# XSS-GUARD: Precise Dynamic Prevention of Cross Site Scripting (XSS) Attacks

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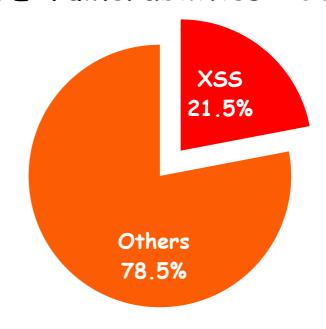


#### XSS attacks: number one threat

#### CVE Vulnerabilities 2004

# XSS 10.9% Others 89.1%

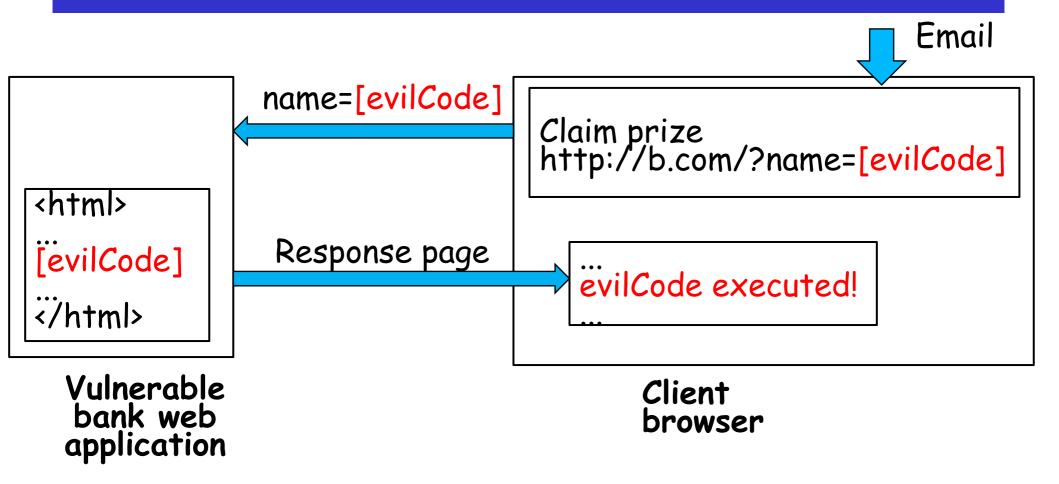
#### CVE Vulnerabilities 2006



- ...and the trend continues...
  - Second half of 2007: 80% of all attacks were XSS
  - January 2007: 70% web applications are vulnerable
     [source: http://en.wikipedia.org]
- Simple attacks lucrative targets
  - <script> alert('xss');</script>



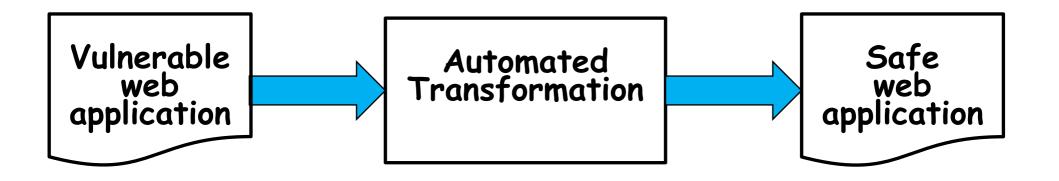
# A typical XSS attack



 Attacker controlled code can steal sensitive information or perform malicious operations.



# Objective



- Automated prevention of XSS attacks: server side
- Robust against subtle attacks
- · Efficient

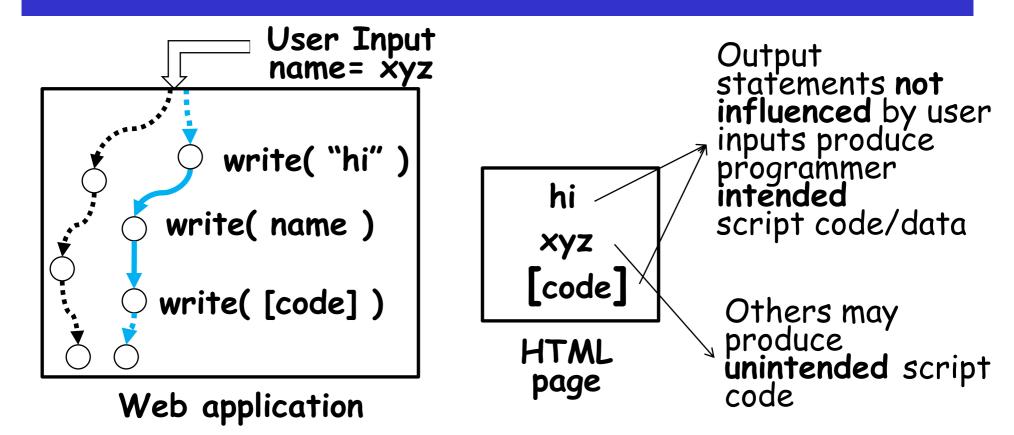


#### Outline of this talk

- Introduction
- Web application transformation technique
- Robust script identification at server side
- · XSS-GUARD
  - Examples
  - Evaluation results
- Related work and summary



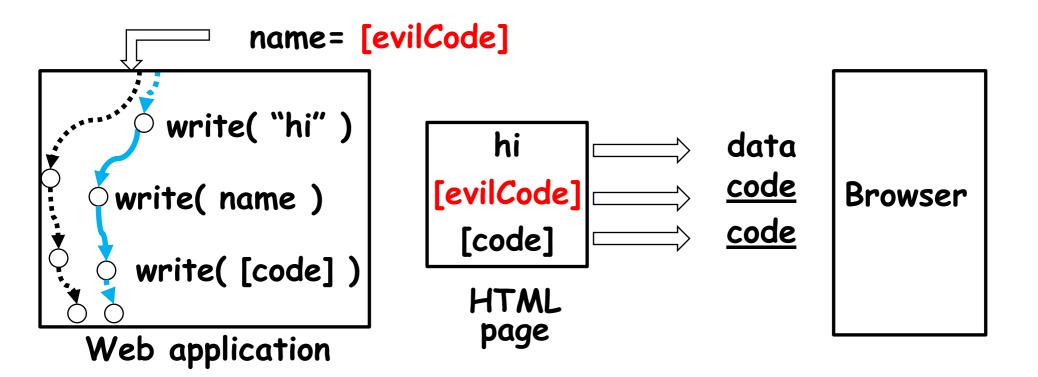
# HTML page: A web application's view



- Page generated by output statements in a control path.
- Web application's view: intended regions & others
  - Other regions could lead to unintended script code.



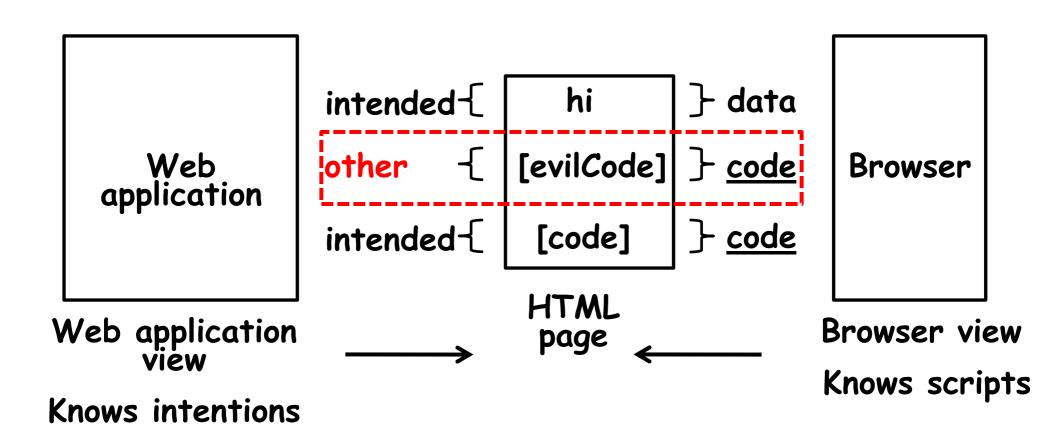
## HTML page: A browser's view



- Browser does not differentiate between injected and programmer intended scripts.
- Browser's view: a collection of script code & data.



### A complete view



· An effective defense would require both these views!



#### Idea

If a web application knows intended scripts and all the scripts (including injected) for a generated HTML page, it can remove unintended scripts.

Question: How to compute intended scripts?



# Computing intended code

```
name= xyz
       name c = aaa
                                     hi
                                              data
                                                       Real
   write( realPage, "hi" )
                                                      HTML
                                              data
                                    XYZ
  write ( shadow Page, "hi" )
                                                       page
                                   [code]
                                              code
write( realPage, name )
write( shadowPage, name_c )
                                     hi
                                               data
                                                      Shadow
  write( realPage, [code] )
                                               data
                                    aga
write( shadowPage, [code] )
                                                       page
                                   [code]
                                               code
```

# Web application

- Replicate output statements uninfluenced by user inputs to create **shadow page**.
- Other output statements replicated but act on benign inputs (as intended).



## Computing intended code...cont.

```
name = [xssCode]

name_c = aaaaaaaa

write( realPage, "hi" )

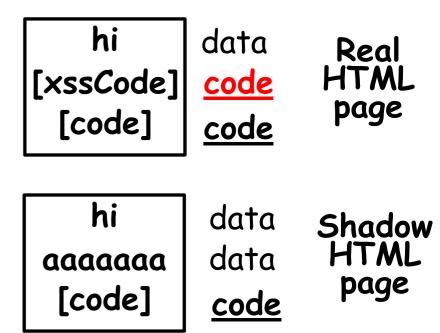
write( shadowPage, "hi" )

write( realPage, name )

write( shadowPage, name_c )

write( realPage, [code] )

write( shadowPage, [code] )
```



#### Web application

- Real page contains injected script, but shadow retains only intended script.
- For a real page, its shadow page has intended scripts.



# Shadow page captures intended code

- Real HTML page = output statements with user inputs
- Shadow HTML page = mirror above output statements with benign user inputs
- Transform web application to create shadow (intended) page for each real page
  - Define "benign input" for each "real input".
  - Mirror the "actual input" processing on "benign input".
  - Replicate output statements with above processed inputs.

#### For details on transformation, please refer to

 CANDID: Preventing SQL Injection Attacks using Dynamic Candidate Evaluations, S. Bandhakavi, P. Bisht, P. Madhusudan, V.N. Venkatakrishnan, ACM CCS 2007, submitted to ACM TISSEC 2008



#### Idea Revisited

If a web application knows

intended scripts

and

all scripts (including injected)
for a generated HTML page, it can
remove unintended scripts.

Question: How to compute intended scripts?

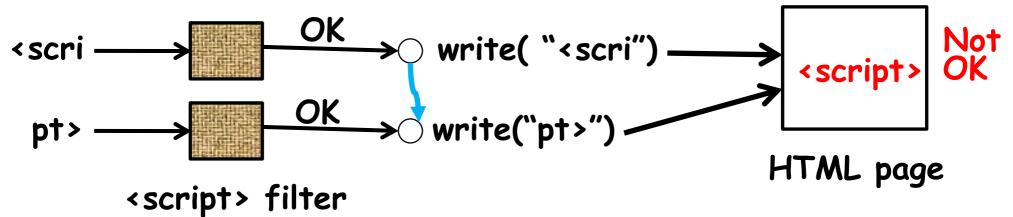
- By computing shadow pages.

Question: How can application identify all scripts? What about filters?



# Filters effective first layer...but lack context

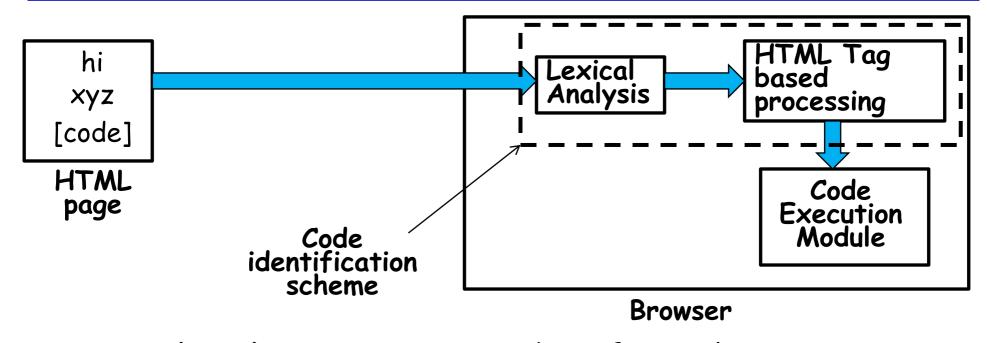
- Ineffective against subtle cases
  - MySpace Samy worm used eval('inner' + 'HTML') to evade "innerHTML" filter.
- Large attack surface
  - tags and URI schemes, attributes, event handlers, alternate encoding...
- Filters analyze inputs without their context of use.



- Alternate scheme: find scripts in output (HTML page)
  - Inputs embedded in context of use in HTML page



# How Firefox identifies scripts?

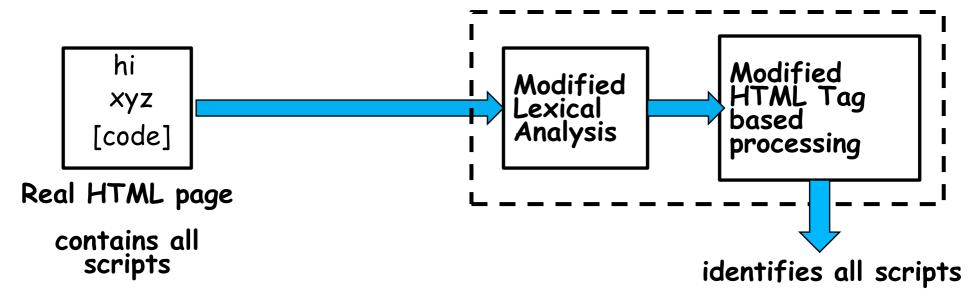


- Lexical analysis component identifies tokens.
- HTML tag based processing identifies scripts in:
  - External resource download e.g., <script src=...>
  - Inlined scripts/event handlers e.g., <body onload=...>
  - URI schemes that can have scripts e.g., javascript/data



# Leveraging browser's code identification mechanism

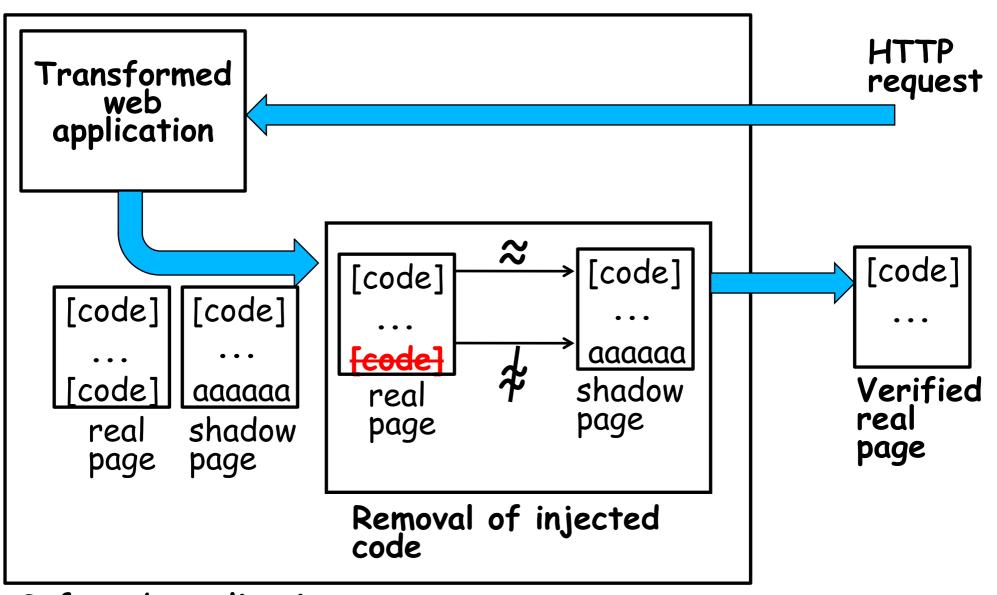
- A browser performs precise identification of scripts.
- Robust
  - alternate encodings
  - large attack surface
- Our approach leverages this at the server side.



- Modifications record all scripts in real HTML page.



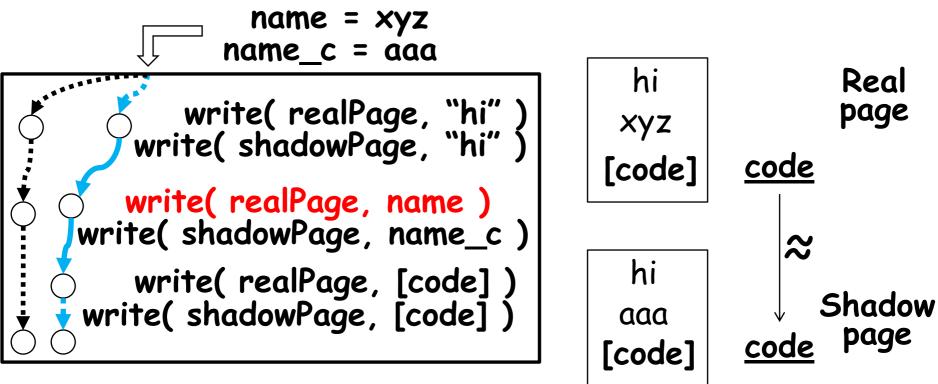
#### XSS-GUARD: End - to - End



Safe web application

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## XSS-GUARD: intended scripts



#### Web application

 All intended scripts in real page have equivalent scripts in shadow page.



## XSS-GUARD: Attack prevention

```
name = <script>...</script>
                                                    Real
         name_c = aaaaaaaaaaaaaaaa
                                                    page
                                                    hi
     write( realPage, "hi" )
write( shadowPage, "hi" )
                                         <script> </script>
                                                              code
                                                  [code]
    write( realPage, name )
 write( shadowPage, name_c )
                                                    hi
     write( realPage, [code] )
                                                               data
                                          aaaaaaaaaaaaaa
 write( shadowPage, [code] )
                                                  [code]
Web application
                                                   Shadow
```

 Injected script in real page does not have equivalent script in shadow page, and is removed.



page

#### XSS-GUARD: Subtle attack case

```
name = aa;evil(...);
                                                              Real
              name_c = aaaaaaaaaa
                                                              page
                                           <script> x = aa;
evil(...);</script>
  write( realPage, "<script>x = " )
                                                                     code
write( shadowPage, "<script>x=" )
   write( realPage, name )
                                           <script> x =
aaaaaaaa
write( shadowPage, name_c )
                                                                     code
  write( realPage, "</script>" )
write( shadowPage, "</script>" )
                                                             Shadow
                                                               page
```

 Any unintended addition to existing scripts is successfully prevented.

web application



#### Effectiveness Evaluation

Against real world exploits

CVE-2007-5120/5121	JSPWiki	defended
CVE-2007-2450	Tomcat HTML Manager	defended
CVE-2007-3386	Tomcat Host Manager	defended
CVE-2007-3383/3384/2449 CVE-2006-7196	Tomcat Example Web Applications	defended

- Defended 32 applicable exploits out of 92: R. Hansen XSS cheatsheet.
- False negatives: non-Firefox attacks
  - Current implementation can be extended
  - initial experiments: Defended 35 / 56 non-Firefox attacks



#### Performance Evaluation

Performance overhead (response time)

Exploits from CVE	5 - 24 %
Varied response sizes (1KB - 75KB)	3 - 14 %
Parse tree comparison of scripts (1 - 5)	37 - 42 %

- Parse tree comparison is rarely done: in presence of attacks, or scripts embedding user inputs.
- These numbers indicate worst case performance -
  - Negligible network latency in experiments (LAN setup)
  - Can be further improved by limiting the transformation to only relevant statements.



#### Some Related Work

- Vulnerability analysis: find vulnerable source-sink pairs e.g., saner: Livshits et al. Usenix 2005, Pixy N. Jovanovic et al. S&P2006, Y. Xie et al. Usenix 2006, D. Balzarotti et al. CCS 2007...
  - Useful but limited to detection
- Server side solutions: filter based or track taint & disallow at sink: W. Xu et al. Usenix 2006, ...
  - Centralized defense but do not know all scripts
- Client side solutions: Firewall like mechanisms to prevent malicious actions at client
  - Noxes E. Kirda, et al. SAC 2006, P. Vogt et al. NDSS 2007
  - User controlled protection but do not know intended scripts
- Client-Server collaborative solutions: Clients enforce application specified policies
  - BEEP T. Jim, et al. WWW 2007, Tahoma R. Cox et al. S&P 2006, Browsershield C. Reis et al. OSDI 2006
  - Can determine intended and all scripts but deployment issues



#### Contributions and future work

- A robust server side solution to prevent XSS attacks.
- A mechanism to compute programmer intended code, useful in defending other code injection attacks.
- Leveraged browser's mechanisms at server side.

Thanks for your attention!

Questions?



# Backup Slides



#### Taint vs Candidate evaluations

- Taint tracking captures "trust" notion.
- Candidate computation captures "intended structure".
- We found taint tracking and limiting tainted constructs in parse trees, a very powerful idea. There are no differences in effectiveness but there are subtle differences -
- X = taintedX taintedX + 100.5;
  - X will be treated as tainted: false positive?
- X\_c = taintedX\_c taintedX\_c + 100.5;
  - X\_c will contain equivalent structure to X.



## Miscellaneous examples

- Non-Firefox construct based false negative
  - <IMG SRC='vbscript:msgbox("XSS")'> Firefox does not understand vbscript URIs.
- Firefox quirk based exploit
  - <script/xss src=""> </script>



# Changes done to handle non-Firefox quirks

- Others (3)
- Attacks based on src attributes (2)
  - <XML SRC="xsstest.xml" ID=I></XML> <SPAN DATASRC=#I DATAFLD=C DATAFORMATAS=HTML></SPAN>
  - <SCRIPT a=`>`
    SRC="http://ha.ckers.org/xss.js"></SCRIPT>
- Javascript / data URI scheme based exploits (30)
  - <img src="javascript:alert('XSS')">Is ignored in Firefox, but valid in other browsers.
  - We forced all the attributes to be parsed irrespective of Firefox applicability.



### Candidate Transformation



# Performance numbers (detailed)



#### Related work

- Filters: don't have adequate context
- Output encoding: may forbid all HTML
- Taint: Effective, focuses on taintedness, rather than semantics
- Server side solutions :
- Client side solutions: Firewall like behavior, may disallow legitimate scripts, or allow attacks on trusted servers.
- Server-client side solutions



# Real page and shadow page comparison with offsets



# False positive in subtle case



# Example with existing filters - prevent attack

