models \Phi(x/\theta, y/\theta, \vec{e}/\theta, \vec{f}/\theta).$$

Equivalently,

$$A/\theta \models \Phi(x/\theta, y/\theta, \vec{0}/\theta, \vec{1}/\theta),$$

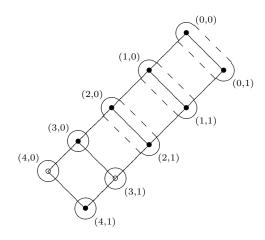


FIGURE 2. One congruence in $L_5 \times L_2$.

and we obtain $x/\theta = y/\theta$. This implies $(x,y) \in \theta$, hence $\varphi \subseteq \theta$, and we have the result.

7. Final considerations

We wish to mention that Theorem 1.1 is in some sense a consequence of Beth's definability theorem of first-order logic. This is because the first evidence that the Determining Property implies Definable Factor Congruences was obtained by a simple application of this theorem. Once we are sure about the concrete existence of a certain first-order formula¹ we know that all efforts dedicated to finding it are not sterile. Due to the expressive power of terms in algebra, it is common that first-order definitions obtained in this area are more manageable than in the general case of model theory, and this fact makes the task of searching for them a little easier.

Another tool that was important for our research for this paper, but was not included in the final presentation, is the Pierce sheaf. Working on this construction, we realized that there was some sense in developing a structural study of central elements when factor congruences are not compact.

Finally, we are indebted to Ross Willard for the terms $s_i(...)$ and $t_i(...)$, which appear in his (unpublished) version of a Mal'cev condition for Boolean Factor Congruences.

APPENDIX A. A PRESERVATION RESULT

For the rest of the section, N will be an even natural number.

Lemma A.1. For every word α in the alphabet $\{1, \ldots, N\}$ of length no greater than N, let $\tau_{\alpha} = \tau_{\alpha}(x, y, \vec{z}, x_1, y_1, \ldots, x_n, y_n)$ be a formula preserved by direct products and by taking direct factors. Define:

$$E_m := \bigwedge_{\substack{m \le |\alpha| \le N \\ |\alpha| \ even}} \left(\bigwedge_{\gamma \ne \varepsilon} \tau_{\alpha\gamma} \to \tau_{\alpha} \right), \qquad O_m := \bigwedge_{\substack{m \le |\alpha| \le N \\ |\alpha| \ odd}} \left(\bigwedge_{\gamma \ne \varepsilon} \tau_{\alpha\gamma} \to \tau_{\alpha} \right).$$

¹As is well known, the formulas given by Beth's theorem can be effectively found.

Then,

(1) For
$$2 \le m \le N$$
, m even, if

$$(\exists y_1 \forall x_1 \dots \exists y_n \forall x_n \ E_m) \land (\exists x_1 \forall y_1 \dots \exists x_n \forall y_n \ O_{m+1})$$

is preserved by direct factors, so is

$$(\exists y_1 \forall x_1 \dots \exists y_n \forall x_n \ E_m) \wedge (\exists x_1 \forall y_1 \dots \exists x_n \forall y_n \ O_{m-1}).$$

(2) For
$$4 \le m \le N$$
, m even, if

$$(\exists y_1 \forall x_1 \dots \exists y_n \forall x_n \ E_m) \wedge (\exists x_1 \forall y_1 \dots \exists x_n \forall y_n \ O_{m-1})$$

is preserved by direct factors, so is

$$(\exists y_1 \forall x_1 \dots \exists y_n \forall x_n \ E_{m-2}) \wedge (\exists x_1 \forall y_1 \dots \exists x_n \forall y_n \ O_{m-1}).$$

Note that every subindex varies over words of length less than or equal to N, so an expression of the form " $\bigwedge_{\gamma \neq \varepsilon} \tau_{\alpha \gamma}$ " should be read as " $\bigwedge \{\tau_{\alpha \gamma} : \gamma \neq \varepsilon \text{ and } |\alpha \gamma| \leq N\}$ ". Therefore, if $m \geq N$, $O_m = true$ (empty conjunction) and $E_N = \bigwedge_{|\beta|=N} \tau_{\beta}$. Also, recall that the i-th component of an element a in a direct product $\Pi_i A_i$ is called a^i .

We will now state and prove two lemmas that will be helpful in order to prove Lemma A.1. In the following we will assume that the tuple \vec{z} has length equal to 1, since proofs are exactly the same and this simplification makes them easier to read.

Lemma A.2. Let m be an even integer, $A_0, A_1 \in \mathcal{V}$ and $c, d, e, a_1, \ldots, a_{2n} \in A_0 \times A_1$ such that $2 \leq m \leq N$, $A_0 \times A_1 \models E_m(c, d, e, a_1, \ldots, a_{2n})$ and $A_1 \models O_{m+1}(c^1, d^1, e^1, a_1^1, \ldots, a_{2n}^1)$. Then $A_0 \models E_m(c^0, d^0, e^0, a_1^0, \ldots, a_{2n}^0)$ and if α has length m, then

(A.1)
$$A_0 \models \left(\bigwedge_{\gamma \neq \varepsilon} \tau_{\alpha\gamma}\right)(c^0, d^0, e^0, a_1^0, \dots, a_{2n}^0) \Rightarrow A_0 \times A_1 \models \left(\bigwedge_{\mu} \tau_{\alpha\mu}\right)(c, d, e, a_1, \dots, a_{2n}).$$

Proof. By induction on m. If m = N, the first part is immediate since E_N is a conjunction of formulas preserved by taking direct factors and hence preserved by taking direct factors. The second part is contained in the hypothesis.

To make the proof more readable, we will omit the string of parameters. Take an even m such that $2 \le m < N$ and suppose the lemma is proved for m+2. Assume (A.2)

$$A_0 \times A_1 \models E_m$$
, noting that $E_m = E_{m+2} \wedge \bigwedge_{|\alpha|=m} \Big(\bigwedge_{\gamma \neq \varepsilon} \tau_{\alpha\gamma} \to \tau_{\alpha} \Big);$

(A.3)

$$A_1 \models O_{m+1}, \quad \text{noting that} \quad O_{m+1} = O_{m+3} \land \bigwedge_{|\alpha| = m+1} \Big(\bigwedge_{\gamma \neq \varepsilon} \tau_{\alpha\gamma} \to \tau_{\alpha} \Big).$$

By the first part of the inductive hypothesis we thus have $A_0 \models E_{m+2}$. We have to see that $A_0 \models \bigwedge_{|\alpha|=m} (\bigwedge_{\gamma \neq \varepsilon} \tau_{\alpha\gamma} \to \tau_{\alpha})$. Suppose now that for some α of length m

$$(A.4) A_0 \models \bigwedge_{\gamma \neq \varepsilon} \tau_{\alpha\gamma}.$$

In particular, for each $i, j \leq N$ we have

$$A_0 \models \bigwedge_{\gamma \neq \varepsilon} \tau_{\alpha i j \gamma}.$$

We will prove that $A_0 \models \tau_{\alpha}$, and the second part of the lemma will be proved along the way.

By the second part of the inductive hypothesis (i.e., (A.1)) we obtain, for all j,

$$A_0 \times A_1 \models \bigwedge_{\mu} \tau_{\alpha i j \mu}$$
 or, in other symbols, $A_0 \times A_1 \models \bigwedge_{\gamma \neq \varepsilon} \tau_{\alpha i \gamma}$.

Since this last formula is preserved by taking direct factors, we have

$$(A.5) A_1 \models \bigwedge_{\gamma \neq \varepsilon} \tau_{\alpha i \gamma}.$$

Using (A.3) (note $|\alpha i| = m+1$), we have $A_1 \models \tau_{\alpha i}$ for all i. This, together with (A.5) yields $A_1 \models \bigwedge_{\gamma \neq \varepsilon} \tau_{\alpha \gamma}$. Now we apply (A.4), thus obtaining

$$A_0 \times A_1 \models \bigwedge_{\gamma \neq \varepsilon} \tau_{\alpha \gamma}.$$

Applying (A.2),

$$A_0 \times A_1 \models \tau_{\alpha}$$
.

The last two formulas jointly say

(A.6)
$$A_0 \times A_1 \models \left(\bigwedge_{\mu} \tau_{\alpha\mu} \right).$$

We have proved (A.4) \Rightarrow (A.6), which is the second conclusion. Since τ_{α} is preserved by taking direct factors, we obtain $A_0 \models \tau_{\alpha}$, which is the first conclusion.

Lemma A.3. Let m be an even integer, $A_0, A_1 \in \mathcal{V}$ and $c, d, e, a_1, \ldots, a_{2n} \in A_0 \times A_1$ such that $2 \leq m \leq N$, $A_0 \times A_1 \models O_{m-1}(c, d, e, a_1, \ldots, a_{2n})$ and $A_1 \models E_m(c^1, c^1, e^1, a_1^1, \ldots, a_{2n}^1)$. Then $A_0 \models O_{m-1}(c^0, d^0, e^0, a_1^0, \ldots, a_{2n}^0)$ and if α has length m-1, then

$$A_0 \models \left(\bigwedge_{\gamma \neq \varepsilon} \tau_{\alpha \gamma} \right) (c^0, d^0, e^0, a_1^0, \dots, a_{2n}^0) \Rightarrow A_0 \times A_1 \models \left(\bigwedge_{\mu} \tau_{\alpha \mu} \right) (c, d, e, a_1, \dots, a_{2n}).$$

Proof. By induction on m. If m = N, the hypotheses are:

$$(A.7) A_0 \times A_1 \models O_{N-1} = \bigwedge_{|\alpha|=N-1} \Big(\bigwedge_{\gamma \neq \varepsilon} \tau_{\alpha\gamma} \to \tau_{\alpha}\Big),$$

(A.8)
$$A_1 \models E_N = \bigwedge_{|\beta|=N} \tau_{\beta}.$$

Assume that for some α of length N-1,

$$A_0 \models \bigwedge_{\gamma \neq \varepsilon} \tau_{\alpha \gamma} = \bigwedge_i \tau_{\alpha i}.$$

By using (A.8) and preservation by direct products, we have

$$A_0 \times A_1 \models \bigwedge \tau_{\alpha i}$$
.

By (A.7) we have $A_0 \times A_1 \models \tau_{\alpha}$, thus obtaining

$$A_0 \times A_1 \models \bigwedge_{\mu} \tau_{\alpha\mu}.$$

We have proved the second part of the lemma. Passing to A_0 we obtain the first part.

Now take an even m such that $2 \le m < N$ and suppose the lemma is proved for m+2. Assume

$$A_0 \times A_1 \models O_{m-1} = O_{m+1} \wedge \bigwedge_{|\alpha|=m-1} \left(\bigwedge_{\gamma \neq \varepsilon} \tau_{\alpha\gamma} \to \tau_{\alpha} \right),$$

$$A_1 \models E_m = E_{m+2} \wedge \bigwedge_{|\alpha|=m} \left(\bigwedge_{\gamma \neq \varepsilon} \tau_{\alpha\gamma} \to \tau_{\alpha} \right).$$

By the inductive hypothesis we thus have $A_0 \models O_{m+1}$. The rest of the argument closely parallels the proof of Lemma A.2.

Proof of Lemma A.1. To see (1), suppose

$$A_0 \times A_1 \models \Big(\big(\exists y_1 \forall x_1 \dots \exists y_n \forall x_n \ E_m \big) \land \big(\exists x_1 \forall y_1 \dots \exists x_n \forall y_n \ O_{m-1} \big) \Big) (c, d, e).$$

Thus we have (Skolem) functions G_1, \ldots, G_n such that G_i is (i-1)-ary and

$$A_0 \times A_1 \models \forall \vec{y} O_{m-1}(c, d, e, G_1, y_1, \dots, G_n(y_1, \dots, y_{n-1}), y_n).$$

Since O_{m-1} implies O_{m+1} , we have

$$A_0 \times A_1 \models \Big(\big(\exists y_1 \forall x_1 \dots \exists y_n \forall x_n \ E_m \big) \land \big(\exists x_1 \forall y_1 \dots \exists x_n \forall y_n \ O_{m+1} \big) \Big) (c, d, e).$$

Then, since this formula is preserved by taking direct factors, by hypothesis,

(A.9)
$$A_0 \models (\exists y_1 \forall x_1 \dots \exists y_n \forall x_n \ E_m)(c^0, d^0, e^0),$$

(A.10)
$$A_1 \models (\exists y_1 \forall x_1 \dots \exists y_n \forall x_n \ E_m)(c^1, d^1, e^1).$$

Thus we have functions F_1, \ldots, F_n such that

$$A_1 \models \forall \vec{x} E_m(c^1, d^1, e^1, x_1, F_1, \dots, x_n, F_n(x_1, \dots, x_{n-1})).$$

Now, for $j=1,\ldots,n,$ define j-ary functions p_j from A_0 to $A_0\times A_1,$ $p_j=p_j(a_1,\ldots,a_j)$:

$$p_1 := (a_1, F_1),$$

$$p_2 := (a_2, F_2(G_1^1)),$$

$$p_j := (a_j, F_j(G_1^1, G_2(p_1)^1, \dots, G_{j-1}(p_1, \dots, p_{j-2})^1)).$$

The reader may check that this selection ensures, for each $\vec{a} \in A_0^n$:

$$A_0 \times A_1 \models O_{m-1}(c, d, e, G_1, p_1, \dots, G_n(p_1, \dots, p_{n-1}), p_n),$$

$$A_1 \models E_m(c^1, d^1, e^1, G_1^1, p_1^1, \dots, G_n(p_1, \dots, p_{n-1})^1, p_n^1).$$

We may apply Lemma A.3 and obtain

$$A_0 \models O_{m-1}(c^0, d^0, e^0, G_1^0, p_1^0, \dots, G_n(p_1, \dots, p_{n-1})^0, p_n^0).$$

Equivalently.

$$A_0 \models O_{m-1}(c^0, d^0, e^0, G_1^0, a_1, G_2(p_1)^0, a_2, \dots, G_n(p_1, \dots, p_{n-1})^0, a_n).$$

Now, defining $H_j: A_0^{j-1} \to A_0$ as follows:

$$H_1 := G_1^0,$$

$$H_2(y_1) := G_2(p_1(y_1))^0,$$

$$H_j(y_1, \dots, y_{j-1}) := G_j(p_1(y_1), \dots, p_{j-1}(y_1, \dots, y_{j-1}))^0,$$

we see at once that

$$A_0 \models \forall \vec{y} O_{m-1}(c^0, d^0, e^0, H_1, y_1, \dots, H_n(y_1, \dots, y_{n-1}), y_n),$$

and then

$$A_0 \models (\exists x_1 \forall y_1 \dots \exists x_n \forall y_n \ O_{m-1})(c^0, d^0, e^0).$$

This, together with (A.9), proves this case.

Part (2) is entirely analogous to the former, and it's proved by using Lemma A.2.

The hypotheses of the next theorem are the same as in Lemma A.1, and we repeat them for ease of reference.

Theorem A.4. For every word α in the alphabet $\{1, \ldots, N\}$ of length no greater than N, let $\tau_{\alpha} = \tau_{\alpha}(x, y, \vec{z}, x_1, y_1, \ldots, x_n, y_n)$ be a formula preserved by taking direct products and direct factors. Define:

$$E_m := \bigwedge_{\substack{m \leq |\alpha| \leq N \\ |\alpha| \ even}} \bigg(\bigwedge_{\gamma \neq \varepsilon} \tau_{\alpha\gamma} \to \tau_{\alpha} \bigg), \qquad O_m := \bigwedge_{\substack{m \leq |\alpha| \leq N \\ |\alpha| \ odd}} \bigg(\bigwedge_{\gamma \neq \varepsilon} \tau_{\alpha\gamma} \to \tau_{\alpha} \bigg).$$

Then the formula

$$(A.11) \qquad (\exists y_1 \forall x_1 \dots \exists y_n \forall x_n \ E_2) \land (\exists x_1 \forall y_1 \dots \exists x_n \forall y_n \ O_1)$$

is preserved by taking direct factors and direct products.

Proof. First observe that

$$(\exists y_1 \forall x_1 \dots \exists y_n \forall x_n \ E_N) \land (\exists x_1 \forall y_1 \dots \exists x_n \forall y_n \ O_{N+1}) = \exists y_1 \forall x_1 \dots \exists y_n \forall x_n \ \bigwedge_{|\beta|=N} \tau_{\beta}$$

is preserved by direct factors. This is immediate since the conjunction and quantification of formulas preserved by taking direct factors is again preserved by taking direct factors. Successive application of Lemma A.1 yields that (A.11) is preserved by taking direct factors.

The proof that (A.11) is preserved by direct products is a straightforward calculation. \Box

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